

Dexter Kozen's Influence
on the Theory of
Markov Processes

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Ithaca 1985

- ◆ Department of Computer Science hired two new profs that year:
- ◆ Dexter Kozen and me.
- ◆ Dexter was famous for the so-called BKR algorithm but

Semantics and Logic

- ◆ also for modal logic (PDL)
- ◆ mu-calculus and
- ◆ probabilistic programs.

Two key papers

- ◆ Semantics of probabilistic programs, JCSS, 1981
- ◆ A probabilistic PDL, JCSS, 1985.

New ideas

- ◆ Take probability theory on **continuous state spaces** seriously.
- ◆ Hence, work with sigma-algebras, measure theory and integration.
- ◆ A new Stone-type duality.

What is Stone duality?

- ◆ Stone representation theorem: every boolean algebra is isomorphic to a concrete boolean algebra of sets.
- ◆ But we know “elements are a hack!”

- ◆ If we look at all the boolean algebras and the maps between them then
- ◆ it looks exactly like a certain collection of topological spaces and the maps between them
- ◆ going backwards!

- ◆ This kind of thing happens several times in mathematics:
- ◆ C^* algebras and compact Hausdorff spaces (Gelfand duality)
- ◆ Finite-dimensional vector spaces and themselves (self duality)

- ◆ Happens in computer science too:
- ◆ We can define what a program does by going forwards: $St \times Act \dashrightarrow St$
- ◆ or backwards:
- ◆ precondition \dashleftarrow predicate $\times Act$.

- ◆ Dexter found that this works for probabilistic programs too:
- ◆ Forwards semantics: measure transformers (Markov kernels)
- ◆ Backward semantics: transformers of random variables (value functions).

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What does this have to do with relations?

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The category of Markov processes (1999), ENTCS, P.

Dexter's Probabilistic Language

$S ::= x_i := f(\vec{x}) \mid S_1; S_2 \mid \text{if } \mathbf{B} \text{ then } S_1 \text{ else } S_2 \mid \text{while } \mathbf{B} \text{ do } S.$

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He also gave *backward semantics* in terms of how random variables transform.

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In Dexter's probabilistic setting:

If f is the value of a random variable after a command described by a Markov kernel h , then $\int f(x')h(x, dx')$ is the value before.

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Duality: $\mathbf{SRel} \equiv \mathbf{SPT}^{op}$.

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In the late 1990s I, working with Josée Desharnais, Abbas Edalat, Radha Jagadeesan and Vineet Gupta were developing the theory of Labelled Markov Processes (MDPs).

A key concept was *bisimulation*: when do two processes behave *exactly* the same?

Logical characterization: when they agree on all the formulas of a simple (modal) logic.

But equivalences are suspect, one should work with metrics.

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How does one make this precise?

The Kozen Analogy

Logic	Probability
State s	Distribution μ
Predicate ϕ	Random variable f
Satisfaction $s \models \phi$	Pairing $\int f d\mu$

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The study of behavioural metrics is now flourishing with many papers every year in concurrency theory and also in machine learning and algorithms.

But it all started with a
random walk
through
Markov processes.

Happy Birthday Dexter!