

COMP250: Induction proofs.

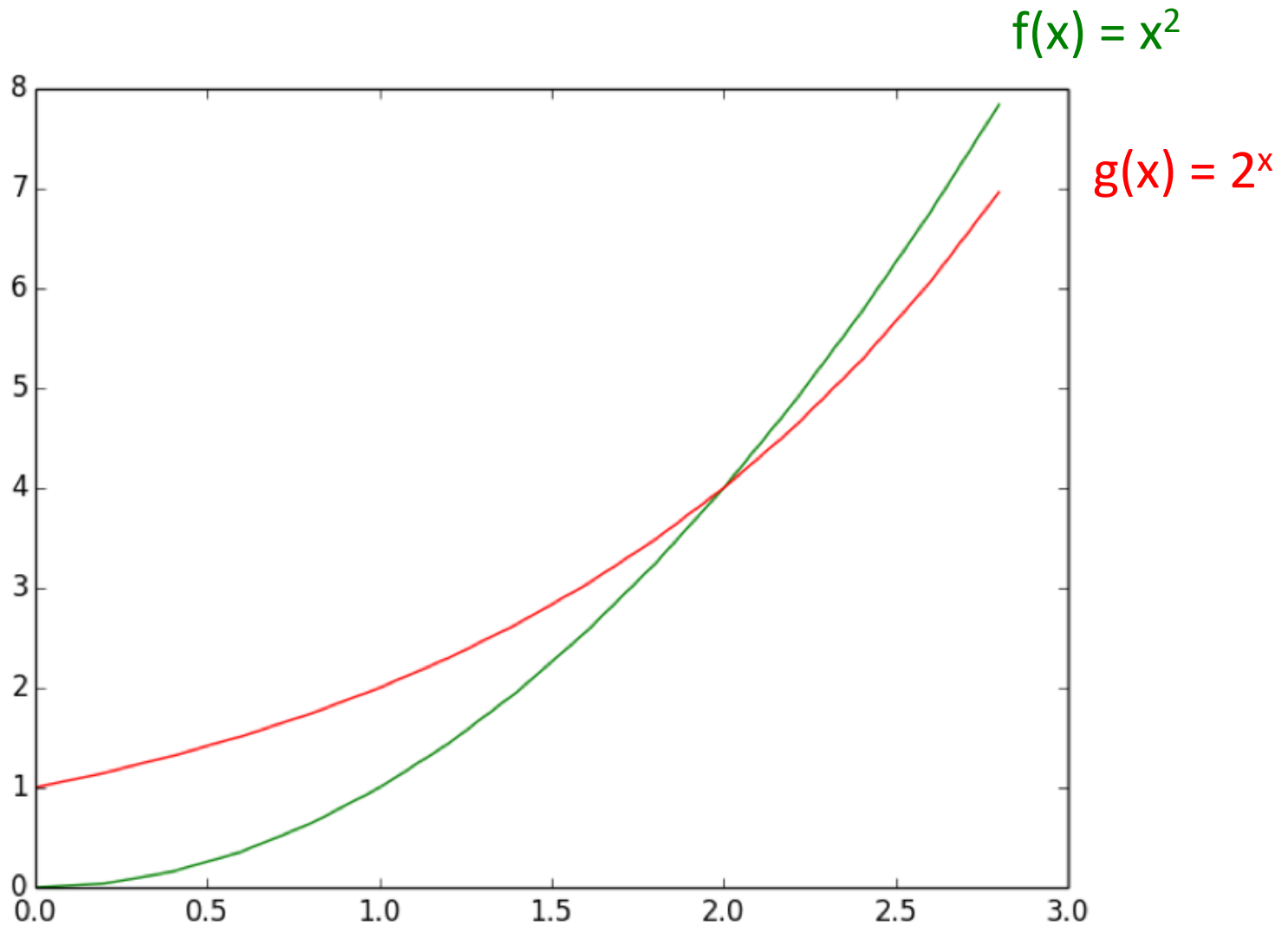
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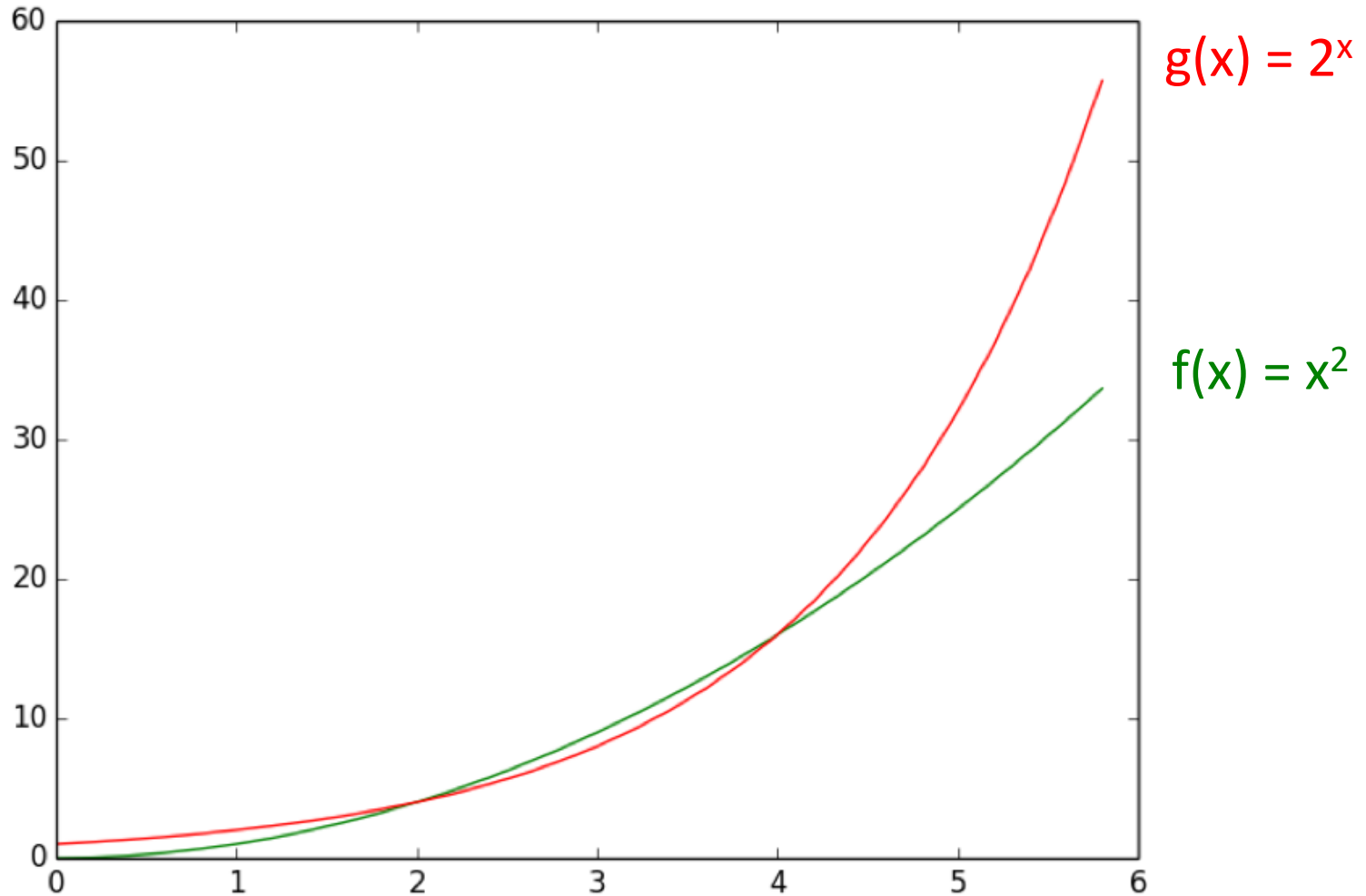
McGill University

Based on slides from (Langer,2012)

for any $n \geq 2$, $n^2 \geq 2^n$?



for any $n \geq 5$, $n^2 \leq 2^n$?



Motivation

How to prove these?

$$\text{for any } n \geq 1, \quad 1 + 2 + 3 + 4 + \cdots + n = \frac{n \cdot (n + 1)}{2}$$

$$\text{for any } n \geq 1, \quad 1 + 3 + 5 + 7 + \cdots + (2 \cdot n - 1) = n^2$$

$$\text{for any } n \geq 5, \quad n^2 \leq 2^n$$

And in general, any statement of the form:

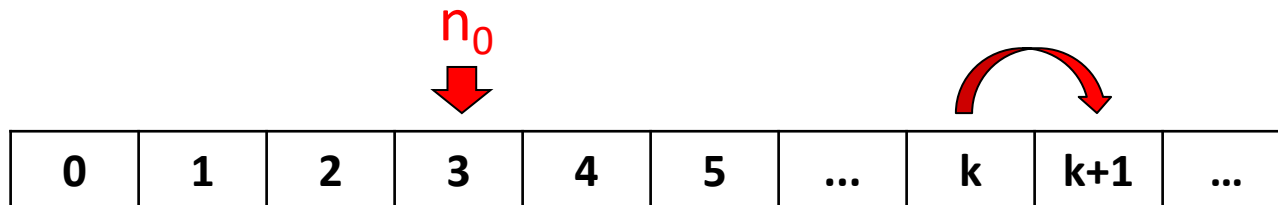
“for all $n \geq n_0$, $P(n)$ ” where $P(n)$ is some proposition.

Mathematical induction

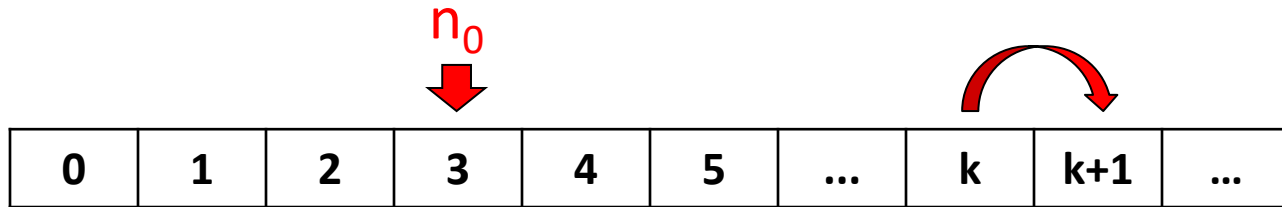
Many statement of the form “*for all $n \geq n_0$, $P(n)$* ” can be proven with a logical argument call *mathematical induction*.

The proof has two components:

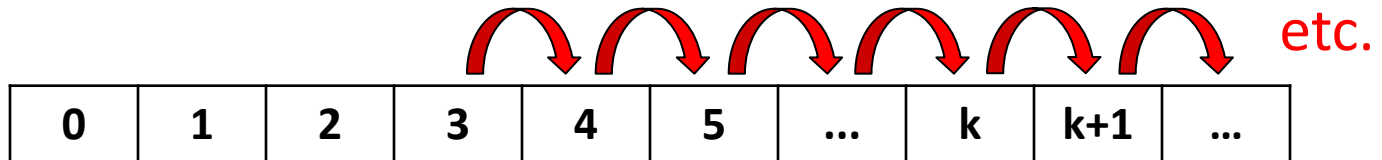
- **Base case:** $P(n_0)$
- **Induction step:** for any $n \geq n_0$, if $P(n)$ then $P(n+1)$



Principle




Implies



Example 1

Claim: *for any* $n \geq 1$, $1 + 2 + 3 + 4 + \dots + n = \frac{n \cdot (n + 1)}{2}$

Proof:

- Base case: $n = 1$ $1 = \frac{1 \cdot 2}{2}$ 

- Induction step:

for any $k \geq 1$, *if* $1 + 2 + 3 + 4 + \dots + k = \frac{k \cdot (k + 1)}{2}$

then $1 + 2 + 3 + 4 + \dots + k + (k + 1) = \frac{(k + 1) \cdot (k + 2)}{2}$

Example 1


Assume $1 + 2 + 3 + 4 + \dots + k = \frac{k \cdot (k + 1)}{2}$

then $1 + 2 + 3 + 4 + \dots + k + (k + 1)$

$$= \frac{k \cdot (k + 1)}{2} + (k + 1)$$

$$= \frac{k \cdot (k + 1) + 2 \cdot (k + 1)}{2}$$

$$= \frac{(k + 2) \cdot (k + 1)}{2}$$



Induction hypothesis

Summary

Base case: $P(1)$ ✓

Induction step: for any $k \geq 1$, if $P(k)$ then $P(k+1)$ ✓

Thus for all $n \geq 1$, $P(n)$ □

Example 2

Claim: *for any* $n \geq 1$, $1 + 3 + 5 + 7 + \cdots + (2 \cdot n - 1) = n^2$

Proof:

• Base case: $n = 1$ $1 = 1^2$ 


• Induction step:

for any $k \geq 1$, *if* $1 + 3 + 5 + 7 + \cdots + (2 \cdot k - 1) = k^2$

then $1 + 3 + 5 + 7 + \cdots + (2 \cdot (k + 1) - 1) = (k + 1)^2$

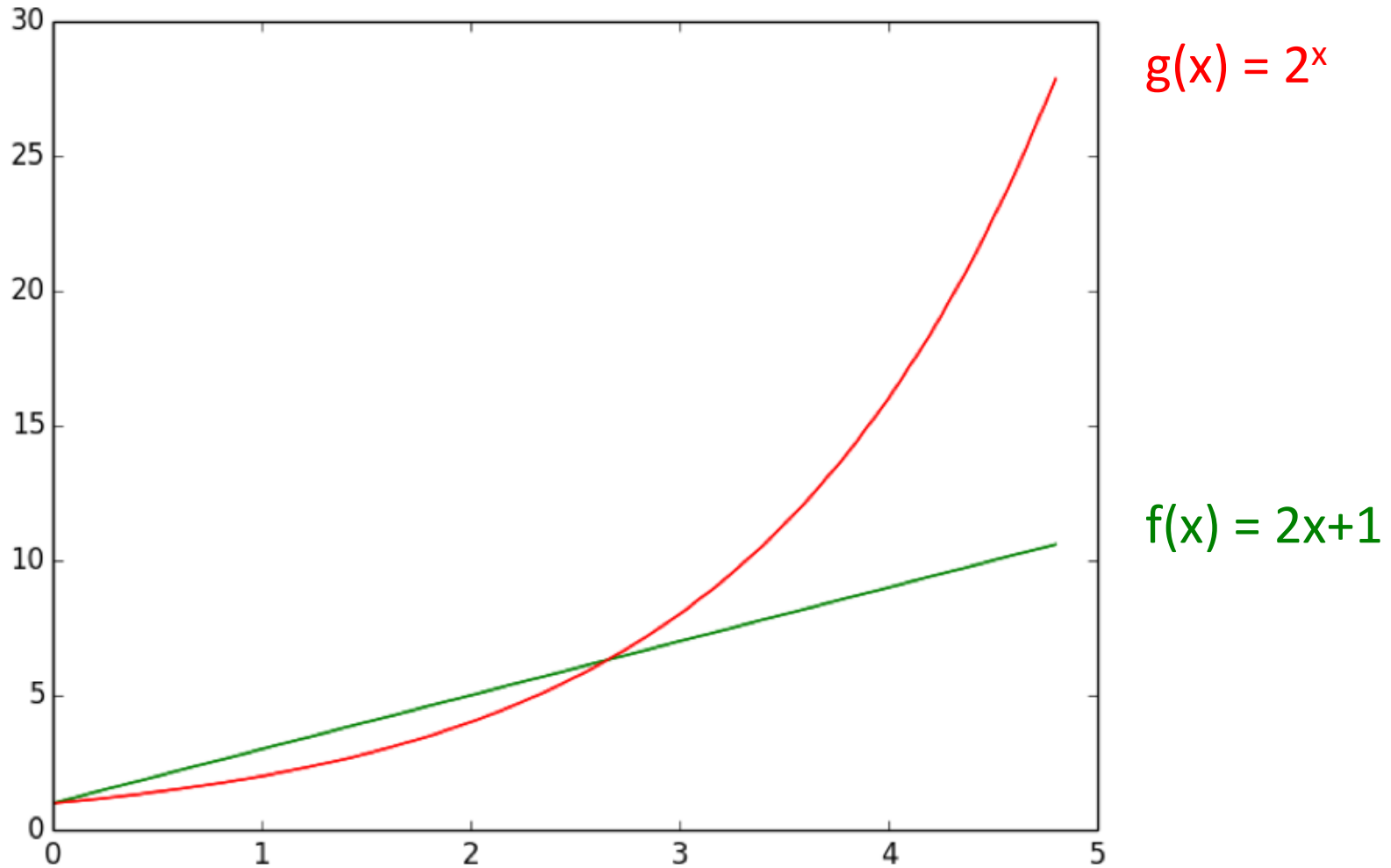
Example 2

Assume $1 + 3 + 5 + 7 + \cdots + (2 \cdot k - 1) = k^2$

then $1 + 3 + 5 + 7 + \cdots + (2 \cdot k - 1) + (2 \cdot (k + 1) - 1)$  Induction hypothesis

$$= k^2 + 2 \cdot (k + 1) - 1$$
$$= k^2 + 2 \cdot k + 1$$
$$= (k + 1)^2$$

Example 3



Example 3

Claim: $\text{for any } n \geq 3, \quad 2 \cdot n + 1 < 2^n$

Proof:

- Base case: $n = 3 \quad 2 \cdot 3 + 1 = 7 < 2^3 = 8$



- Induction step:

$\text{for any } k \geq 3, \quad \text{if } 2 \cdot k + 1 < 2^k$

$\text{then } 2 \cdot (k + 1) + 1 < 2^{k+1}$

Example 3

Assume $2 \cdot k + 1 < 2^k$


then $2 \cdot (k + 1) + 1$


$$= 2 \cdot k + 2 + 1$$

$$< 2^k + 2$$

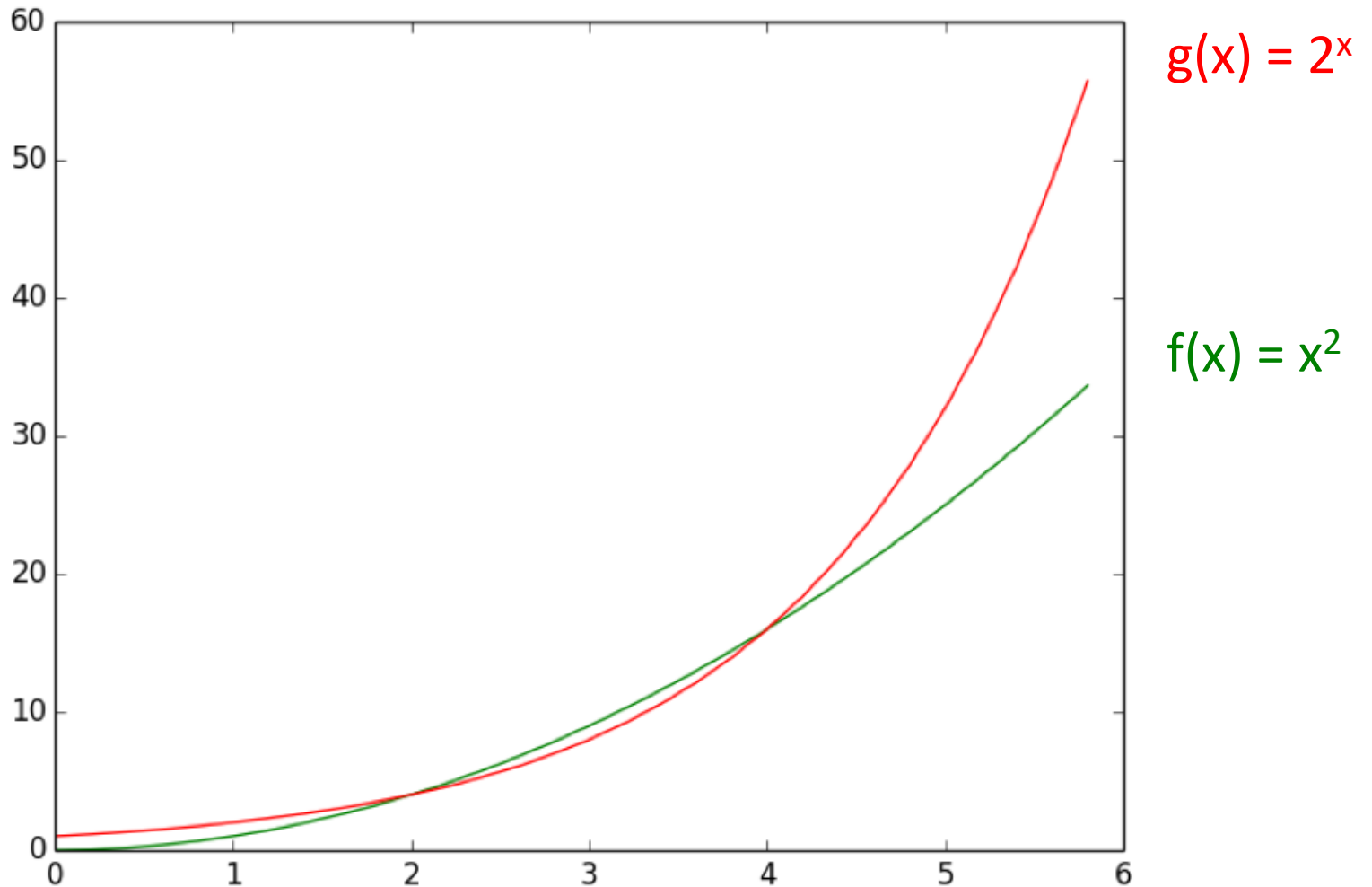
$$\leq 2^k + 2^k \quad \text{for } k \geq 1$$

$$= 2^{k+1}$$

 Induction hypothesis

 Stronger than we need,
but that works!


Example 4



Example 4

Claim: *for any $n \geq 5$, $n^2 \leq 2^n$*

Proof:

• Base case: $n = 5$ $25 \leq 32$ 

• Induction step:

*for any $k \geq 5$, if $k^2 \leq 2^k$
then $(k+1)^2 \leq 2^{k+1}$*

Example 4

Assume $k^2 \leq 2^k$

then $(k+1)^2$

$$= k^2 + 2 \cdot k + 1$$

$$\leq 2^k + 2 \cdot k + 1$$

$$\leq 2^k + 2^k$$

$$= 2^{k+1}$$



Induction hypothesis

From previous example

Example 5

Fibonacci sequence:

$$\text{Fib}_0 = 0 \quad \text{base case}$$

$$\text{Fib}_1 = 1 \quad \text{base case}$$

$$\text{Fib}_n = \text{Fib}_{n-1} + \text{Fib}_{n-2} \quad \text{for } n > 1 \quad \text{recursive case}$$

Claim: For all $n \geq 0$, $\text{Fib}_n < 2^n$

Base case: $\text{Fib}_0 = 0 < 2^0 = 1$, $\text{Fib}_1 = 1 < 2^1 = 2$

Q: Why should we check both Fib_0 and Fib_1 ?

Induction step: for any $i \leq k$, if $\text{Fib}_i < 2^i$ then $\text{Fib}_{k+1} < 2^{k+1}$

Example 5

Assume that *for all $i \leq k$, $Fib_i < 2^i$* (Note variation of induction hypothesis)

$$\begin{aligned} \text{Then } Fib_{k+1} &= Fib_k + Fib_{k-1} \\ &< 2^k + 2^{k-1} \\ &\leq 2^k + 2^k \\ &= 2^{k+1} \end{aligned}$$



Induction hypothesis (x2)

for $k \geq 1$