

Compiler Design

An Introduction to Equality Saturation and Its Applications

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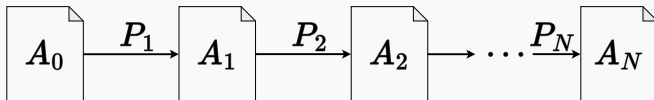
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Pass Sequences

An optimizer performs a sequence of passes (a.k.a. transformations, rewrites) on the program to find the optimal version.





Optimization Passes

- Constant folding
- Dead code elimination
- Common subexpression elimination
- Loop unrolling
- Loop fusion
- Loop tiling
- Loop vectorization



Phase Ordering Problem

The order in which passes are applied affects the final program and its performance.



Phase Ordering Problem

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- $P_1. \underline{xy/z \rightarrow x(y/z)}$
- $P_2. \underline{x \cdot 2 \rightarrow x \ll 1}$
- $P_3. \underline{xy \rightarrow yx}$
- $P_5. \underline{x/x \rightarrow 1}$
- $P_6. \underline{x \cdot 1 \rightarrow x}$



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- $(a \cdot 2)/2 \xrightarrow{P_2} (a \ll 1)/2$
- $(a \cdot 2)/2 \xrightarrow{P_1} a \cdot (2/2) \xrightarrow{P_5} a \cdot 1 \xrightarrow{P_6} a$



Phase Ordering Problem

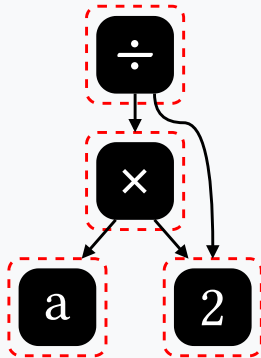
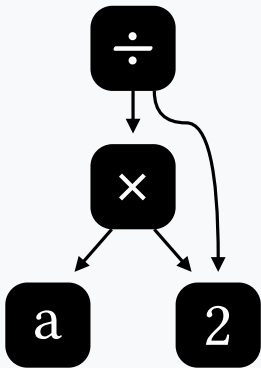
- $(a \cdot 2)/2 \xrightarrow{P_2} (a \ll 1)/2$
- $(a \cdot 2)/2 \xrightarrow{P_1} a \cdot (2/2) \xrightarrow{P_5} a \cdot 1 \xrightarrow{P_6} a$

Problems:

1. The compiler gets stuck in a local optimum if it applies P_2 before P_1 .
2. The passes are destructive: we lose the original and intermediate versions of the program.

Equivalence Graphs

We build an **e-graph** to store all defined equivalent versions of the program.

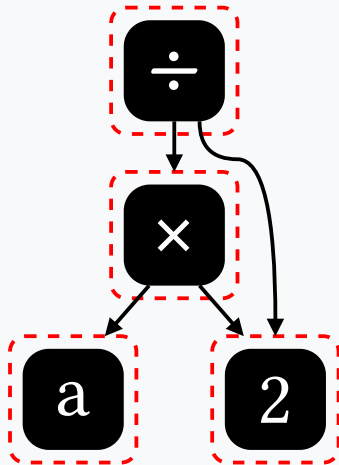


$$(a \cdot 2) / 2$$

Equivalence Graphs

Store many equivalent versions of a program.

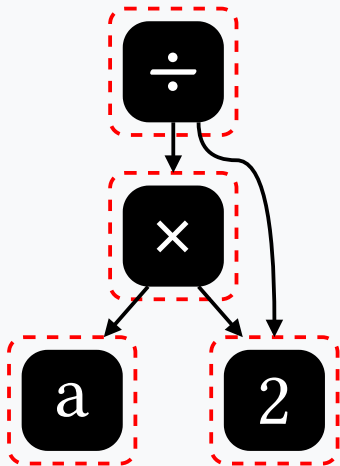
1. **e-node**: Operation node.
2. **e-class**: Set of equivalent e-nodes.
3. **e-graph**: Set of e-classes.



Equality Saturation

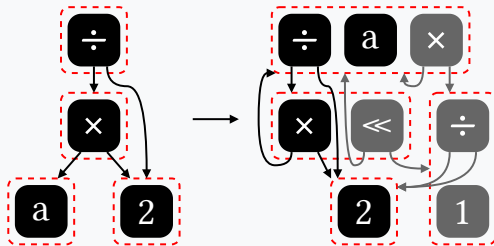
Process:

1. Build an **e-graph** of $(a \cdot 2)/2$.
2. Apply the rewrite rules to the e-graph until a fixed-point.
3. Extract the optimized expression via a cost model.

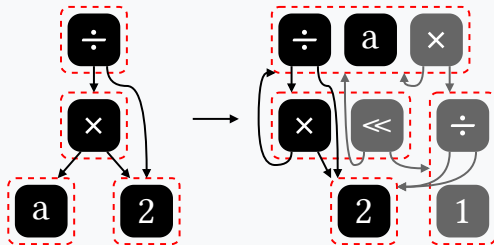


Equality Saturation

- $P_1. \frac{xy/z \rightarrow x(y/z)}{}$
- $P_2. \frac{x \cdot 2 \rightarrow x \lll 1}{}$
- $P_3. \frac{xy \rightarrow yx}{}$
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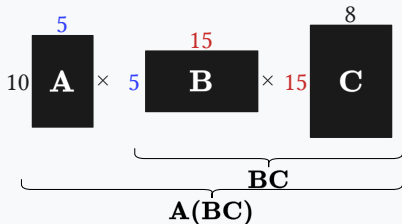
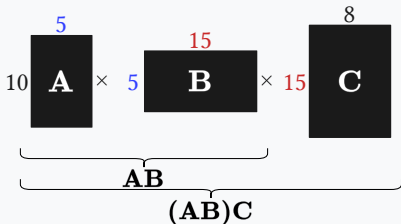
Equality Saturation



Problems are solved:

1. The global optimum is found by extracting the optimal expression via a cost model.
2. The original and intermediate versions of the program are preserved.

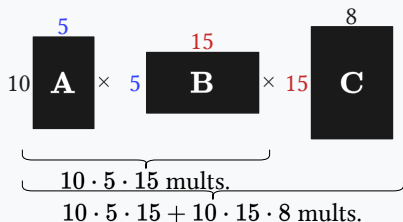
MatMul Associativity



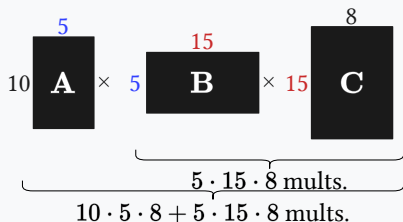
$$(AB)C = A(BC)$$

How do we find out which one is more efficient?

The Order of Operations Matters



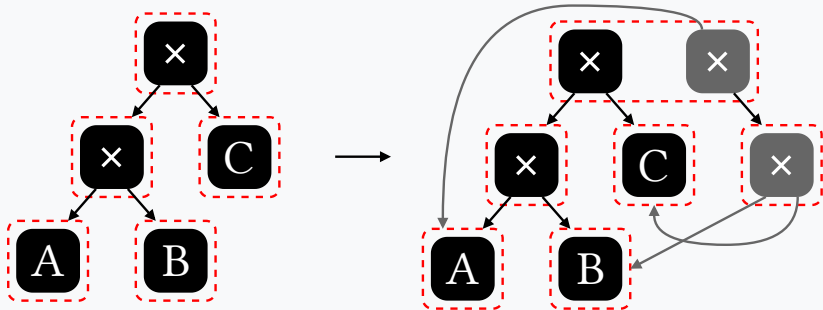
cost **(AB) C** = 1950
multiplications.



cost **A (BC)** = 1000
multiplications.

MatMul Associativity

$$(AB)C = A(BC)$$





Polynomial Evaluation: Horner's Method

Reduce from $O(n^2)$ to $O(n)$ multiplications.

$$\begin{aligned} P(x) &= a_0 + a_1x + a_2x^2 + \cdots + a_{n-1}x^{n-1} + a_nx^n \\ &= a_0 + x(a_1 + x(a_2 + \cdots + x(a_{n-1} + xa_n))) \end{aligned}$$



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- Exponentiation: $x^0 \Rightarrow 1$ and $x^n \Rightarrow x \cdot x^{n-1}$

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- Exponentiation: $x^0 \Rightarrow 1$ and $x^n \Rightarrow x \cdot x^{n-1}$
- Commutativity: $x + y \Leftrightarrow y + x$ and $x \cdot y \Leftrightarrow y \cdot x$
- Associativity: $(x + y) + z \Leftrightarrow x + (y + z)$ and $(x \cdot y) \cdot z \Leftrightarrow x \cdot (y \cdot z)$
- Distributivity: $x \cdot (y + z) \Leftrightarrow x \cdot y + x \cdot z$
- Identity: $x \cdot 1 \Rightarrow x$



References

- Equality Saturation: a New Approach to Optimization. [Tate et al., 2009]
- DialEgg: Dialect-Agnostic MLIR Optimizer using Equality Saturation with Egglog. [Zayed and Dubach, 2025]
- Latent Idiom Recognition for a Minimalist Functional Array Language using Equality Saturation. [der Cruysse and Dubach, 2023]



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Dialegg: Dialect-agnostic mlir optimizer using equality saturation with egglog.

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