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Operations on Attributes

The Domain Algebra

- Horizontal (scalar) operations
 - Arithmetic
 - Constants
 - Renaming
- Vertical (aggregation) operations
 - Reduction & Equivalence reduction
 - Functional mapping & Partial functional mapping

Domain Algebra (principle of closure)

I Horizontal operations

e.g., **let** *TotMk* **be** *Asst* + *Exam*;

| <i>NewMk</i> | | | |
|------------------|-------------|---------------|--------------|
| (<i>Student</i> | <i>Asst</i> | <i>Exam</i>) | <i>TotMk</i> |
| Smith | 25. | 60. | 85. |
| Jones | 28. | 66. | 94. |
| Brown | 20. | 50. | 70. |
| Hung | 24. | 58. | 82. |
| Raman | 24. | 66. | 90. |

N.B. *Virtual* domains (principle of abstraction)

Actualization via Project

```
let TotMk be Asst + Exam;  
Result ← [ Student, TotMk ] in NewMk;
```

Any expression allowed that can be actualized on each tuple separately.

Domain Algebra

Frequently Asked Questions

- What if *Asst* and *Exam* are also in some other relation?

Nothing is affected: TotMk could be actualized there, too.

- What if *Asst* and *Exam* each come from different relations?

Use relational algebra to put the relations together before actualizing.

- What if we update *NewMk* after actualizing *TotMk*?

TotMk is an actual attribute of Result and unaffected by changes to other relations.

Domain Algebra

Constants

let *One* **be** 1;

Renaming

let *Final* **be** *Exam*;

Course < – [*Student*, *Asst*, *Final*, *TotMk*,
One] **in** *NewMk*;

| <i>Course</i> | | | | |
|------------------|-------------|--------------|--------------|--------------|
| (<i>Student</i> | <i>Asst</i> | <i>Final</i> | <i>TotMk</i> | <i>One</i>) |
| Smith | 25. | 60. | 85. | 1 |
| Jones | 28. | 66. | 94. | 1 |
| Brown | 20. | 50. | 70. | 1 |
| Hung | 24. | 58. | 82. | 1 |
| Raman | 24. | 66. | 90. | 1 |

Domain Algebra

II Vertical operations e.g. of *Reduction*

let *Total* **be** *red + of TotMk*;

let *Count* **be** *red + of One*;

let *AvgMk* **be** *Total/Count*;

(let *AvgMk* **be** *(red + of TotMk)/(red + of 1)*;))

| <i>Course</i> (<i>Student</i> | <i>Asst</i> | <i>Final</i> | <i>TotMk</i> | <i>One</i>) | <i>Total</i> | <i>Count</i> | <i>AvgMk</i> |
|-----------------------------------|-------------|--------------|--------------|--------------|--------------|--------------|--------------|
| Smith | 25. | 60. | 85. | 1 | 421 | 5 | 84.2 |
| Jones | 28. | 66. | 94. | 1 | 421 | 5 | 84.2 |
| Brown | 20. | 50. | 70. | 1 | 421 | 5 | 84.2 |
| Hung | 24. | 58. | 82. | 1 | 421 | 5 | 84.2 |
| Raman | 24. | 66. | 90. | 1 | 421 | 5 | 84.2 |

Vertical operations

Operators: $+$, \times , **max**, **min**, **and**, **or**

e.g., **let** *MaxMk* **be red max of** *TotMk*;

| <i>Course</i> | | | | | |
|------------------|-------------|--------------|--------------|--------------|--------------|
| (<i>Student</i> | <i>Asst</i> | <i>Final</i> | <i>TotMk</i> | <i>One</i>) | <i>MaxMk</i> |
| Smith | 25. | 60. | 85. | 1 | 94 |
| Jones | 28. | 66. | 94. | 1 | 94 |
| Brown | 20. | 50. | 70. | 1 | 94 |
| Hung | 24. | 58. | 82. | 1 | 94 |
| Raman | 24. | 66. | 90. | 1 | 94 |

Of course, the attributes outside the (..) are *virtual*, so there is no “waste”: the programmer actualizes and controls the waste.

Class \leftarrow [*AvgMk*, *MaxMK*] **in** *Course*;

| <i>Course</i> | |
|----------------|----------------|
| (<i>AvgMk</i> | <i>MaxMk</i>) |
| 84.2 | 94 |

Domain Algebra

II Vertical operations

e.g. of *Equivalence Reduction*

let $STot$ be equiv + of Mark by $S\#$;

let $CTot$ be equiv + of Mark by $C\#$;

StuCour

| $(S\#$ | $C\#$ | Mark) | $STot$ | $CTot$ |
|--------|-------|-------|--------|--------|
| 1 | 1 | 73. | 233. | 213. |
| 1 | 2 | 82. | 233. | 147. |
| 1 | 3 | 78. | 233. | 78. |
| 2 | 1 | 64. | 64. | 213. |
| 3 | 1 | 76. | 141. | 213. |
| 3 | 2 | 65. | 141. | 147. |

Domain algebra—Vertical

e.g. of *Functional Mapping*

let Cum be fun + of Amount order Year

Sales

| <i>(President</i> | <i>Year</i> | <i>Amount)</i> | <i>Cum</i> |
|-------------------|-------------|----------------|------------|
| Smith | 1994 | 150 | 150 |
| Smith | 1995 | 175 | 325 |
| Smith | 1996 | 200 | 525 |
| Brown | 1996 | 200 | 525 |
| Brown | 1997 | 210 | 735 |
| Brown | 1998 | 225 | 960 |

Vertical operations

Operators: $+$, $-$, \times , \div , **max**, **min**, **||**, **pred**, **succ**

e.g., **let** *Alt* **be fun** – **of** *Num* **order** *Seq*;

| <i>(Num</i> | <i>Seq)</i> | <i>Alt</i> |
|-------------|-------------|------------|
| 27 | 1 | 27 |
| 13 | 2 | -14 |
| 12 | 3 | 26 |
| 16 | 4 | -10 |
| 9 | 5 | 19 |

Vertical operations

| <i>Text</i> | | <i>Seq)</i> | <i>Next</i> |
|---------------|--|-------------|--------------|
| (<i>Word</i> | | | |
| Algebraic | | 1 | data |
| data | | 2 | processing |
| processing | | 3 | techniques |
| techniques | | 4 | can |
| can | | 5 | enable |
| enable | | 6 | applications |
| applications | | 7 | programmers |
| programmers | | 8 | to |
| to | | 9 | work |
| work | | 10 | with |
| with | | 11 | units |
| units | | 12 | of |
| of | | 13 | data |
| data | | 14 | larger |
| larger | | 15 | than |
| than | | 16 | a |
| a | | 17 | single |
| single | | 18 | computer |
| computer | | 19 | word |
| word | | 20 | Algebraic |

let *Next* be fun succ of *Word* order *Seq*;

Domain algebra—Vertical

e.g. of *Partial Functional Mapping*

let *DCum* be par + of *Amount* order *Year* by *Div*;

DivSales

| <i>(Div</i> | <i>Year</i> | <i>Amount)</i> | <i>DCum</i> |
|-------------|-------------|----------------|-------------|
| A | 1997 | 80 | 80 |
| A | 1998 | 110 | 190 |
| B | 1997 | 60 | 60 |
| B | 1998 | 75 | 135 |
| C | 1997 | 90 | 90 |
| C | 1998 | 110 | 200 |

Combining Relational and Domain Algebras

e.g., Matrix multiplication

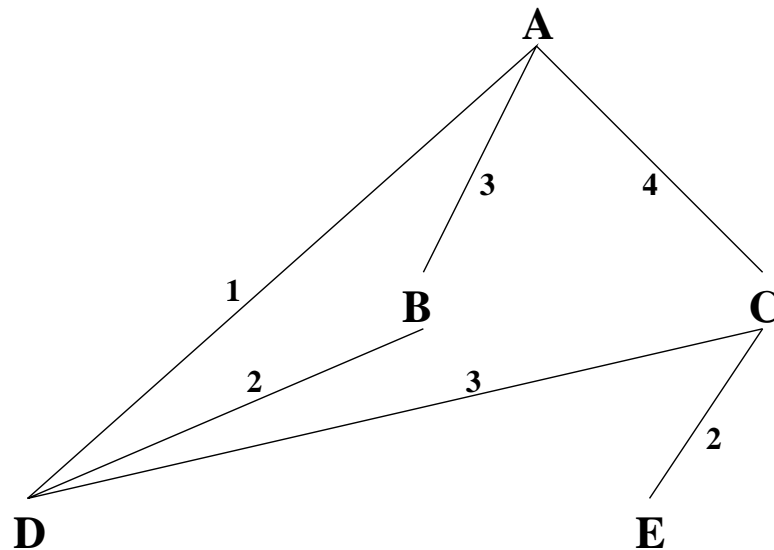
| $A(\text{Val}A \quad I \quad J)$ | $B(\text{Val}B \quad J \quad K)$ | $AB(\text{Val}AB \quad I \quad K)$ |
|----------------------------------|----------------------------------|------------------------------------|
| 1 | 1 | 3 |
| 2 | 1 | 5 |
| 2 | 2 | 2 |
| 1 | 2 | 3 |
| 3 | 3 | 2 |
| 1 | 3 | 1 |
| | | 4 |
| | | 5 |
| | | 1 |

$$A \begin{pmatrix} 1 & 0 & 2 \\ 0 & 2 & 1 \\ 3 & 0 & 1 \end{pmatrix} B \begin{pmatrix} 1 & 1 & 0 \\ 1 & 0 & 0 \\ 1 & 2 & 1 \end{pmatrix} AB \begin{pmatrix} 3 & 5 & 2 \\ 3 & 2 & 1 \\ 4 & 5 & 1 \end{pmatrix}$$

let $ValAB$ be equiv + of $ValA \times ValB$ by I, K ;
 $AB \leftarrow [ValAB, I, K]$ in $(A \text{ ijoin } B)$;

| $A \text{ ijoin } B$ | | | | | | | |
|----------------------|-----|-----|-----|---------|--------------------|---------|--|
| $(ValA$ | I | J | K | $ValB)$ | $ValA \times ValB$ | $ValAB$ | |
| 1 | 1 | 1 | 1 | 1 | 1 | 3 | |
| 1 | 1 | 1 | 2 | 1 | 1 | : | |
| 3 | 3 | 1 | 1 | 1 | 3 | 4 | |
| 3 | 3 | 1 | 2 | 1 | 3 | : | |
| 2 | 2 | 2 | 1 | 1 | 2 | | |
| 2 | 1 | 3 | 1 | 1 | 2 | 3 | |
| 2 | 1 | 3 | 2 | 2 | 4 | : | |
| 2 | 1 | 3 | 3 | 1 | 2 | | |
| 1 | 2 | 3 | 1 | 1 | 1 | | |
| 1 | 2 | 3 | 2 | 2 | 2 | | |
| 1 | 2 | 3 | 3 | 1 | 1 | | |
| 1 | 3 | 3 | 1 | 1 | 1 | 4 | |
| 1 | 3 | 3 | 2 | 2 | 2 | : | |
| 1 | 3 | 3 | 3 | 1 | 1 | | |

Bill of Materials



| <i>PartOf</i> (A | <i>S</i> | <i>Q</i>) | <i>Explo</i> (A | <i>S</i> | <i>Q</i>) | <i>Q''</i> | <i>Q'''</i> | | |
|------------------|----------|------------|-----------------|----------|------------|------------|-------------|---|---|
| A | B | 3 | A | B | 3 | | 3 | | |
| A | C | 4 | A | C | 4 | | 4 | | |
| A | D | 1 | A | D | 1 | 18 | 19 | | |
| B | D | 2 | B | D | 2 | | 2 | | |
| C | D | 3 | C | D | 3 | | 3 | | |
| C | E | 2 | C | E | 2 | | 2 | | |
| | | | | | | A | E | 8 | 8 |

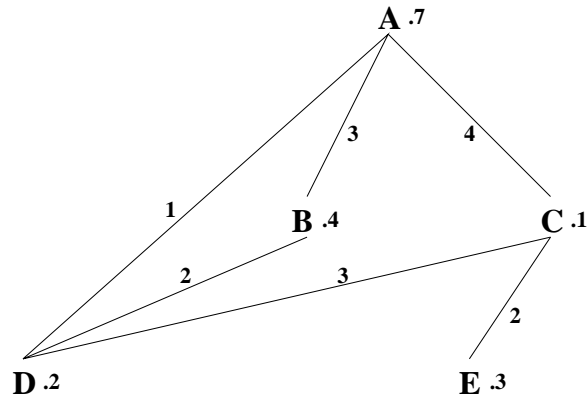
let A' be A ; let S' be S ; let Q' be Q ;

let Q'' be equiv + of $Q \times Q'$ by A, S' ;

let Q''' be $Q + Q''$; let Q be Q''' ;

Explo is $[A, S, Q]$ in $[A, S, Q''']$ in (*PartOf* $[A, S$ ujoin $A, S']$
 $[A, S', Q'']$ in (*Explo* $[S$ ijoin $A']$ $[A', S', Q']$ in *PartOf*));

Bill of Materials with Costs



| <i>Cost(A C)</i> | <i>Explo(A S Q)</i> | <i>C</i> | <i>QC</i> | <i>QCEA</i> |
|------------------|---------------------|----------|-----------|-------------|
| A .7 | A B 3 | .4 | 1.2 | 7.8 |
| B .4 | A C 4 | .1 | .4 | 7.8 |
| C .1 | A D 19 | .2 | 3.8 | 7.8 |
| D .2 | A E 8 | .3 | 2.4 | 7.8 |
| E .3 | B D 2 | .2 | .4 | .4 |
| | C D 3 | .2 | .6 | 1.2 |
| | C E 2 | .3 | .6 | 1.2 |

let A' be A ; let QC be $Q \times C$;

let $QCEA$ be equiv + of $Q \times C$ by A ;

let $Tcost$ be $QCEA + C$;

$SubCost \leftarrow [A, QCEA]$ in ($Explo [S \text{ ijoin } A'] [A', C]$ in $Cost$);

$FinCost \leftarrow [A, Tcost]$ in ($SubCost \text{ ujoin } Cost$);

| <i>SubCost ujoin Cost(A QCEA C)</i> | <i>Tcost</i> |
|-------------------------------------|--------------|
| A 7.8 .7 | 8.5 |
| B .4 .4 | .8 |
| C 1.2 .1 | 1.3 |
| D .2 | .2 |
| E .3 | .3 |