## Probability in Programming

Prakash Panangaden School of Computer Science McGill University

• Correctness of software:

- Correctness of software:
- with respect to precise specifications.

- Correctness of software:
- with respect to precise specifications.
- Efficiency of code:

- Correctness of software:
- with respect to precise specifications.
- Efficiency of code:
- based on well-designed algorithms.

 Precise specifications have to be made in a formal language

- Precise specifications have to be made in a formal language
- with a rigourous definition of meaning.

- Precise specifications have to be made in a formal language
- with a rigourous definition of meaning.
- Logic is the "calculus" of computer science.

- Precise specifications have to be made in a formal language
- with a rigourous definition of meaning.
- Logic is the "calculus" of computer science.
- It comes with a framework for reasoning.

- Precise specifications have to be made in a formal language
- with a rigourous definition of meaning.
- Logic is the "calculus" of computer science.
- It comes with a framework for reasoning.
- Many kinds of logic: propositional, predicate, modal, .....

• Probability is also a framework for reasoning

- Probability is also a framework for reasoning
- quantitatively.

- Probability is also a framework for reasoning
- quantitatively.
- But is this relevant for computer programmers?

- Probability is also a framework for reasoning
- quantitatively.
- But is this relevant for computer programmers?
- · Yes!

- Probability is also a framework for reasoning
- quantitatively.
- But is this relevant for computer programmers?
- · Yes!
- Probabilistic reasoning is everywhere.

#### Some quotations

### Some quotations

• The true logic of the world is the calculus of probabilities — James Clerk Maxwell

### Some quotations

- The true logic of the world is the calculus of probabilities James Clerk Maxwell
- The theory of probabilities is at bottom nothing but common sense reduced to calculus — Pierre Simon Laplace

• Some algorithms use probability as a computational resource: randomized algorithms.

- Some algorithms use probability as a computational resource: randomized algorithms.
- Software for interacting with physical systems have to cope with noise and uncertainty: telecommunications, robotics, vision, control systems, ....

- Some algorithms use probability as a computational resource: randomized algorithms.
- Software for interacting with physical systems have to cope with noise and uncertainty: telecommunications, robotics, vision, control systems, ....
- Big data and machine learning: probabilistic reasoning has had a revolutionary impact.

Sample space X: the set of things that can possibly happen.

Sample space X: the set of things that can possibly happen.

Event: subset of the sample space;  $A, B \subset X$ .

Sample space X: the set of things that can possibly happen.

Event: subset of the sample space;  $A, B \subset X$ .

Probability:  $\Pr: X \to [0, 1], \sum_{x \in X} \Pr(x) = 1.$ 

Sample space X: the set of things that can possibly happen.

Event: subset of the sample space;  $A, B \subset X$ .

Probability:  $\Pr: X \to [0, 1], \sum_{x \in X} \Pr(x) = 1.$ 

Probability of an event A:  $\Pr(A) = \sum_{x \in A} \Pr(x)$ .

Sample space X: the set of things that can possibly happen.

Event: subset of the sample space;  $A, B \subset X$ .

Probability:  $\Pr: X \to [0, 1], \sum_{x \in X} \Pr(x) = 1.$ 

Probability of an event A:  $\Pr(A) = \sum_{x \in A} \Pr(x)$ .

A, B are independent:  $\Pr(A \cap B) = \Pr(A) \cdot \Pr(B)$ .

Sample space X: the set of things that can possibly happen.

Event: subset of the sample space;  $A, B \subset X$ . Probability:  $\Pr: X \to [0, 1], \sum_{x \in X} \Pr(x) = 1$ . Probability of an event A:  $\Pr(A) = \sum_{x \in A} \Pr(x)$ .

A, B are independent:  $\Pr(A \cap B) = \Pr(A) \cdot \Pr(B)$ .

Subprobability:  $\sum_{x \in X} \Pr(x) \le 1$ .

Imagine a town where every birth is equally likely to give a boy or a girl.  $Pr(boy) = Pr(girl) = \frac{1}{2}$ .

Imagine a town where every birth is equally likely to give a boy or a girl.  $Pr(boy) = Pr(girl) = \frac{1}{2}$ .

Each birth is an *independent* random event.

Imagine a town where every birth is equally likely to give a boy or a girl.  $Pr(boy) = Pr(girl) = \frac{1}{2}$ .

Each birth is an *independent* random event.

There is a family with two children.

#### A Puzzle

Imagine a town where every birth is equally likely to give a boy or a girl.  $Pr(boy) = Pr(girl) = \frac{1}{2}$ .

Each birth is an *independent* random event.

There is a family with two children.

One of them is a boy (not specified which one), what is the probability that the other one is a boy?

#### A Puzzle

Imagine a town where every birth is equally likely to give a boy or a girl.  $Pr(boy) = Pr(girl) = \frac{1}{2}$ .

Each birth is an *independent* random event.

There is a family with two children.

One of them is a boy (not specified which one), what is the probability that the other one is a boy?

Since the births are independent, the probability that the other child is a boy should be  $\frac{1}{2}$ . Right?

Wrong!

Wrong!

Initially, there are 4 *equally likely* situations: bb, bg, gb, gg.

Wrong!

Initially, there are 4 *equally likely* situations: bb, bg, gb, gg.

The possibility gg is ruled out with the additional information.

Wrong!

Initially, there are 4 *equally likely* situations: bb, bg, gb, gg.

The possibility gg is ruled out with the additional information.

So of the three *equally likely* scenarios: bb, bg, gb, only one has the other child being a boy.

Wrong!

Initially, there are 4 *equally likely* situations: bb, bg, gb, gg.

The possibility gg is ruled out with the additional information.

So of the three *equally likely* scenarios: bb, bg, gb, only one has the other child being a boy.

The correct answer is  $\frac{1}{3}$ .

Wrong!

Initially, there are 4 *equally likely* situations: bb, bg, gb, gg.

The possibility gg is ruled out with the additional information.

So of the three *equally likely* scenarios: bb, bg, gb, only one has the other child being a boy.

The correct answer is  $\frac{1}{3}$ .

If I had said, "The *elder* child is a boy", then the probability that the other child is a boy is indeed  $\frac{1}{2}$ 

Conditioning = revising probability in the presence of new information.

Conditioning = revising probability in the presence of new information.

Conditional probability/expectation is *the* heart of probabilistic reasoning.

Conditioning = revising probability in the presence of new information.

Conditional probability/expectation is *the* heart of probabilistic reasoning.

Conditional probability is tricky!

Conditioning = revising probability in the presence of new information.

Conditional probability/expectation is *the* heart of probabilistic reasoning.

Conditional probability is tricky!

Analogous to *inference* in ordinary logic.

Definition: if A and B are events, the **conditional probability** of A given B, written Pr(A | B) is defined by

 $\Pr(A \mid B) = \Pr(A \cap B) / \Pr(B).$ 

Definition: if A and B are events, the **conditional probability** of A given B, written Pr(A | B) is defined by

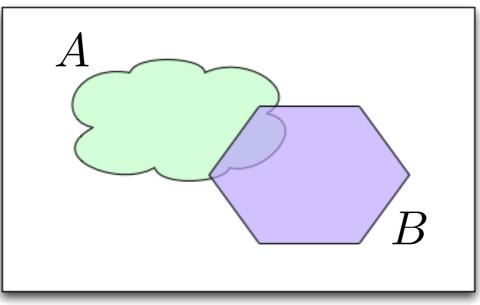
 $\Pr(A \mid B) = \Pr(A \cap B) / \Pr(B).$ 

We are told that the outcome is one of the possibilities in B. We now need to change our guess for the outcome being in A.

Definition: if A and B are events, the **conditional probability** of A given B, written Pr(A | B) is defined by

 $\Pr(A \mid B) = \Pr(A \cap B) / \Pr(B).$ 

We are told that the outcome is one of the possibilities in B. We now need to change our guess for the outcome being in A.



#### Bayes' Rule

# $\Pr(A \mid B) = \frac{\Pr(B \mid A) \cdot \Pr(A)}{\Pr(B)}.$

How to revise probabilities.

#### Bayes' Rule

# $\Pr(A \mid B) = \frac{\Pr(B \mid A) \cdot \Pr(A)}{\Pr(B)}.$

How to revise probabilities.

Proof is just from the definition.

Two coins, one fake (two heads) one OK.

Two coins, one fake (two heads) one OK.

One coin chosen with equal probability and then tossed to yield a H.

Two coins, one fake (two heads) one OK.

One coin chosen with equal probability and then tossed to yield a H.

What is the probability the coin was fake?

Two coins, one fake (two heads) one OK.

One coin chosen with equal probability and then tossed to yield a H.

What is the probability the coin was fake? Answer:  $\frac{2}{3}$ .

Two coins, one fake (two heads) one OK.

One coin chosen with equal probability and then tossed to yield a H.

What is the probability the coin was fake? Answer:  $\frac{2}{3}$ .

$$\Pr(H \mid \text{Fake}) = 1, \Pr(\text{Fake}) = \frac{1}{2}, \Pr(H) = \frac{1}{2} \cdot 1 + \frac{1}{2} \cdot \frac{1}{2} = \frac{3}{4}.$$

Two coins, one fake (two heads) one OK.

One coin chosen with equal probability and then tossed to yield a H.

What is the probability the coin was fake? Answer:  $\frac{2}{3}$ .

$$\Pr(H \mid \text{Fake}) = 1, \Pr(\text{Fake}) = \frac{1}{2}, \Pr(H) = \frac{1}{2} \cdot 1 + \frac{1}{2} \cdot \frac{1}{2} = \frac{3}{4}.$$

Hence  $\Pr(\text{Fake} \mid H) = (\frac{1}{2})/(\frac{3}{4}) = \frac{2}{3}$ .

Two coins, one fake (two heads) one OK.

One coin chosen with equal probability and then tossed to yield a H.

What is the probability the coin was fake? Answer:  $\frac{2}{3}$ .

$$\Pr(H \mid \text{Fake}) = 1, \Pr(\text{Fake}) = \frac{1}{2}, \Pr(H) = \frac{1}{2} \cdot 1 + \frac{1}{2} \cdot \frac{1}{2} = \frac{3}{4}.$$

Hence  $\Pr(\text{Fake} \mid H) = (\frac{1}{2})/(\frac{3}{4}) = \frac{2}{3}$ .

Similarly  $\Pr(\text{Fake} \mid HHH) = \frac{8}{9}$ .

Two coins, one fake (two heads) one OK.

One coin chosen with equal probability and then tossed to yield a H.

What is the probability the coin was fake? Answer:  $\frac{2}{3}$ .

$$\Pr(H \mid \text{Fake}) = 1, \Pr(\text{Fake}) = \frac{1}{2}, \Pr(H) = \frac{1}{2} \cdot 1 + \frac{1}{2} \cdot \frac{1}{2} = \frac{3}{4}$$

Hence  $\Pr(\text{Fake} \mid H) = (\frac{1}{2})/(\frac{3}{4}) = \frac{2}{3}$ .

Similarly  $\Pr(\text{Fake} \mid HHH) = \frac{8}{9}$ .  $\Pr(\text{Fake} \mid \underbrace{H \dots H}) = \frac{1}{1 + (\frac{1}{2})^n}$ .

Bayes' rule shows how to update the *prior* probability of A with the new information that the outcome was B: this gives the *posterior* probability of A given B.

A random variable r is a real-valued function on X.

A random variable r is a real-valued function on X.

The expectation value of r is  $\mathbb{E}[r] = \sum_{x \in X} \Pr(x) r(x).$ 

A random variable r is a real-valued function on X.

The expectation value of r is 
$$\mathbb{E}[r] = \sum_{x \in X} \Pr(x) r(x).$$

The conditional expectation value of r given A is:

A random variable r is a real-valued function on X.

The expectation value of r is 
$$\mathbb{E}[r] = \sum_{x \in X} \Pr(x) r(x).$$

The conditional expectation value of r given A is:

$$\mathbb{E}[r \mid A] = \sum_{x \in X} r(x) \mathsf{Pr}(\{x\} \mid A).$$

A random variable r is a real-valued function on X.

The expectation value of r is 
$$\mathbb{E}[r] = \sum_{x \in X} \Pr(x) r(x).$$

The conditional expectation value of r given A is:

$$\mathbb{E}[r \mid A] = \sum_{x \in X} r(x) \mathsf{Pr}(\{x\} \mid A).$$

Conditional probability is a special case of conditional expectation.

Two people roll dice independently.

Two people roll dice independently.



Two people roll dice independently.

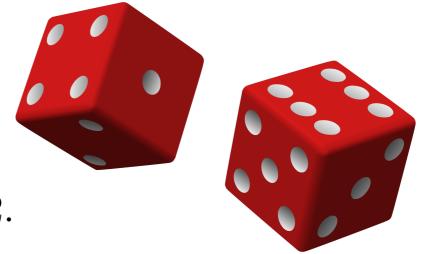
The first one keeps rolling until he gets a 1 *immediately* followed by a 2.



Two people roll dice independently.

The first one keeps rolling until he gets a 1 *immediately* followed by a 2.

The second one keeps rolling until she gets a 1 and a 1.

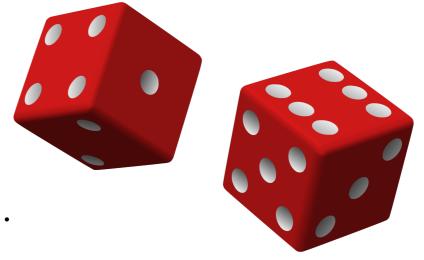


Two people roll dice independently.

The first one keeps rolling until he gets a 1 *immediately* followed by a 2.

The second one keeps rolling until she gets a 1 and a 1.

Do they have the same *expected* number of rolls?



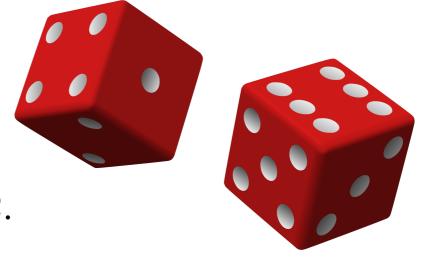
Two people roll dice independently.

The first one keeps rolling until he gets a 1 *immediately* followed by a 2.

The second one keeps rolling until she gets a 1 and a 1.

Do they have the same *expected* number of rolls?

If not, who is expected to finish first?



Two people roll dice independently.

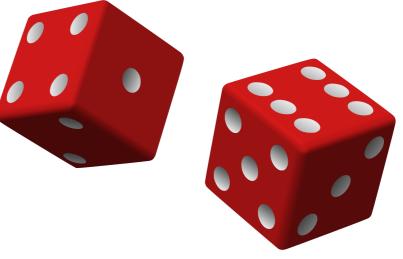
The first one keeps rolling until he gets a 1 *immediately* followed by a 2.

The second one keeps rolling until she gets a 1 and a 1.

Do they have the same *expected* number of rolls?

If not, who is expected to finish first?

What is the expected number of rolls for each one?



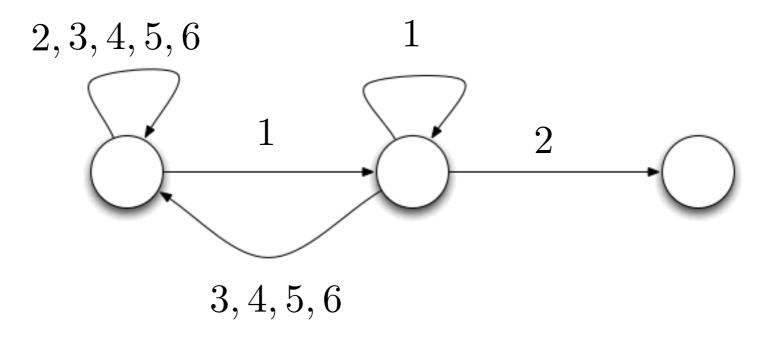
Is there a better way?

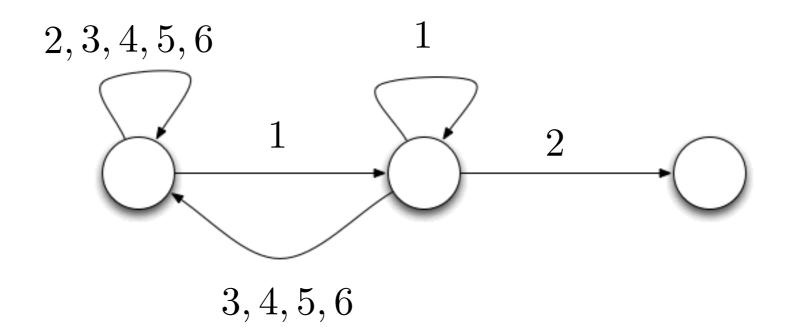
Is there a better way?

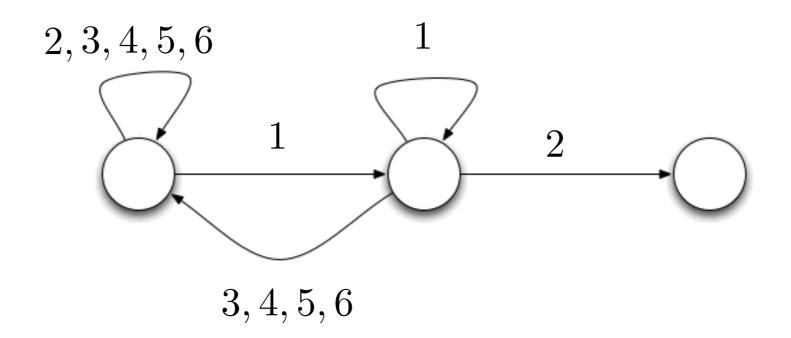
Use conditional expectations and think in terms of state-transition diagrams:

Is there a better way?

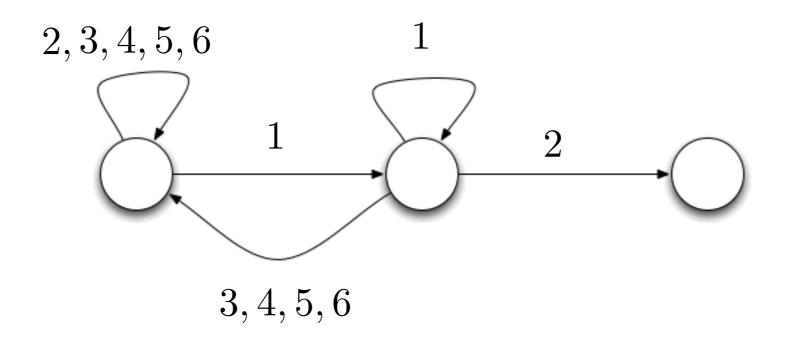
Use conditional expectations and think in terms of state-transition diagrams:



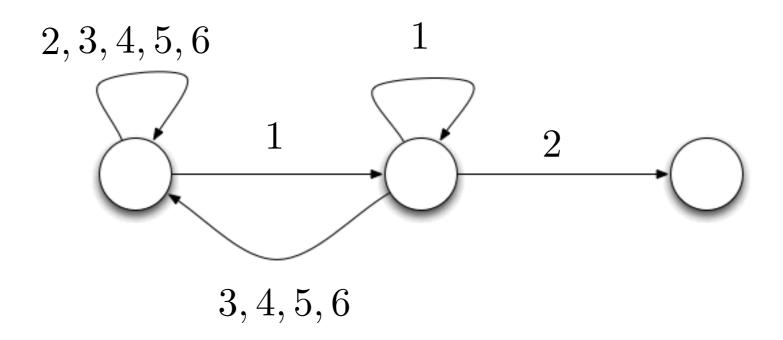




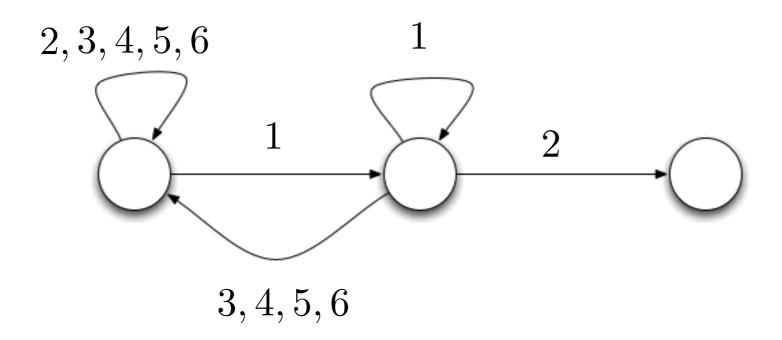
#### Let $x = \mathbb{E}[\text{Finish} \mid \text{Start}]$



#### Let $x = \mathbb{E}[\text{Finish} \mid \text{Start}]$ Let $y = \mathbb{E}[\text{Finish} \mid 1]$

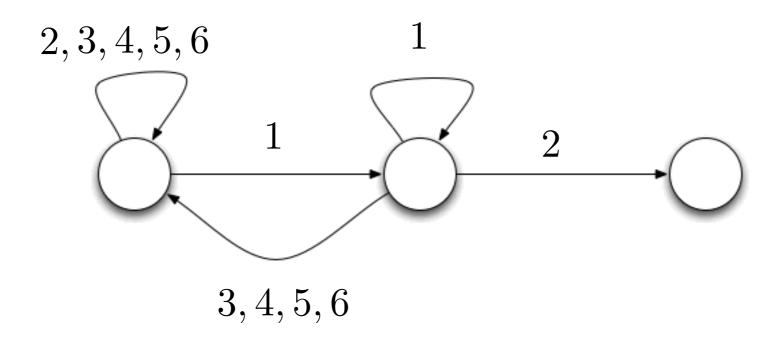


Let  $x = \mathbb{E}[\text{Finish} \mid \text{Start}]$  Let  $y = \mathbb{E}[\text{Finish} \mid 1]$  $x = \frac{5}{6} \cdot (1+x) + \frac{1}{6} \cdot (1+y)$ 



Let  $x = \mathbb{E}[\text{Finish} \mid \text{Start}]$  Let  $y = \mathbb{E}[\text{Finish} \mid 1]$ 

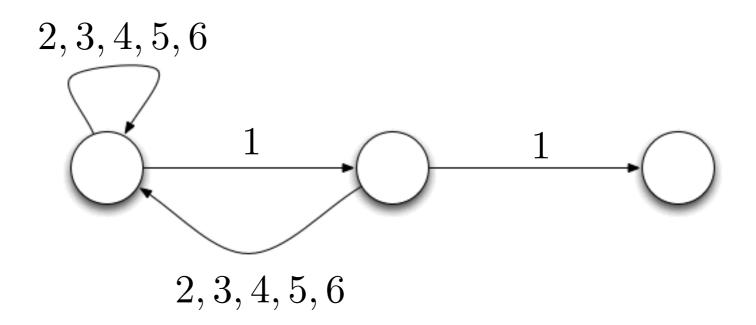
$$x = \frac{5}{6} \cdot (1+x) + \frac{1}{6} \cdot (1+y)$$
$$y = \frac{1}{6} \cdot 1 + \frac{1}{6} \cdot (1+y) + \frac{2}{3} \cdot (1+x)$$

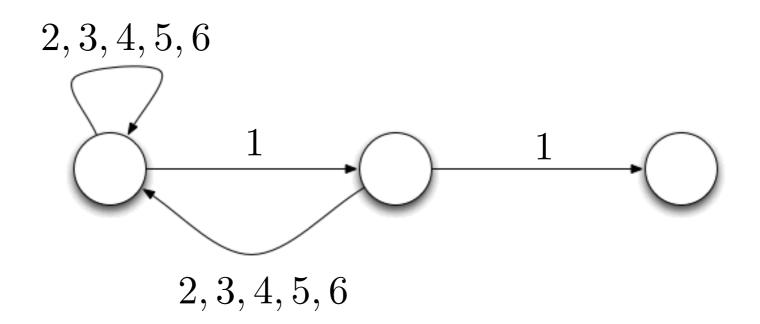


Let  $x = \mathbb{E}[\text{Finish} \mid \text{Start}]$  Let  $y = \mathbb{E}[\text{Finish} \mid 1]$ 

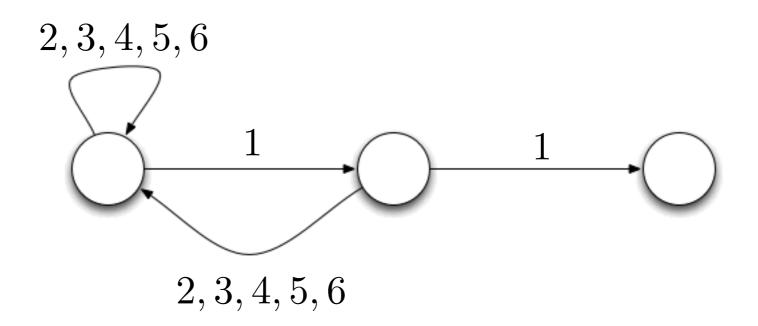
$$x = \frac{5}{6} \cdot (1+x) + \frac{1}{6} \cdot (1+y)$$
$$y = \frac{1}{6} \cdot 1 + \frac{1}{6} \cdot (1+y) + \frac{2}{3} \cdot (1+x)$$

Easy to solve: x = 30, y = 36.

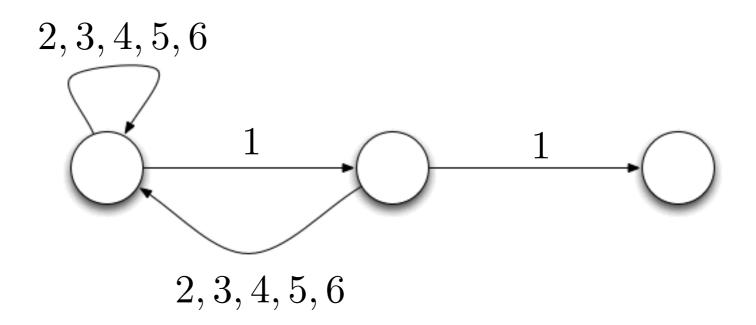




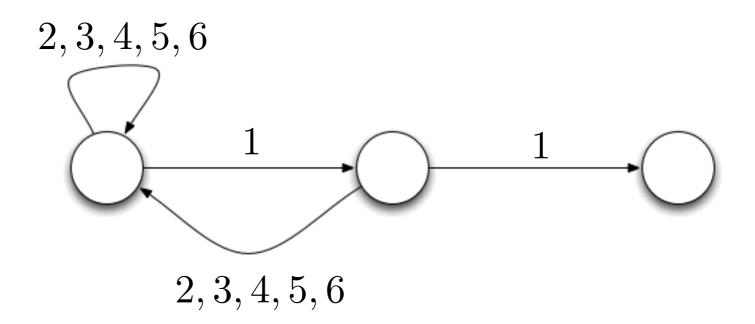
#### Let $x = \mathbb{E}[\text{Finish} \mid \text{Start}]$



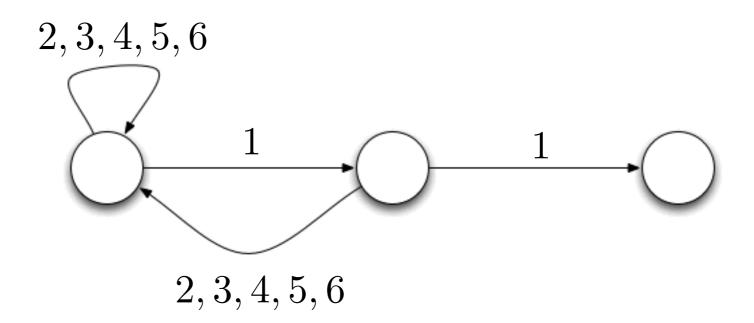
Let  $x = \mathbb{E}[\text{Finish} \mid \text{Start}]$  Let  $y = \mathbb{E}[\text{Finish} \mid 1]$ 



Let  $x = \mathbb{E}[\text{Finish} \mid \text{Start}]$  Let  $y = \mathbb{E}[\text{Finish} \mid 1]$  $x = \frac{1}{6} \cdot (1+y) + \frac{5}{6} \cdot (1+x)$ 

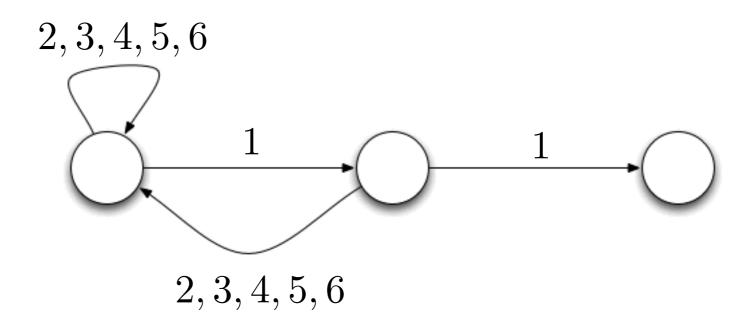


Let  $x = \mathbb{E}[\text{Finish} \mid \text{Start}]$  Let  $y = \mathbb{E}[\text{Finish} \mid 1]$  $x = \frac{1}{6} \cdot (1+y) + \frac{5}{6} \cdot (1+x)$  $y = \frac{1}{6} \cdot 1 + \frac{5}{6} \cdot (1+x)$ 



Let  $x = \mathbb{E}[\text{Finish} \mid \text{Start}]$  Let  $y = \mathbb{E}[\text{Finish} \mid 1]$  $x = \frac{1}{6} \cdot (1+y) + \frac{5}{6} \cdot (1+x)$  $y = \frac{1}{6} \cdot 1 + \frac{5}{6} \cdot (1+x)$ 

Easy to solve: x = 42, y = 36.



Let  $x = \mathbb{E}[\text{Finish} \mid \text{Start}]$  Let  $y = \mathbb{E}[\text{Finish} \mid 1]$  $x = \frac{1}{6} \cdot (1+y) + \frac{5}{6} \cdot (1+x)$  $y = \frac{1}{6} \cdot 1 + \frac{5}{6} \cdot (1+x)$ 

Easy to solve: x = 42, y = 36.

Did you expect this to be the slower one?

The *state* of a program is the correspondence between names and values.

The *state* of a program is the correspondence between names and values.  $[X \mapsto 3, Y \mapsto 4, Z \mapsto -2.5]$ 

The *state* of a program is the correspondence between names and values.  $[X \mapsto 3, Y \mapsto 4, Z \mapsto -2.5]$ 

Running a part of a program changes the state.

The *state* of a program is the correspondence between names and values.  $[X \mapsto 3, Y \mapsto 4, Z \mapsto -2.5]$ 

Running a part of a program changes the state.

if X > 1 then Y = Y + Z else Y = Z

The *state* of a program is the correspondence between names and values.  $[X \mapsto 3, Y \mapsto 4, Z \mapsto -2.5]$ 

Running a part of a program changes the state.

if X > 1 then Y = Y + Z else Y = Z

 $[X \mapsto 3, Y \mapsto 4, Z \mapsto -2.5] \longrightarrow [X \mapsto 3, Y \mapsto 1.5, Z \mapsto -2.5]$ 

The *state* of a program is the correspondence between names and values.  $[X \mapsto 3, Y \mapsto 4, Z \mapsto -2.5]$ 

Running a part of a program changes the state.

if X > 1 then Y = Y + Z else Y = Z

 $[X \mapsto 3, Y \mapsto 4, Z \mapsto -2.5] \longrightarrow [X \mapsto 3, Y \mapsto 1.5, Z \mapsto -2.5]$ 

Ordinary programs define state-transformer *functions*.

The *state* of a program is the correspondence between names and values.  $[X \mapsto 3, Y \mapsto 4, Z \mapsto -2.5]$ 

Running a part of a program changes the state.

if X > 1 then Y = Y + Z else Y = Z

 $[X \mapsto 3, Y \mapsto 4, Z \mapsto -2.5] \longrightarrow [X \mapsto 3, Y \mapsto 1.5, Z \mapsto -2.5]$ 

Ordinary programs define state-transformer *functions*. When one combines program pieces one can *compose* the functions to find the combined effect.

The *state* of a program is the correspondence between names and values.  $[X \mapsto 3, Y \mapsto 4, Z \mapsto -2.5]$ 

Running a part of a program changes the state.

if X > 1 then Y = Y + Z else Y = Z

 $[X \mapsto 3, Y \mapsto 4, Z \mapsto -2.5] \longrightarrow [X \mapsto 3, Y \mapsto 1.5, Z \mapsto -2.5]$ 

Ordinary programs define state-transformer *functions*. When one combines program pieces one can *compose* the functions to find the combined effect.

How do we understand probabilistic programs?

The *state* of a program is the correspondence between names and values.  $[X \mapsto 3, Y \mapsto 4, Z \mapsto -2.5]$ 

Running a part of a program changes the state.

if X > 1 then Y = Y + Z else Y = Z

 $[X \mapsto 3, Y \mapsto 4, Z \mapsto -2.5] \longrightarrow [X \mapsto 3, Y \mapsto 1.5, Z \mapsto -2.5]$ 

Ordinary programs define state-transformer *functions*. When one combines program pieces one can *compose* the functions to find the combined effect.

How do we understand probabilistic programs?

As distribution transformers.

Initial distribution:  $[X \mapsto (0, 1.0), C \mapsto (0, 1.0)]$ 

Initial distribution:  $[X \mapsto (0, 1.0), C \mapsto (0, 1.0)]$ 

Final distribution:  $[X \mapsto (1, 0.5)(-1, 0.5), C \mapsto (0, 0.5)(1, 0.5)]$ 

Initial distribution:  $[X \mapsto (0, 1.0), C \mapsto (0, 1.0)]$ 

Final distribution:  $[X \mapsto (1, 0.5)(-1, 0.5), C \mapsto (0, 0.5)(1, 0.5)]$ 

A Markov chain has S: states and a probability distribution transformer T.

Initial distribution:  $[X \mapsto (0, 1.0), C \mapsto (0, 1.0)]$ 

Final distribution:  $[X \mapsto (1, 0.5)(-1, 0.5), C \mapsto (0, 0.5)(1, 0.5)]$ 

A Markov chain has S: states and a probability distribution transformer T.

 $T: St \times St \to [0, 1] \text{ or } T: St \to Dist(St).$ 

Initial distribution:  $[X \mapsto (0, 1.0), C \mapsto (0, 1.0)]$ 

Final distribution:  $[X \mapsto (1, 0.5)(-1, 0.5), C \mapsto (0, 0.5)(1, 0.5)]$ 

A Markov chain has S: states and a probability distribution transformer T.

 $T: St \times St \to [0, 1] \text{ or } T: St \to Dist(St).$ 

 $T(s_1, s_2)$  is the conditional probability of being in state  $s_2$  after the transition given that the state was  $s_1$  before.

Initial distribution:  $[X \mapsto (0, 1.0), C \mapsto (0, 1.0)]$ 

Final distribution:  $[X \mapsto (1, 0.5)(-1, 0.5), C \mapsto (0, 0.5)(1, 0.5)]$ 

A Markov chain has S: states and a probability distribution transformer T.

 $T: St \times St \to [0, 1] \text{ or } T: St \to Dist(St).$ 

 $T(s_1, s_2)$  is the conditional probability of being in state  $s_2$  after the transition given that the state was  $s_1$  before.

Markov property: the transition probability only depends on the current state, not on the whole history.

When one combines probabilistic program pieces one can *multiply* the transition matrices to find the combined effect.

When one combines probabilistic program pieces one can *multiply* the transition matrices to find the combined effect.

We are understanding the program by stepping forwards.

When one combines probabilistic program pieces one can *multiply* the transition matrices to find the combined effect.

We are understanding the program by stepping forwards.

This is called "forwards" or state-transformer semantics.

Maybe we do not want to track every detail of the state as it changes.

Maybe we do not want to track every detail of the state as it changes.

Perhaps we want to know if a *property* holds, e.g. X > 0.

Maybe we do not want to track every detail of the state as it changes.

Perhaps we want to know if a *property* holds, e.g. X > 0.

We write  $\{P\}$  step  $\{Q\}$  to mean that P holds before the step and Q holds after the step.

Maybe we do not want to track every detail of the state as it changes.

Perhaps we want to know if a *property* holds, e.g. X > 0.

We write  $\{P\}$  step  $\{Q\}$  to mean that P holds before the step and Q holds after the step.

$$\{X > 0\} X = X - 5 \{X > 0\} ??$$

Maybe we do not want to track every detail of the state as it changes.

Perhaps we want to know if a *property* holds, e.g. X > 0.

We write  $\{P\}$  step  $\{Q\}$  to mean that P holds before the step and Q holds after the step.

$$\{X > 0\} X = X - 5 \{X > 0\} ??$$

We cannot say anything for sure after the step!

$$\{X > 5\} X = X - 5 \{X > 0\}$$

$${X > 5} X = X - 5 {X > 0}$$

$${X > 5} X = X - 5 {X > 0}$$

We must read this differently: if we want X > 0after the step, we must make sure X > 5 before the step.

$${X > 5} X = X - 5 {X > 0}$$

We must read this differently: if we want X > 0after the step, we must make sure X > 5 before the step.

This is called "predicate-transformer" semantics.

$${X > 5} X = X - 5 {X > 0}$$

We must read this differently: if we want X > 0after the step, we must make sure X > 5 before the step.

This is called "predicate-transformer" semantics.

Can we do something like this for probabilistic programs?

Classical logic	Generalization
Truth values $\{0,1\}$	Probabilities [0, 1]
Predicate	Random variable
State	Distribution
The satisfaction relation $\models$	Integration $\int$

Suppose that your probabilistic program describes a search for some resource.

Suppose that your probabilistic program describes a search for some resource.

Suppose that  $r: St \to \mathbb{R}$  is the "expected reward" in each state.

Suppose that your probabilistic program describes a search for some resource.

Suppose that  $r: St \to \mathbb{R}$  is the "expected reward" in each state.

We write  $\hat{T}r$  for a new reward function defined by  $(\hat{T}r)(s) = \sum_{s' \in S} T(s, s')r(s').$ 

Suppose that your probabilistic program describes a search for some resource.

Suppose that  $r: St \to \mathbb{R}$  is the "expected reward" in each state.

We write 
$$\hat{T}r$$
 for a new reward function defined by  $(\hat{T}r)(s) = \sum_{s' \in S} T(s, s')r(s').$ 

This tells you the expected reward *before* the transition assuming that r is the reward after the transition.

• A general framework to reason about situations computationally and quantitatively.

- A general framework to reason about situations computationally and quantitatively.
- Most important class of models: graphical models.

- A general framework to reason about situations computationally and quantitatively.
- Most important class of models: graphical models.
- They capture dependence and independence and

- A general framework to reason about situations computationally and quantitatively.
- Most important class of models: graphical models.
- They capture dependence and independence and
- conditional independence.

## Complex systems

# Complex systems

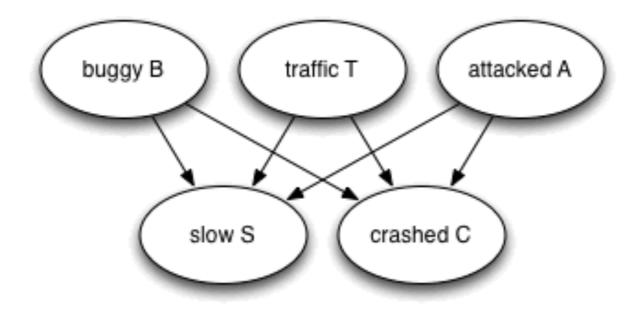
• Complex systems have many variables.

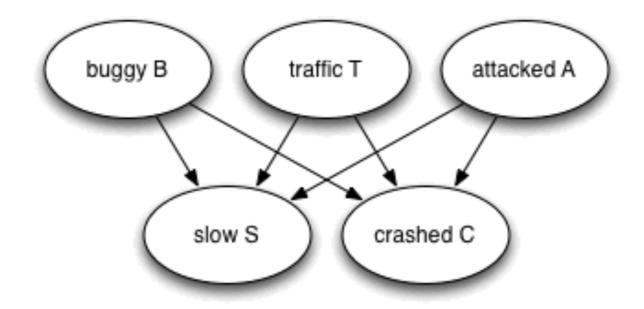
- Complex systems have many variables.
- They are related in intricate ways: correlated or independent.

- Complex systems have many variables.
- They are related in intricate ways: correlated or independent.
- They may be conditionally independent or dependent.

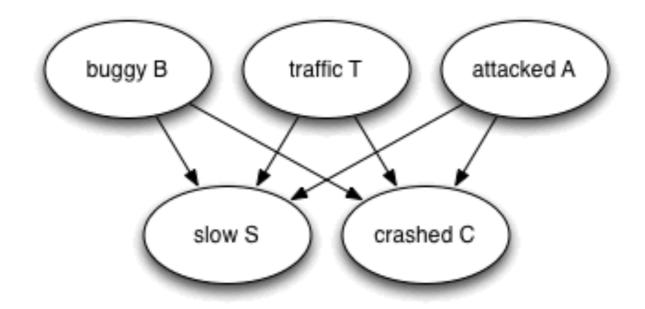
- Complex systems have many variables.
- They are related in intricate ways: correlated or independent.
- They may be conditionally independent or dependent.
- There may be causal connections.

- Complex systems have many variables.
- They are related in intricate ways: correlated or independent.
- They may be conditionally independent or dependent.
- There may be causal connections.
- We want to represent several random variables that may be connected in different ways.

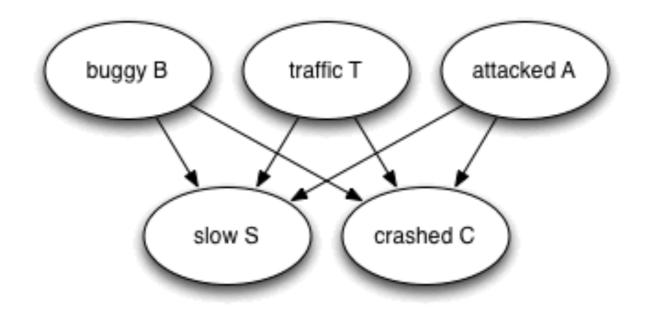




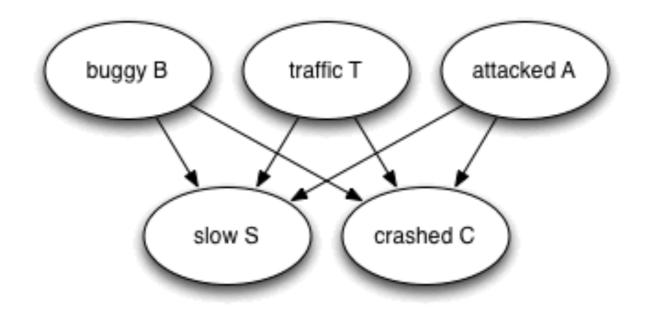
Trying to analyze what is wrong with a web site:



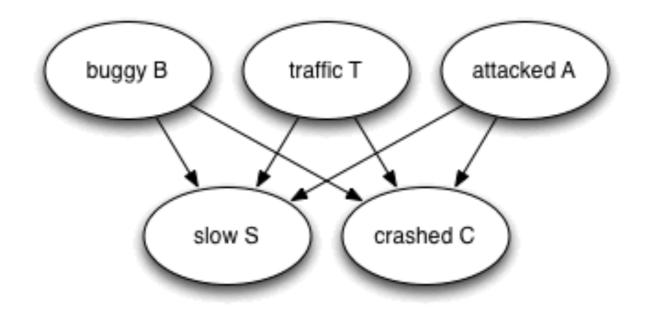
Trying to analyze what is wrong with a web site: could be buggy (yes/no)



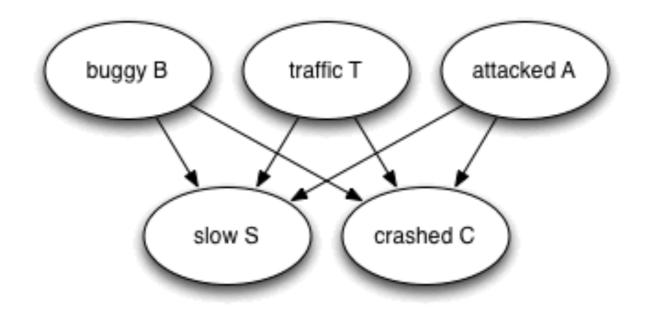
Trying to analyze what is wrong with a web site: could be buggy (yes/no) could be under attack (yes/no)



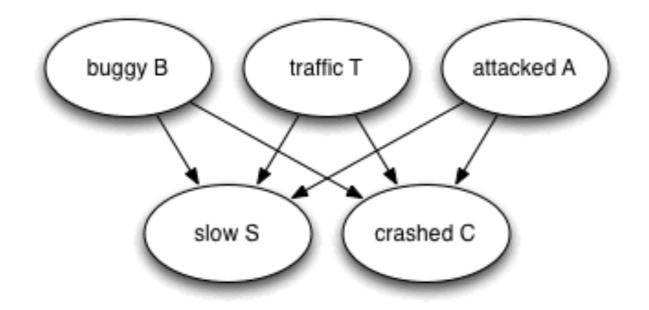
Trying to analyze what is wrong with a web site: could be buggy (yes/no) could be under attack (yes/no) traffic (very high/high/moderate/low)

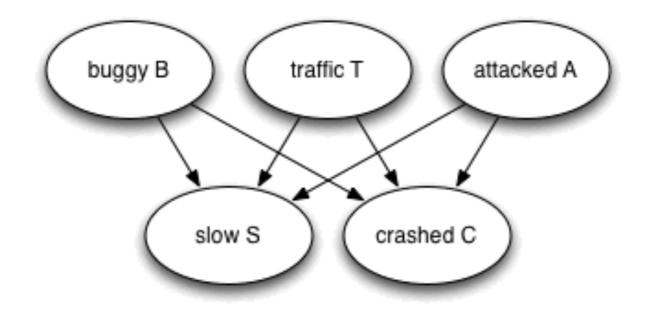


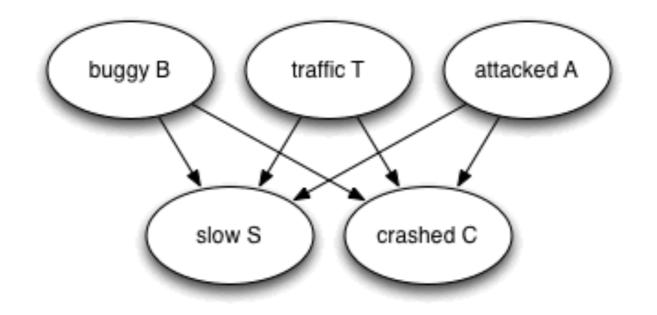
Trying to analyze what is wrong with a web site: could be buggy (yes/no) could be under attack (yes/no) traffic (very high/high/moderate/low) Symptoms: crashed or slow



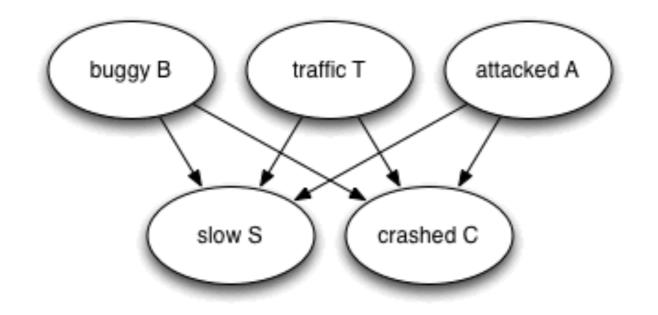
Trying to analyze what is wrong with a web site: could be buggy (yes/no) could be under attack (yes/no) traffic (very high/high/moderate/low) Symptoms: crashed or slow There are 64 states.





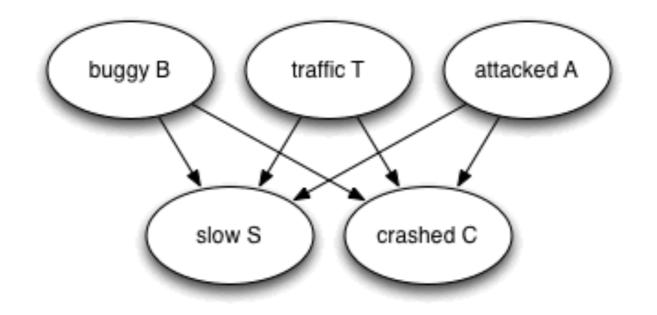


Everything in the top row is independent of each other.



Everything in the top row is independent of each other.

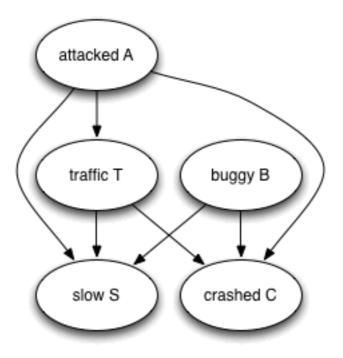
We write  $A \perp B$  for A is independent of B.



Everything in the top row is independent of each other.

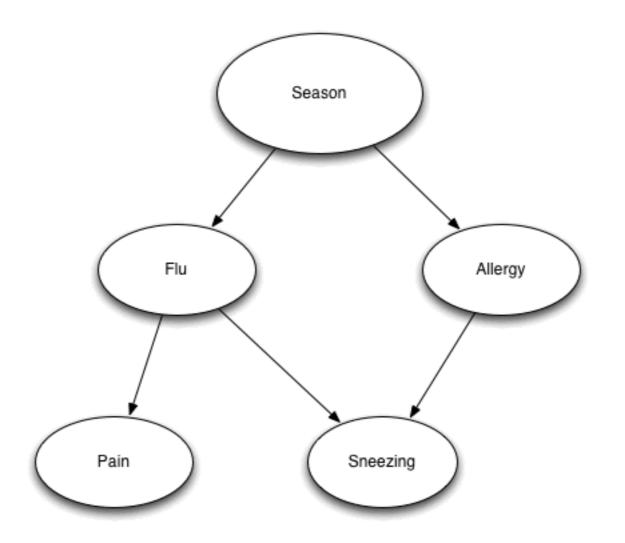
We write  $A \perp B$  for A is independent of B.

Perhaps too simple: if attacked then the traffic should be heavy.

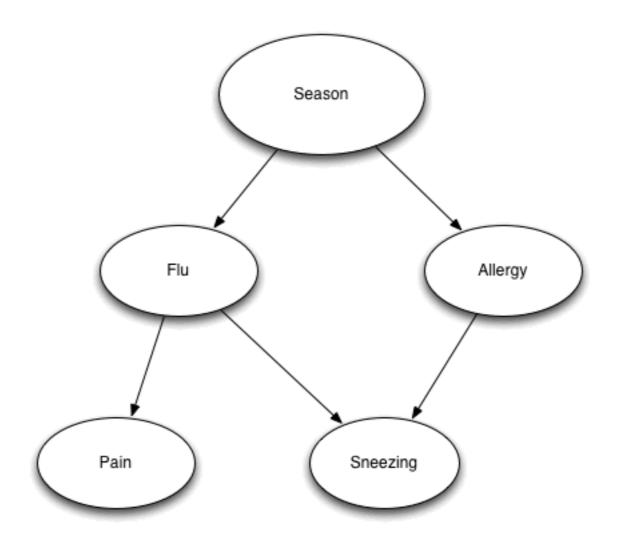


This version: traffic is affected by being attacked.

#### Medical example (from Koller and Friedman)

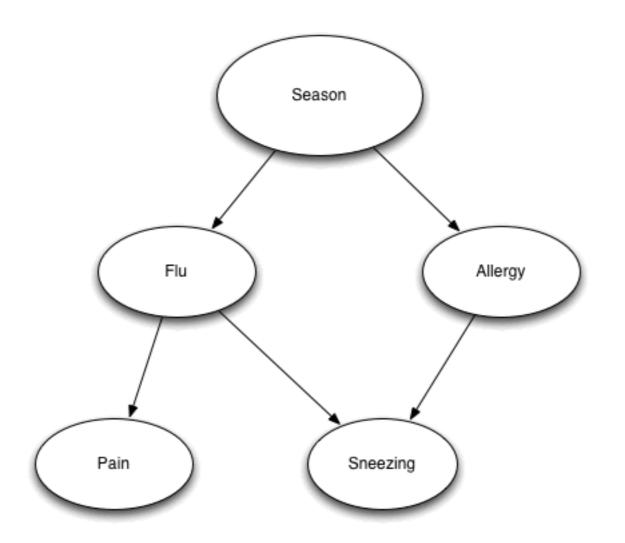


#### Medical example (from Koller and Friedman)



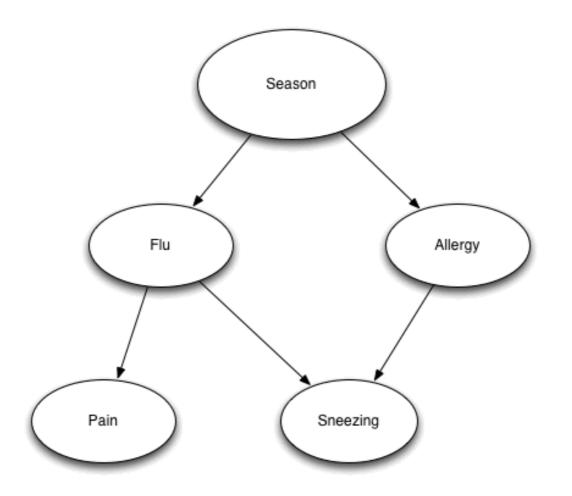
Flu and allergy are correlated through season.

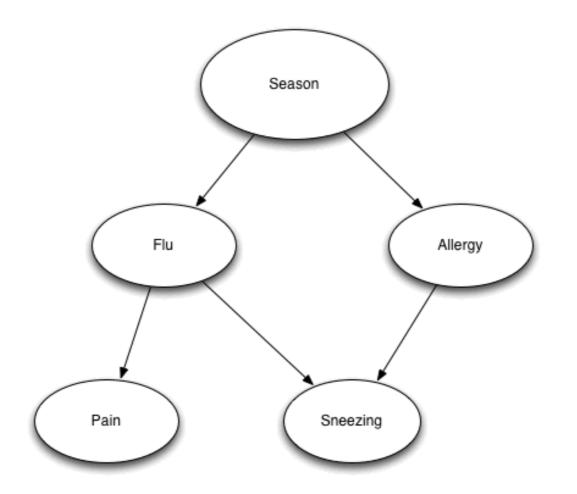
#### Medical example (from Koller and Friedman)

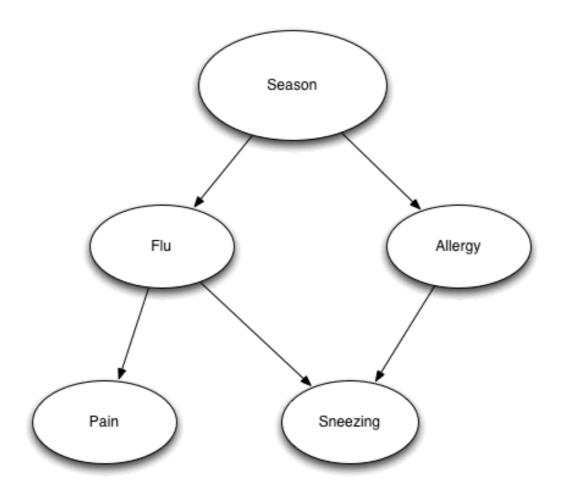


Flu and allergy are correlated through season.

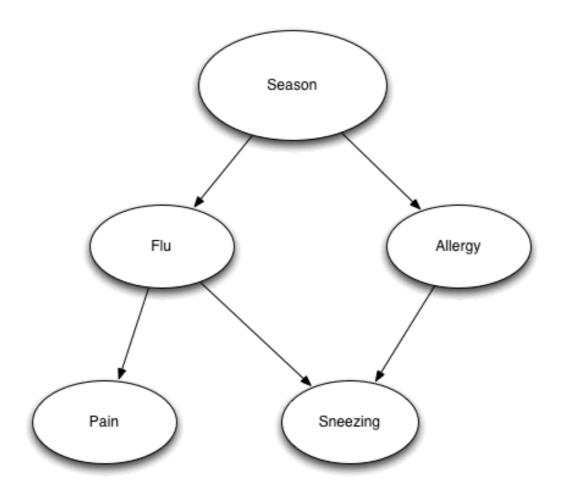
Given the season, they are independent:  $(A \perp F \mid S)$ 



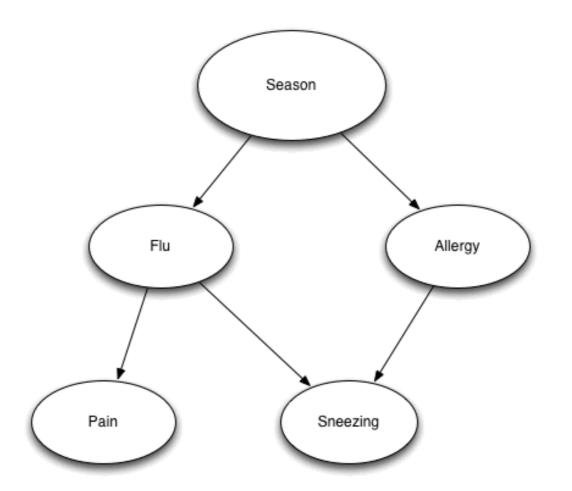




 $(F \perp A \mid S), (Sn \perp S \mid F, A), (P \perp A, Sn \mid F), (P \mid Sn \mid F)$ 



 $(F \perp A \mid S), (Sn \perp S \mid F, A), (P \perp A, Sn \mid F), (P \mid Sn \mid F)$  $P(Sn \mid F, A, S) = P(Sn \mid F, A)$ 



 $(F \perp A \mid S), (Sn \perp S \mid F, A), (P \perp A, Sn \mid F), (P \mid Sn \mid F)$ 

$$P(Sn \mid F, A, S) = P(Sn \mid F, A)$$

Sn depends on S but it is *conditionally* independent given A and F.

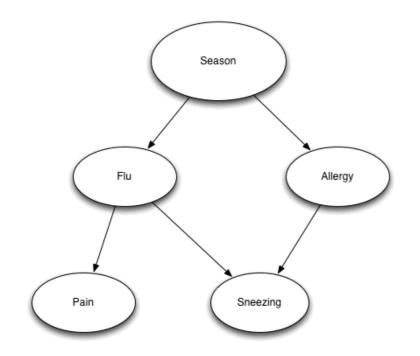
Why do we care about conditional independence?

Why do we care about conditional independence?

Because we can *factorize* the joint distributions.

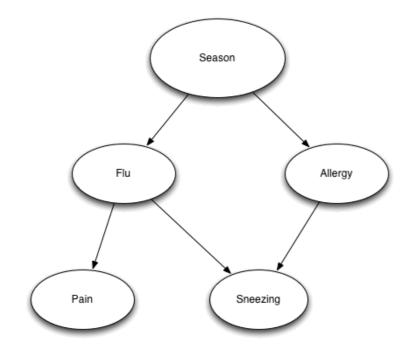
Why do we care about conditional independence?

Because we can *factorize* the joint distributions.



Why do we care about conditional independence?

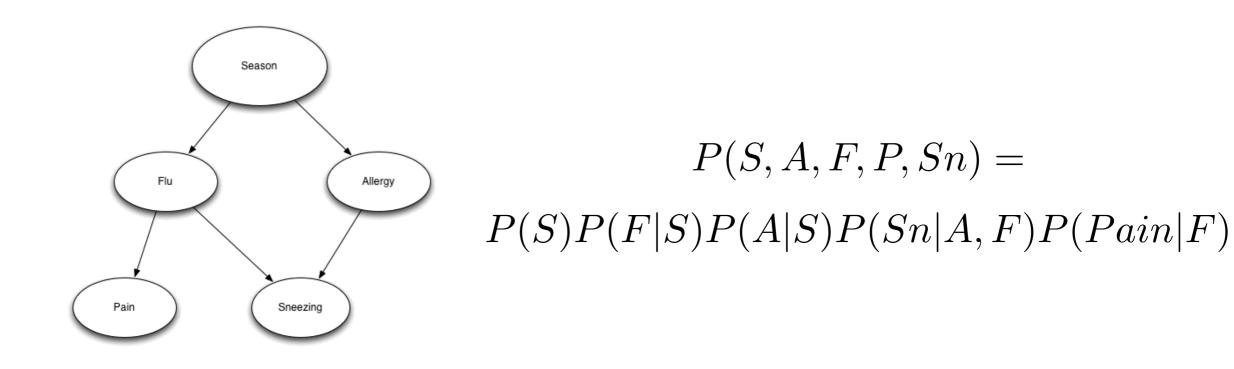
Because we can *factorize* the joint distributions.



P(S, A, F, P, Sn) = P(S)P(F|S)P(A|S)P(Sn|A, F)P(Pain|F)

Why do we care about conditional independence?

Because we can *factorize* the joint distributions.



A huge advantage for representing, computing and reasoning.

Continuous state spaces: robotics, telecommunication, control systems, sensor systems.

Continuous state spaces: robotics, telecommunication, control systems, sensor systems.

Requires measure theory and analysis.

Continuous state spaces: robotics, telecommunication, control systems, sensor systems.

Requires measure theory and analysis. Why?

Continuous state spaces: robotics, telecommunication, control systems, sensor systems.

Requires measure theory and analysis. Why?

It is subtle to answer questions like:

Continuous state spaces: robotics, telecommunication, control systems, sensor systems.

Requires measure theory and analysis. Why?

It is subtle to answer questions like:

I choose a real number between 0 and 1 uniformly, what is the probability of getting a rational number?

Continuous state spaces: robotics, telecommunication, control systems, sensor systems.

Requires measure theory and analysis. Why?

It is subtle to answer questions like:

I choose a real number between 0 and 1 uniformly, what is the probability of getting a rational number? Answer: 0!

Continuous state spaces: robotics, telecommunication, control systems, sensor systems.

Requires measure theory and analysis. Why?

It is subtle to answer questions like:

I choose a real number between 0 and 1 uniformly, what is the probability of getting a rational number? Answer: 0!

If every single point has probability 0 how can I have anything interesting?

Continuous state spaces: robotics, telecommunication, control systems, sensor systems.

Requires measure theory and analysis. Why?

#### It is subtle to answer questions like:

I choose a real number between 0 and 1 uniformly, what is the probability of getting a rational number? Answer: 0!

If every single point has probability 0 how can I have anything interesting?

Is it possible to assign a probability to every set?

Continuous state spaces: robotics, telecommunication, control systems, sensor systems.

Requires measure theory and analysis. Why?

#### It is subtle to answer questions like:

I choose a real number between 0 and 1 uniformly, what is the probability of getting a rational number? Answer: 0!

If every single point has probability 0 how can I have anything interesting?

Is it possible to assign a probability to every set? Answer: No!

Continuous state spaces: robotics, telecommunication, control systems, sensor systems.

Requires measure theory and analysis. Why?

#### It is subtle to answer questions like:

I choose a real number between 0 and 1 uniformly, what is the probability of getting a rational number? Answer: 0!

If every single point has probability 0 how can I have anything interesting?

Is it possible to assign a probability to every set? Answer: No!

There is a rich and fascinating theory of programming and reasoning about probabilistic systems.

#### Logic and Probability are your weapons. Go forth and conquer the software world!

Thank you!