1. (5 Points) Prove that every Turing Recognizable language $L$ satisfies $L \leq_m A_{TM}$.

2. (10 Points) Prove that for any two languages $A$ and $B$, there exists a language $J$ such that $A \leq_T J$ and $B \leq_T J$.

3. (10 Points) Prove that for any languages $A$, there exists a language $J$ such that $A \leq_T J$ and $J \not\leq_T A$.

4. (15 Points) Let $\Gamma = \{0, 1, \sqcup\}$ be the tape alphabet for all Turing Machines in this problem. Define the busy beaver function $BB : \mathbb{N} \to \mathbb{N}$ as follows. For each value of $k$, consider all $k$-state Turing Machines that halt when started with a blank tape. Let $BB(k)$ be the maximum number of 1’s that remain on the tape among all of those machines. Prove that $BB$ is not a computable function.

5. (15 points) Let $f : \mathbb{N} \to \mathbb{N}$ be defined as

$$f(x) = \begin{cases} 3x + 1 & x \text{ is odd} \\ x/2 & x \text{ is even} \end{cases}$$

For every natural number $x$, if you start with $x$ and iterate $f$, you obtain a sequence

$$x, f(x), f(f(x)), f(f(f(x))), \ldots.$$  

Stop if you ever hit 1. For example, if $x = 17$, we get the sequence

$$17, 52, 26, 13, 40, 20, 10, 5, 16, 8, 4, 2, 1.$$  

It has been checked by computers that if we start with any number $x < 10^{17}$, we will eventually hit 1 and terminate, but it is unknown whether this is true for every integer $x$. This is known
as Collatz’s conjecture\(^1\). Suppose that \(\text{HALT}_{TM}\) were decidable, and \(R\) was a TM that was deciding \(\text{HALT}_{TM}\). Construct an Turing Machine \(M\) based on \(R\) such that it accepts \(\varepsilon\) if Collatz’s conjecture is true (all numbers \(x\) hit 1 eventually), and rejects it if Collatz’s conjecture is false (there is some \(x\) that does not hit 1).

6. (15 points) Let

\[
X = \{\langle M, w \rangle | M \text{ is a single-tape TM that never modifies the } w \text{ part of the tape, and it accepts } w \}.
\]

Is \(X\) decidable? Prove your answer.

7. (15 Points) Prove that a language \(L\) is Turing Recognizable if and only if it can be expressed as

\[
L = \{x | \exists y \text{ such that } \langle x, y \rangle \in R\}
\]

where \(R\) is a decidable language. You need to prove that every language of this form is Turing recognizable, and that every Turing recognizable language can be described as above for some decidable language \(R\).

8. (15 Points) Determine whether the following language is decidable or undecidable:

\[
X = \{\langle M \rangle | \text{On every input } w, \text{ } M \text{ eventually leaves the start state}\}.
\]

\(^1\)It is interesting to view this in regard to Q4 in the previous assignment. It is fairly easy to construct a Turing Machine \(M\) with 100 states that given \(1^x\) as input, applies \(f\) recursively to this, and halts and accepts if we hit 1. Note that \(M\) is a decider if and only if Collatz’s conjecture is true, which currently is unknown.