

Compiler Design

Lecture 18:

Instruction Selection via Peephole Matching

Christophe Dubach

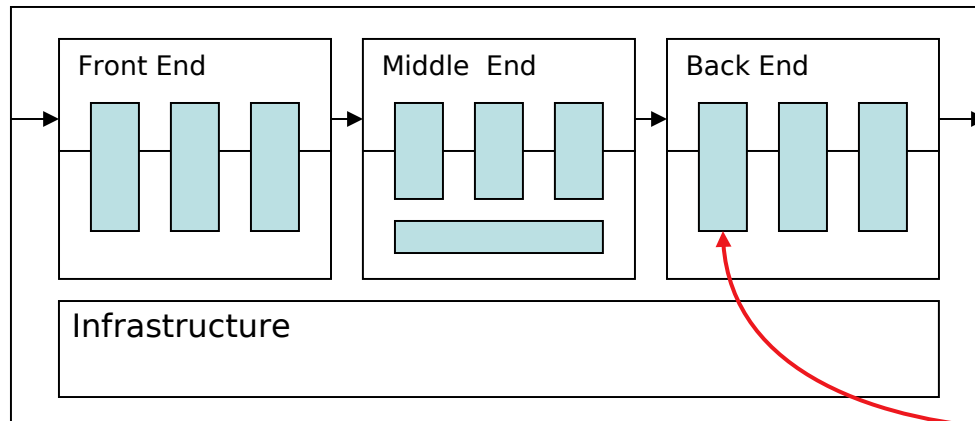
Winter 2025

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The Problem

Writing a compiler is a lot of work

- Would like to reuse components whenever possible
- Would like to **automate** construction of components



Today's lecture:
Automating
Instruction
Selection

- Front end construction is largely automated
- Middle is largely hand crafted
- (Parts of) back end can be automated

Definitions

Instruction selection

- Process of mapping **IR** into assembly code
 - Assumes a fixed storage mapping & code shape
 - Combining operations, using address modes

Instruction scheduling

- Process of reordering operations to hide latencies
 - Assumes a fixed program (*set of operations*)
 - Changes demand for registers

Register allocation

- Process of deciding which values will reside in registers
 - Changes the storage mapping, may add false sharing
 - Concerns about placement of data & memory operations

The Problem

Modern computers have many ways to do anything

Consider register-to-register copy

- Obvious operation is `MOVE  $r_j \leftarrow r_i$`
- Many others exist

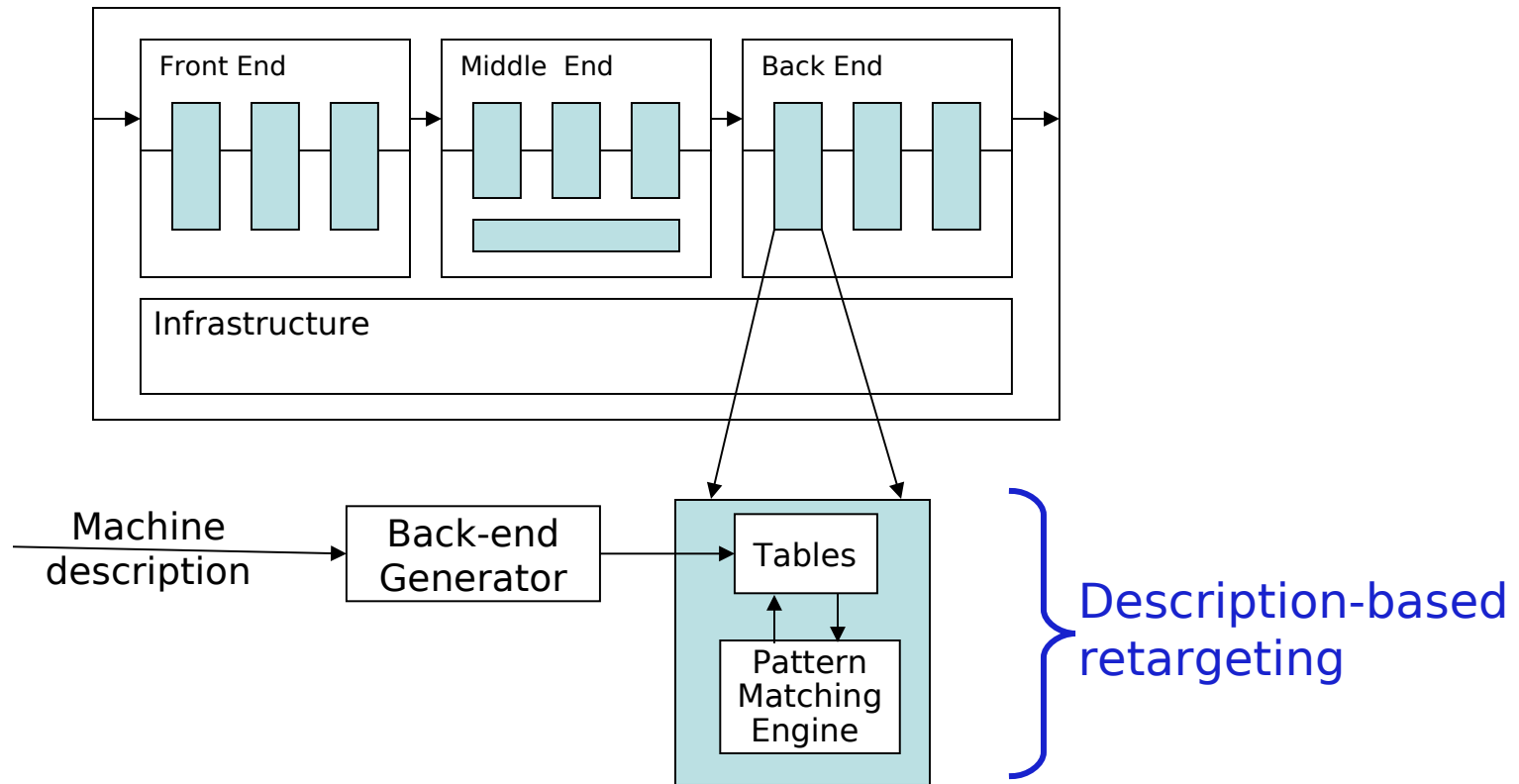
<code>ADDI $r_j \leftarrow r_i, 0$</code>	<code>OR $r_j \leftarrow r_i, r_0$</code>	<code>SLL $r_j \leftarrow r_i, 0$</code>
<code>ADD $r_j \leftarrow r_i, r_0$</code>	<code>MULI $r_j \leftarrow r_i, 1$</code>	<code>SRL $r_j \leftarrow r_i, 0$</code>
<code>ORI $r_j \leftarrow r_i, 0$</code>	<code>XORI $r_j \leftarrow r_i, 0$</code>	... and others ...

- Human would ignore all of these
- Algorithm must look at all of them & find low-cost encoding
 - Take context into account (*busy functional unit?*)

And this is an overly-simplified example

The Goal

Want to automate generation of instruction selectors



Machine description should also help with scheduling & allocation

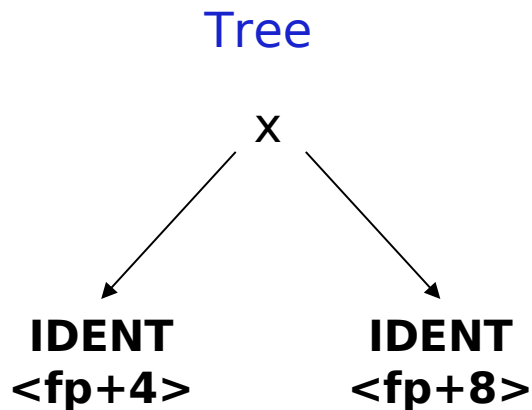
The Big Picture

Need pattern matching techniques

- Must produce good code *(some metric for good)*
- Must run quickly

Our treewalk (visitor) code generator ran quickly

How good was the code?



The Big Picture

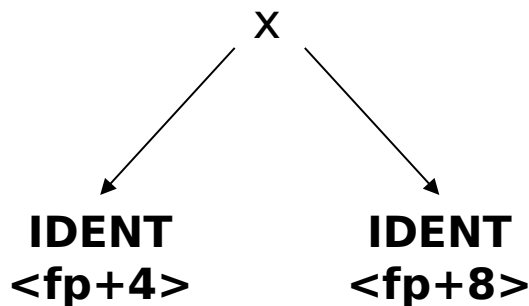
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How good was the code?

Tree



Treewalk Code

```
MOVE r1 ← 4
ADD  r2 ← r1 , $fp
LW   r3 ← 0(r2)
MOVE r4 ← 8
ADD  r5 ← r4 , $fp
LW   r6 ← 0(r5)
MUL  r7 ← r3, r6
```

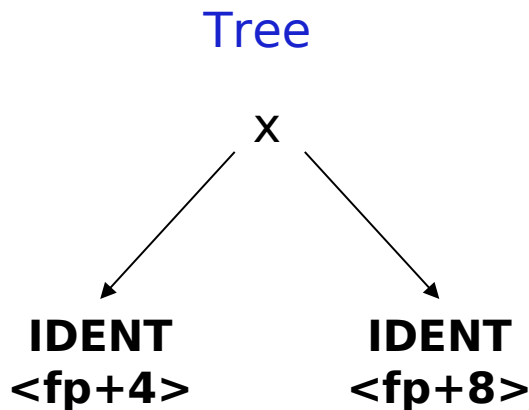
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Treewalk Code

```
MOVE r1 ← 4
ADD r2 ← r1 , $fp
LW r3 ← 0(r2)
MOVE r4 ← 8
ADD r5 ← r4 , $fp
LW r6 ← 0(r5)
MUL r7 ← r3 , r6
```

Desired Code

```
LW r3 ← 4($fp)
LW r6 ← 8($fp)
MUL r7 ← r3 , r6
```

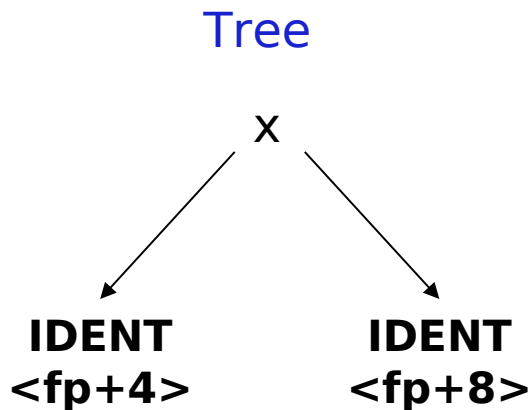

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```
MOVE r1 ← 4
ADD r2 ← r1 , $fp
LW r3 ← 0(r2)
MOVE r4 ← 8
ADD r5 ← r4 , $fp
LW r6 ← 0(r5)
MUL r7 ← r3 , r6
```

Desired Code

```
LW r3 ← 4($fp)
LW r6 ← 8($fp)
MUL r7 ← r3 , r6
```

Okay wasn't too hard, could do it in the code generator

The Big Picture

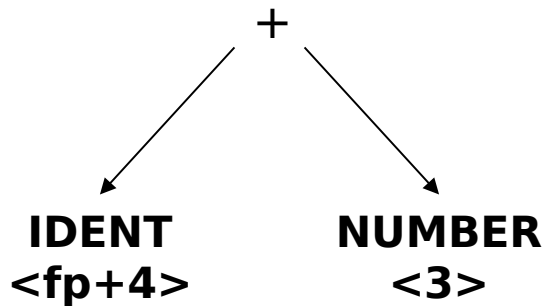
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The Big Picture

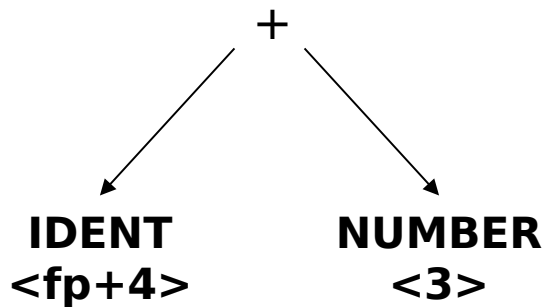
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How good was the code?

Tree



Treewalk Code

```
MOVE r1 ← 4
ADD  r2 ← r1 , $fp
LW   r3 ← 0(r2)
MOVE r4 ← 3
ADD  r5 ← r3 , r4
```

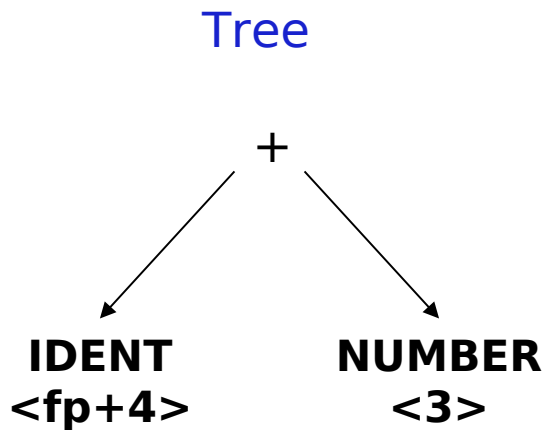
The Big Picture

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How good was the code?



Treewalk Code

```
MOVE r1 ← 4
ADD  r2 ← r1 , $fp
LW   r3 ← 0(r2)
MOVE r4 ← 3
ADD  r5 ← r3 , r4
```

Desired Code

```
LW    r3 ← 4($fp)
ADDI  r5 ← r3 , 3
```

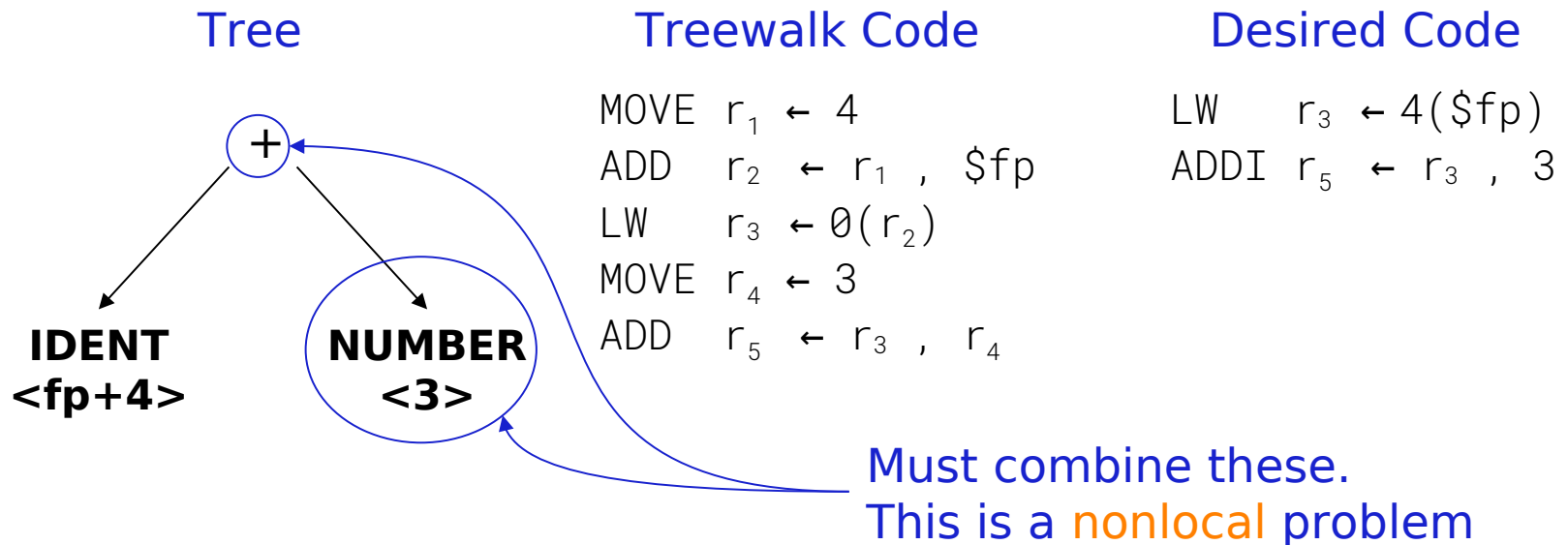
The Big Picture

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How do we perform this kind of matching ?

Tree-oriented IR suggests pattern matching on trees

- Tree-patterns as input, matcher as output
- Each pattern maps to a target-machine instruction sequence
- Use dynamic programming or bottom-up rewrite systems

Linear IR suggests using some sort of string matching

- Strings as input, matcher as output
- Each string maps to a target-machine instruction sequence
- Use text matching (Aho-Corasick) or peephole matching

In practice, both work well. Today we will look at matchers on Linear IR.

Peephole Matching

- Basic idea:
Compiler can discover local-*ish* improvements
 - Look at a small set of adjacent operations
 - Move a “peephole” over code & search for improvement



Marco Verch , CC BY 2.0 Source: <https://foto.wuestenigel.com/a-mans-hand-opens-the-peephole-on-the-door/>

Peephole Matching

- Basic idea: Compiler can discover local improvements
 - Look at a small set of adjacent operations
 - Move a “peephole” over code & search for improvement
- Classic examples:
 - store followed by load

Original code

```
SW    r1 → 0(r2)  
LW    r15 ← 0(r2)
```

Improved code

```
SW    r1 → 0(r2)  
MOVE  r15 ← r1
```


Peephole Matching

- Basic idea: Compiler can discover local improvements
 - Look at a small set of adjacent operations
 - Move a “peephole” over code & search for improvement
- Classic examples:
 - store followed by load
 - simple algebraic identities

Original code

```
ADDI  r7 ← r2, 0  
MUL   r10 ← r4, r7
```

Improved code

```
MUL   r10 ← r4, r2
```

Peephole Matching

- Basic idea: Compiler can discover local improvements
 - Look at a small set of adjacent operations
 - Move a “peephole” over code & search for improvement
- Classic examples:
 - store followed by load
 - simple algebraic identities
 - jump to a jump

Original code

```
      Jump    L10  
L10: Jump    L11
```

Improved code

```
L10: Jump    L11
```

Peephole Matching

Implementing it

- Early systems used limited set of hand-coded patterns
- Window size ensured quick processing

Modern peephole instruction selectors

- Break problem into three tasks



- Apply symbolic interpretation & simplification systematically

Peephole Matching

Expander

- Turns IR code into a low-level IR (LLIR) such as RTL*
- Operation-by-operation, template-driven rewriting
- LLIR form includes all direct effects
- Significant, albeit constant, expansion of size



*RTL = Register Transfer Language

Peephole Matching

Simplifier

- Looks at LLIR through window and rewrites it
- Uses forward substitution, algebraic simplification, local constant propagation, dead-effect elimination, ...
- Performs local optimization within window



- This is the heart of the peephole system
 - Benefit of peephole optimization shows up in this step

Peephole Matching

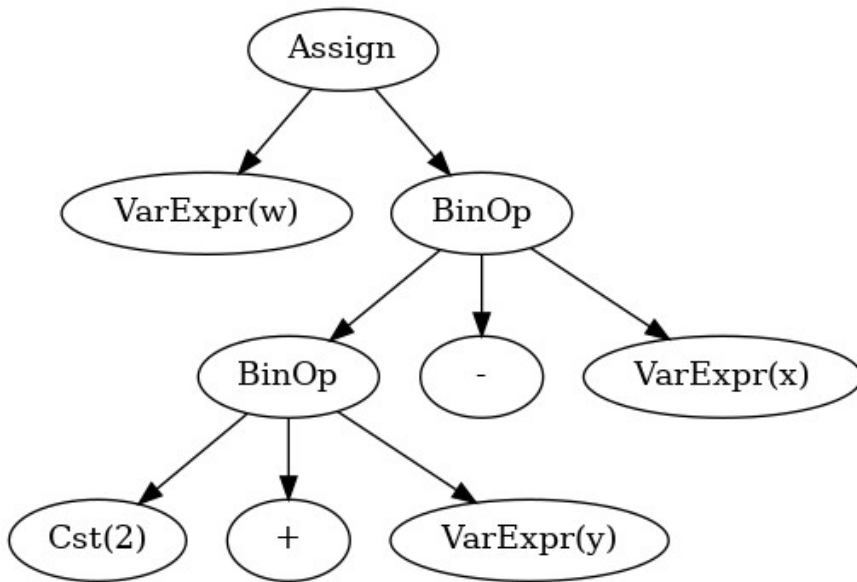
Matcher

- Compares simplified LLIR against a library of patterns
- Picks low-cost pattern that captures effects
- Must preserve LLIR effects, may add new ones
(*e.g., set condition code*)
- Generates the assembly code output

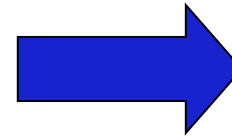


Example: Expander

Original IR



Expand



LLIR Code

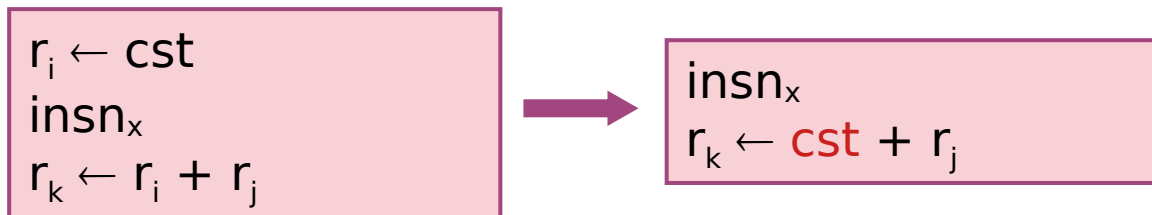
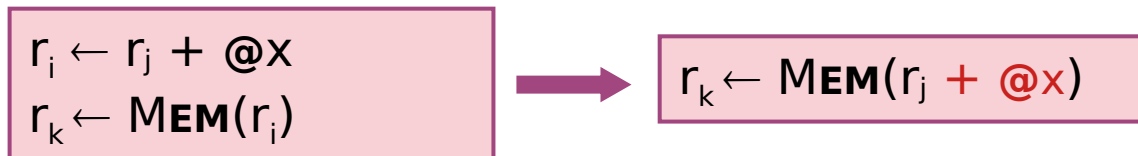
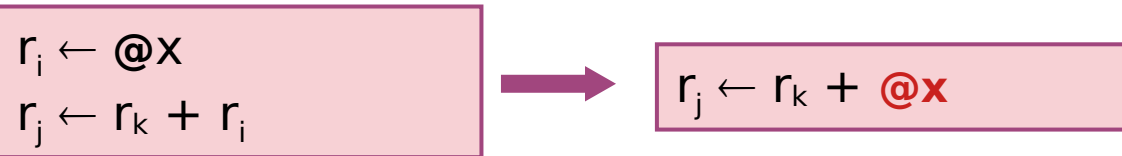
```
r10 ← 2
r11 ← @y
r12 ← fp + r11
r13 ← MEM(r12)
r14 ← r10 + r13
r15 ← @x
r16 ← fp + r15
r17 ← MEM(r16)
r18 ← r14 - r17
r19 ← @w
r20 ← fp + r19
MEM(r20) ← r18
```

assumes x,y,w stack allocated

Each register is
single use

@x, @y, @w = offsets from fp

Example: Simplification rules

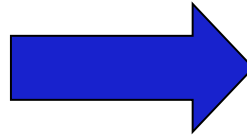


Example: Simplifier

LLIR Code

```
r10 ← 2
r11 ← @y
r12 ← fp + r11
r13 ← MEM(r12)
r14 ← r10 + r13
r15 ← @x
r16 ← fp + r15
r17 ← MEM(r16)
r18 ← r14 - r17
r19 ← @w
r20 ← fp + r19
MEM(r20) ← r18
```

Simplify



LLIR Code

```
r13 ← MEM(fp + @y)
r14 ← 2 + r13
r17 ← MEM(fp + @x)
r18 ← r14 - r17
MEM(fp + @w) ← r18
```

Steps of the Simplifier *(3-operation window)*

LLIR Code

$r_{10} \leftarrow 2$

$r_{11} \leftarrow @y$

$r_{12} \leftarrow fp + r_{11}$

$r_{13} \leftarrow \mathbf{MEM}(r_{12})$

$r_{14} \leftarrow r_{10} + r_{13}$

$r_{15} \leftarrow @x$

$r_{16} \leftarrow fp + r_{15}$

$r_{17} \leftarrow \mathbf{MEM}(r_{16})$

$r_{18} \leftarrow r_{14} - r_{17}$

$r_{19} \leftarrow @w$

$r_{20} \leftarrow fp + r_{19}$

$\mathbf{MEM}(r_{20}) \leftarrow r_{18}$

Steps of the Simplifier *(3-operation window)*

LLIR Code

$r_{10} \leftarrow 2$
 $r_{11} \leftarrow @y$
 $r_{12} \leftarrow fp + r_{11}$
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 $r_{14} \leftarrow r_{10} + r_{13}$
 $r_{15} \leftarrow @x$
 $r_{16} \leftarrow fp + r_{15}$
 $r_{17} \leftarrow \mathbf{MEM}(r_{16})$
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 $r_{19} \leftarrow @w$
 $r_{20} \leftarrow fp + r_{19}$
 $\mathbf{MEM}(r_{20}) \leftarrow r_{18}$

$r_{10} \leftarrow 2$
 $r_{11} \leftarrow @y$
 $r_{12} \leftarrow fp + r_{11}$

Steps of the Simplifier (3-operation window)

LLIR Code

$r_{10} \leftarrow 2$
 $r_{11} \leftarrow @y$
 $r_{12} \leftarrow fp + r_{11}$
 $r_{13} \leftarrow \text{MEM}(r_{12})$
 $r_{14} \leftarrow r_{10} + r_{13}$
 $r_{15} \leftarrow @x$
 $r_{16} \leftarrow fp + r_{15}$
 $r_{17} \leftarrow \text{MEM}(r_{16})$
 $r_{18} \leftarrow r_{14} - r_{17}$
 $r_{19} \leftarrow @w$
 $r_{20} \leftarrow fp + r_{19}$
 $\text{MEM}(r_{20}) \leftarrow r_{18}$

$r_{10} \leftarrow 2$
 $r_{11} \leftarrow @y$
 $r_{12} \leftarrow fp + r_{11}$



$r_{10} \leftarrow 2$
 $r_{11} \leftarrow @y$
 $r_{12} \leftarrow fp + @y$

$r_i \leftarrow @x$
 $r_j \leftarrow r_k + r_i$



$r_j \leftarrow r_k + @x$

Steps of the Simplifier (3-operation window)

LLIR Code

$r_{10} \leftarrow 2$
 $r_{11} \leftarrow @y$
 $r_{12} \leftarrow fp + r_{11}$
 $r_{13} \leftarrow \mathbf{MEM}(r_{12})$
 $r_{14} \leftarrow r_{10} + r_{13}$
 $r_{15} \leftarrow @x$
 $r_{16} \leftarrow fp + r_{15}$
 $r_{17} \leftarrow \mathbf{MEM}(r_{16})$
 $r_{18} \leftarrow r_{14} - r_{17}$
 $r_{19} \leftarrow @w$
 $r_{20} \leftarrow fp + r_{19}$
 $\mathbf{MEM}(r_{20}) \leftarrow r_{18}$

$r_{10} \leftarrow 2$
 $r_{11} \leftarrow @y$
 $r_{12} \leftarrow fp + r_{11}$



$r_{10} \leftarrow 2$
 $r_{12} \leftarrow fp + @y$
 $r_{13} \leftarrow \mathbf{MEM}(r_{12})$

Steps of the Simplifier (3-operation window)

LLIR Code

$r_{10} \leftarrow 2$
 $r_{11} \leftarrow @y$
 $r_{12} \leftarrow fp + r_{11}$
 $r_{13} \leftarrow \text{MEM}(r_{12})$
 $r_{14} \leftarrow r_{10} + r_{13}$
 $r_{15} \leftarrow @x$
 $r_{16} \leftarrow fp + r_{15}$
 $r_{17} \leftarrow \text{MEM}(r_{16})$
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 $r_{19} \leftarrow @w$
 $r_{20} \leftarrow fp + r_{19}$
 $\text{MEM}(r_{20}) \leftarrow r_{18}$

$r_{10} \leftarrow 2$
 $r_{12} \leftarrow fp + @y$
 $r_{13} \leftarrow \text{MEM}(r_{12})$



$r_{10} \leftarrow 2$
 $r_{12} \leftarrow fp + @y$
 $r_{13} \leftarrow \text{MEM}(fp + @y)$

$r_i \leftarrow r_j + @x$
 $r_k \leftarrow \text{MEM}(r_i)$



$r_k \leftarrow \text{MEM}(r_j + @x)$

Steps of the Simplifier (3-operation window)

LLIR Code

$r_{10} \leftarrow 2$
 $r_{11} \leftarrow @y$
 $r_{12} \leftarrow fp + r_{11}$
 $r_{13} \leftarrow \text{MEM}(r_{12})$
 $r_{14} \leftarrow r_{10} + r_{13}$
 $r_{15} \leftarrow @x$
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 $r_{17} \leftarrow \text{MEM}(r_{16})$
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 $r_{19} \leftarrow @w$
 $r_{20} \leftarrow fp + r_{19}$
 $\text{MEM}(r_{20}) \leftarrow r_{18}$

$r_{10} \leftarrow 2$
 $r_{12} \leftarrow fp + @y$
 $r_{13} \leftarrow \text{MEM}(r_{12})$



$r_{10} \leftarrow 2$
 $r_{13} \leftarrow \text{MEM}(fp + @y)$
 $r_{14} \leftarrow r_{10} + r_{13}$

Steps of the Simplifier

(3-operation window)

LLIR Code

$r_{10} \leftarrow 2$
 $r_{11} \leftarrow @y$
 $r_{12} \leftarrow fp + r_{11}$
 $r_{13} \leftarrow \text{MEM}(r_{12})$
 $r_{14} \leftarrow r_{10} + r_{13}$
 $r_{15} \leftarrow @x$
 $r_{16} \leftarrow fp + r_{15}$
 $r_{17} \leftarrow \text{MEM}(r_{16})$
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 $r_{20} \leftarrow fp + r_{19}$
 $\text{MEM}(r_{20}) \leftarrow r_{18}$

$r_{10} \leftarrow 2$
 $r_{13} \leftarrow \text{MEM}(fp + @y)$
 $r_{14} \leftarrow r_{10} + r_{13}$

$r_{10} \leftarrow 2$
 $r_{13} \leftarrow \text{MEM}(fp + @y)$
 $r_{14} \leftarrow 2 + r_{13}$

$r_i \leftarrow \text{cst}$
 insn_x
 $r_k \leftarrow r_i + r_j$

insn_x
 $r_k \leftarrow \text{cst} + r_j$

Steps of the Simplifier (3-operation window)

LLIR Code

$r_{10} \leftarrow 2$
 $r_{11} \leftarrow @y$
 $r_{12} \leftarrow fp + r_{11}$
 $r_{13} \leftarrow \text{MEM}(r_{12})$
 $r_{14} \leftarrow r_{10} + r_{13}$
 $r_{15} \leftarrow @x$
 $r_{16} \leftarrow fp + r_{15}$
 $r_{17} \leftarrow \text{MEM}(r_{16})$
 $r_{18} \leftarrow r_{14} - r_{17}$
 $r_{19} \leftarrow @w$
 $r_{20} \leftarrow fp + r_{19}$
 $\text{MEM}(r_{20}) \leftarrow r_{18}$

$r_{10} \leftarrow 2$
 $r_{13} \leftarrow \text{MEM}(fp + @y)$
 $r_{14} \leftarrow r_{10} + r_{13}$



$r_{13} \leftarrow \text{MEM}(fp + @y)$
 $r_{14} \leftarrow 2 + r_{13}$
 $r_{15} \leftarrow @x$

Steps of the Simplifier

(3-operation window)

LLIR Code

1st op rolling out of window

$r_{10} \leftarrow 2$

$r_{11} \leftarrow @y$

$r_{12} \leftarrow fp + r_{11}$

$r_{13} \leftarrow \text{MEM}(r_{12})$

$r_{14} \leftarrow r_{10} + r_{13}$

$r_{15} \leftarrow @x$

$r_{16} \leftarrow fp + r_{15}$

$r_{17} \leftarrow \text{MEM}(r_{16})$

$r_{18} \leftarrow r_{14} - r_{17}$

$r_{19} \leftarrow @w$

$r_{20} \leftarrow fp + r_{19}$

$\text{MEM}(r_{20}) \leftarrow r_{18}$

$r_{13} \leftarrow \text{MEM}(fp + @y)$

$r_{14} \leftarrow 2 + r_{13}$

$r_{15} \leftarrow @x$

$r_{13} \leftarrow \text{MEM}(fp + @y)$

$r_{14} \leftarrow 2 + r_{13}$

$r_{15} \leftarrow @x$

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Steps of the Simplifier (3-operation window)

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r11 ← @y
r12 ← fp + r11
r13 ← MEM(r12)
r14 ← r10 + r13
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r18 ← r14 - r17
r19 ← @w
r20 ← fp + r19
MEM(r20) ← r18
```

$r_{14} \leftarrow 2 + r_{13}$
 $r_{15} \leftarrow @x$
 $r_{16} \leftarrow fp + r_{15}$



$r_{14} \leftarrow 2 + r_{13}$
 $r_{15} \leftarrow @x$
 $r_{16} \leftarrow fp + @x$

$r_i \leftarrow @x$
 $r_j \leftarrow r_k + r_i$



$r_j \leftarrow r_k + @x$

Steps of the Simplifier *(3-operation window)*

LLIR Code

```
r10 ← 2  
r11 ← @y  
r12 ← fp + r11  
r13 ← MEM(r12)  
r14 ← r10 + r13  
r15 ← @x  
r16 ← fp + r15  
r17 ← MEM(r16)  
r18 ← r14 - r17  
r19 ← @w  
r20 ← fp + r19  
MEM(r20) ← r18
```

$r_{14} \leftarrow 2 + r_{13}$
 $r_{15} \leftarrow @x$
 $r_{16} \leftarrow fp + r_{15}$

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 $r_{16} \leftarrow fp + @x$
 $r_{17} \leftarrow \text{MEM}(r_{16})$

Steps of the Simplifier (3-operation window)

LLIR Code

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```

$r_{14} \leftarrow 2 + r_{13}$
 $r_{16} \leftarrow \text{fp} + @x$
 $r_{17} \leftarrow \text{MEM}(r_{16})$



$r_{14} \leftarrow 2 + r_{13}$
 $r_{17} \leftarrow \text{MEM}(\text{fp} + @x)$
 $r_{18} \leftarrow r_{14} - r_{17}$

$r_i \leftarrow r_j + @x$
 $r_k \leftarrow \text{MEM}(r_i)$



$r_k \leftarrow \text{MEM}(r_j + @x)$

Steps of the Simplifier

(3-operation window)

LLIR Code

$r_{10} \leftarrow 2$

$r_{11} \leftarrow @y$

$r_{12} \leftarrow fp + r_{11}$

$r_{13} \leftarrow \text{MEM}(r_{12})$

$r_{14} \leftarrow r_{10} + r_{13}$

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Steps of the Simplifier

(3-operation window)

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$r_{17} \leftarrow \mathbf{MEM}(fp + @x)$

$r_{18} \leftarrow r_{14} - r_{17}$

$r_{19} \leftarrow @w$

$r_{17} \leftarrow \mathbf{MEM}(fp + @x)$

$r_{18} \leftarrow r_{14} - r_{17}$

$r_{19} \leftarrow @w$

$r_{20} \leftarrow fp + r_{19}$

Steps of the Simplifier (3-operation window)

LLIR Code

```
r10 ← 2
r11 ← @y
r12 ← fp + r11
r13 ← MEM(r12)
r14 ← r10 + r13
r15 ← @x
r16 ← fp + r15
r17 ← MEM(r16)
r18 ← r14 - r17
r19 ← @w
r20 ← fp + r19
MEM(r20) ← r18
```

$r_i \leftarrow @x$
 $r_j \leftarrow r_k + r_i$



$r_j \leftarrow r_k + @x$

$r_{18} \leftarrow r_{14} - r_{17}$
 $r_{19} \leftarrow @w$
 $r_{20} \leftarrow fp + r_{19}$



$r_{18} \leftarrow r_{14} - r_{17}$
 $r_{20} \leftarrow fp + @w$
 $MEM(r_{20}) \leftarrow r_{18}$

Steps of the Simplifier (3-operation window)

LLIR Code

```
r10 ← 2
r11 ← @y
r12 ← fp + r11
r13 ← MEM(r12)
r14 ← r10 + r13
r15 ← @x
r16 ← fp + r15
r17 ← MEM(r16)
r18 ← r14 - r17
r19 ← @w
r20 ← fp + r19
MEM(r20) ← r18
```

$r_i \leftarrow r_j + @x$
 $r_k \leftarrow \text{MEM}(r_i)$



$r_k \leftarrow \text{MEM}(r_j + @x)$

$r_{18} \leftarrow r_{14} - r_{17}$
 $r_{20} \leftarrow \text{fp} + @w$
 $\text{MEM}(r_{20}) \leftarrow r_{18}$



$r_{18} \leftarrow r_{14} - r_{17}$
 $\text{MEM}(\text{fp} + @w) \leftarrow r_{18}$

Steps of the Simplifier (3-operation window)

LLIR Code

```
r10 ← 2
r11 ← @y
r12 ← fp + r11
r13 ← MEM(r12)
r14 ← r10 + r13
r15 ← @x
r16 ← fp + r15
r17 ← MEM(r16)
r18 ← r14 - r17
r19 ← @w
r20 ← fp + r19
MEM(r20) ← r18
```

$r_{18} \leftarrow r_{14} - r_{17}$
 $\text{MEM}(\text{fp} + @w) \leftarrow r_{18}$



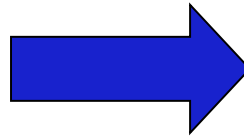
$r_{18} \leftarrow r_{14} - r_{17}$
 $\text{MEM}(\text{fp} + @w) \leftarrow r_{18}$

Example

LLIR Code

```
r10 ← 2
r11 ← @y
r12 ← fp + r11
r13 ← MEM(r12)
r14 ← r10 + r13
r15 ← @x
r16 ← fp + r15
r17 ← MEM(r16)
r18 ← r14 - r17
r19 ← @w
r20 ← fp + r19
MEM(r20) ← r18
```

Simplify



LLIR Code

```
r13 ← MEM(fp + @y)
r14 ← 2 + r13
r17 ← MEM(fp + @x)
r18 ← r14 - r17
MEM(fp + @w) ← r18
```

Example: Matcher

LLIR Code

$r_{13} \leftarrow \text{MEM}(\text{fp} + @y)$

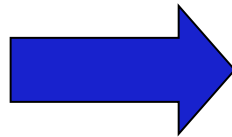
$r_{14} \leftarrow 2 + r_{13}$

$r_{17} \leftarrow \text{MEM}(\text{fp} + @x)$

$r_{18} \leftarrow r_{14} - r_{17}$

$\text{MEM}(\text{fp} + @w) \leftarrow r_{18}$

Match



MIPS Code

LW $r_{13}, @y(\$fp)$

ADDI $r_{14}, r_{13}, 2$

LW $r_{17}, @x(\$fp)$

SUB r_{18}, r_{14}, r_{17}

SW $r_{18}, @w(\$fp)$

- Produces pretty good code

Making It All Work

Details

- LLIR is largely machine independent (RTL)
- Target machine described as LLIR → ASM pattern
- Actual pattern matching
 - Use a hand-coded pattern matcher (GCC)
 - Turn patterns into grammar & use LR parser (VPO)
- Several important compilers use this technology
- It seems to produce good portable instruction selectors

Key strength appears to be late low-level optimization

Next lecture

Instruction selection

- Tree-based pattern matching (LLVM)