Compositional Semantics: Montagovian Semantics and Lambda Calculus

> COMP-599 Oct 26, 2016

Outline

Principle of compositionality Semantic inference First-order logic Lambda calculus

The Principle of Compositionality

Compositionality: The meaning of a phrase depends on the meanings of its parts.

COMP-599 is a fantastically awesome class.

Lexical semantics gives us the meanings and behaviours of each of the words:

• COMP-599, is, a, fantastically, awesome, class

We build up the meaning of the entire sentence through composition.

Idioms - Violation of Compositionality

Idioms are expressions whose meanings cannot be predicted from their parts.

kick the bucket the last straw piece of cake hit the sack

Co-Compositionality

Consider the meaning of *red* when modifying each of the following nouns

- rose
- wine
- cheeks
- hair

Is red really combining compositionally with each of these nouns?

 Co-compositionality (Pustejovsky, 1995) – the meanings of words depend also on the other words that they are composed with

Goal of Compositional Semantics

Derive a meaning representation of a phrase/sentence from its parts

What is a good meaning representation?

Relates the linguistic expression to the world:

 Asserts a proposition that is either true or false relative to the world

Pandas are purple and yellow.

- Conveys information about the world It will snow tomorrow.
- Is a query about the state of the world What is the weather like in Montreal next week?

Semantic Inference

Making *explicit* something that is *implicit* in language (Blackburn and Bos, 2003)

I want to visit the capital of Italy.

The capital of Italy is Rome.

∴ I want to visit Rome.

All wugs are blorks.

All blorks are cute.

: All wugs are cute.

Montagovian Semantics

Montague (1970) started a tradition of using a logical formalism to represent the meaning of a sentence, with a tight connection to syntax.

There is in my opinion no important theoretical difference between natural languages and the artificial languages of logicians; indeed I consider it possible to comprehend the syntax and semantics of both kinds of languages with a single natural and mathematically precise theory. (Montague 1970c, 222)

Natural language inference then can be seen as applying logical rules of inference.

First-Order Predicate Calculus

Domain of discourse

A set of entities that we care about

e.g., the students in the class, the topics we study, classrooms, courses, etc.

Variables

Typically lower-case Stands for potential elements of the domain e.g., *x*, *y*, *z*

First-Order Predicate Calculus

Predicates

- Maps elements of the domain to truth values
- Can be of different valences
- e.g. inCourse(x, y): takes in two elements of the domain, returns true if x is a student in course y, false otherwise

Functions

- Maps elements to other elements
- Can be of different valences
- e.g. instructorOf(x): takes x, returns an element corresponding to x's instructor
- What is a valence 0 function?

First-Order Predicate Calculus

Logical connectives

• All the standard ones $\neg, \Lambda, V, \rightarrow, \leftrightarrow$

Quantifiers

- Existential 3
- Universal ∀

Example Sentences

The capital of Italy is Rome. capitalOf(ITALY) = ROME

All wugs are blorks.

 $\forall x. wug(x) \rightarrow blork(x)$

Interpretating FOL

A particular instance of a FOL consists of:

- Predicate and function names and arity
- A set of sentences in FOL using those predicates and functions

An interpretation or model of a FOL consists of:

Domain of discourse, D

Mapping for the functions to elements of D

Mapping for the predicates to True or False

Exercise

Come up with a FOL characterization of the following:

- Students who study AND do homework will get an A
- Students who only do one of them get a B
- Students who do neither get a C

List the predicates and functions that are necessary. Make constants for the grades (A, B, C).

Come up with an interpretation of this FOL, where you and your neighbours are the elements in the domain of discourse, such that the above FOL formulas are true.

Building Meaning Representations

Target MR: a logical formula in FOL

Still needed:

A procedure to map sentences to a FOL formula compositionally

Tool: Lambda calculus

Lambda Calculus

Basically a way to describe computation using mathematical functions

• The computation we will be doing is to build up a FOL sentence as the meaning representation of a sentence.

Terms in Lambda calculus can be defined recursively:

- A variable (e.g., x)
- $\lambda x. t$, where t is a lambda term
- *ts*, where *t* and *s* are lambda terms

Functional Abstraction and Application

Function application (or **beta reduction**) of term $(\lambda x. t)s$

• Replace all instances of x in t with the expression s e.g., $(\lambda x. x + y)$ 2 simplifies to 2 + y $(\lambda x. xx)(\lambda x. x) = (\lambda x. x)(\lambda x. x) = (\lambda x. x)$

Function application is **left-associative**: abcd = ((ab)c)d

I define this notion intuitively here, and gloss over some details, but these definitions can (should) be formalized, in order to be precise:

http://arxiv.org/pdf/1503.09060.pdf

Power of Lambda Calculus

They allow us to store partial computations of the MR, as we are composing the meaning of the sentence constituent by constituent.

- Whiskers disdained catnip.
- disdained
- disdained catnip

 $\lambda x. \lambda y. disdained(y, x)$ ($\lambda x. \lambda y. disdained(y, x)$) catnip = $\lambda y. disdained(y, catnip)$

Whiskers disdained catnip (λy.disdained(y,catnip))Whiskers = disdained(Whiskers,catnip)

Exercises

What is the result of simplifying the following expressions in lambda calculus through beta reduction? $(\lambda z. z)(\lambda y. y y)(\lambda x. x a)$

 $(((\lambda x.\lambda y.(x y))(\lambda y.y)) w)$

 $(\lambda x. x x) (\lambda y. y x) z$

Syntax-Driven Semantic Composition

Augment CFG trees with lambda expressions

• Syntactic composition = function application

Semantic attachments:

 $A \rightarrow \alpha_1 \dots \alpha_n$

syntactic composition

$$f(\alpha_j.sem,...,\alpha_k.sem)$$

semantic attachment

Proper Nouns

Proper nouns are FOL constants $PN \rightarrow COMP599$ {COMP599}

Actually, we will **type-raise** proper nouns $PN \rightarrow COMP599 \qquad \{\lambda x. x(COMP599)\}$

- It is now a function rather than an argument.
- We will see why we do this next class.

NP rule: $NP \rightarrow PN$ {PN.sem}

Common Nouns

Common nouns are predicates inside a lambda expression of type $\langle e, t \rangle$

 Takes an entity, tells you whether the entity is a member of that class

 $N \rightarrow student \qquad \{\lambda x. Student(x)\}$

Let's talk more about common nouns next class when we also talk about quantifiers.

Intransitive Verbs

We introduce an *event variable e*, and assert that there exists a certain event associated with this verb, with arguments.

 $V \rightarrow rules \qquad \{\lambda x. \exists e. Rules(e) \land Ruler(e, x)\}$

Then, composition is

 $S \rightarrow NP VP$ {NP.sem(VP.sem)}

Let's derive the representation of the sentence "COMP-599 rules"

Neo-Davidsonian Event Semantics Notice that we have changed how we represent events Method 1: multi-place predicate Rules(x)

Method 2: Neo-Davidsonian version with event variable

 $\exists e. Rules(e) \land Ruler(e, x)$

Reifying the event variable makes things more flexible

- Optional elements such as location and time, passives
- Add information to the event variable about tense, modality

Transitive Verbs

Transitive verbs

 $V \rightarrow enjoys \\ \{\lambda w. \lambda z. w(\lambda x. \exists e. Enjoys(e) \land Enjoyer(e, z) \land Enjoyee(e, x))\}$

 $VP \rightarrow V NP$ {V.sem(NP.sem)} $S \rightarrow NP VP$ {NP.sem(VP.sem)}

Exercise: verify that this works with the sentence "Jackie enjoys COMP-599"

Next Class

Quantifiers and common nouns Adjectives, adverbs, and modifiers Underspecification