Compiler Design

Lecture 11: A Brief Tour of MIPS assembly

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Overview

MIPS assembly program sections

.data

Data section: constant and variable definitions go here (including statically allocated arrays)

- format for declarations: name: storage_type size/value
- · create storage for a variable:

```
    var1: .word 3 # one word storage with initial value 3
    arr1: .space 40 # 40 bytes storage (no initial value)
```

.text

Text section: assembly instructions go here

Components of an assembly program

Category	Example
Comment	# I am a comment
Assembler directives	.data, .asciiz
Operation mnemonic	add, addi, lw, bne
Register name	\$zero, \$t3
Address label (declaration)	loop1:
Address label (use)	loop1
Integer constant	8, -4, 0xA9
Character constant	'h', '\t'
String constant	"Hello, world\n"

Hello world example

```
# A simple hello world program
.data
hellostr: .asciiz "Hello, world\n"
.text
 li $v0, 4  # setup print syscall
 la $a0, hellostr # argument to print string
                 # tell the OS to do the system call
 syscall
 li $v0, 10  # setup exit syscall
 syscall
                    # tell the OS to perform the syscall
```

Registers

Registers

- · 32 general-purpose registers
- register preceded by \$ in MIPS assembly language
- two formats: name (\$zero) or number (\$0)
- holds 32 bits value (= 4 bytes = 1 word)
- stack grows from high memory to low memory

Registers

Register	Alternative	Description
number	name	
0	\$zero	the value 0
1	\$at	assembler temporary: reserved by the assembler
2-3	\$v0-\$v1	function return values
4-7	\$a0-\$a3	function arguments: first four parameters for function (no pre-
		served across function call)
8-15	\$t0-\$t7	temporaries (not preserved across function calls)
16-23	\$s0-\$s7	saved temporaries (preserved across function calls)
24-25	\$t8-\$t9	temporaries: (not preserved across function calls)
26-27	\$k0-\$k1	reserved for use by the interrupt/trap handler
28	\$gp	global pointer : base of global data segment
29	\$sp	stack pointer : points to last location on stack
30	\$s8/\$fp	saved value / frame pointer (preserved across function call)
31	\$ra	return address

 Special Hi and Lo registers (not shown above) holds result of multiplication and division (see example later)

Arithmetic

Arithmetic Instructions

- · Most use three operands
- All operands are registers or immediate values (Load/Store architecture)
- · All operands are 4 bytes (a word)

Arithmetic & Move Instruction Examples

```
add
        $t0,$t1,$t2 # $t0 = $t1 + $t2;
                   # add as signed (2's complement) integers
        $t2,$t3,$t4 # $t2 = $t3 - $t4
sub
       $t2,$t3, 5 # $t2 = $t3 + 5; "add immediate"
addi
addu
       $t1,$t6,$t7 # $t1 = $t6 + $t7;
                                        add as unsigned integers
subu
       $t1,$t6,$t7 # $t1 = $t6 + $t7; subtract as unsigned integers
mul
        $t1, $t2, $t3 # $t1 = $t2 * $t3; multiply
mult
       $t3,$t4
# multiply 32-bit quantities in $t3 and $t4, and store 64-bit
# result in special registers Lo and Hi: (Hi,Lo) = $t3 * $t4
        $t5,$t6 # Lo = $t5 / $t6 (integer quotient)
div
               # Hi = $t5 mod $t6 (remainder)
mfhi
        $t0 # move value from special register Hi to $t0: $t0 = Hi
mflo
        $t1 # move value from special register Lo to $t1:
                                                           $t1 = Lo
       $t2,$t3 # $t2 = $t3
move
```

Memory

Load / Store Instructions

- Memory access only allowed with explicit load and store instructions (load/store architecture)
- · All other instructions use register operands
- · Load
 - lw register_destination, mem_source copy a word (4 bytes) at source memory location to destination register
 - lb register_destination, mem_source
 copy a byte to low-order byte of destination register (sign extend higher-order bytes)
 - li register_destination, value load immediate value into destination register (pseudo-instruction)

- Store
 - sw register_source, mem_destination
 store a word (4 bytes) from source register to memory location
 - sb register_source, mem_destination
 store a byte (low-order) from source register to memory location

```
Lata
var1: .word 23 # declare storage for var1; initial value is 23

.text
lw $t0, var1 # load content of mem location into register $t0: $t0 = 23
li $t1, 5 # $t1 = 5 ("load immediate")
sw $t1, var1 # store content of $t1 into mem: *var1 = 5
```

A var1 represents a pointer to a word since it is an address.

Indirect and Based Addressing

- · load address:
 - la \$t0, var1 copy memory address of var1 into register \$t0
- · indirect addressing:
 - lw \$t1, (\$t0)
 load word at memory address contained in \$t0 into \$t2
 - sw \$t2, (\$t0)
 store word in register \$t2 into memory at address contained in \$t0
- based/indexed addressing (useful for field access in struct):
 - lw \$t2, 4(\$t0) load word at memory address (\$t0+4) into register \$t2
 - sw \$t2, -12(\$t0) store content of register \$t2 into memory at address (\$t0-12)

```
Example
```

```
.data
array1: .space 12 # declare 12 bytes of storage
.text
la $t0, array1 # load base address of array into $t0
li $t1, 5  # $t1 = 5 ("load immediate")
sw $t1, ($t0) # first array element set to 5
li $t1, 13 # $t1 = 13
sw $t1, 4($t0) # second array element set to 13
li $t1, -7 # $t1 = -7
sw $t1, 8($t0) # third array element set to -7
```

Exercise

Write a MIPS assembly program corresponding to the following C code:

```
struct point_t {
 int x;
 int y;
};
struct point_t p;
int arr[12];
void foo() {
 p.x = 2;
 p.y = 5;
 arr[3] = 6;
```

Control Structures

Control structures

· Branches:

```
b target # unconditional branch to target
beq $t0,$t1,target # branch to target if $t0 = $t1
blt $t0,$t1,target # branch to target if $t0 < $t1
ble $t0,$t1,target # branch to target if $t0 <= $t1
bgt $t0,$t1,target # branch to target if $t0 > $t1
bge $t0,$t1,target # branch to target if $t0 >= $t1
bne $t0,$t1,target # branch to target if $t0 >= $t1
bne $t0,$t1,target # branch to target if $t0 <> $t1
```

```
Example

addi $t0, $zero, 0 # t0 = 0
addi $t1, $zero, 10 # t1 = 10

loop:
addi $t0, $t0, 1 # t0 = t0+1
blt $t0, $t1, loop # branch to loop if t0<t1 (t0<10)
```

Control structures

· Jumps:

```
j target
# unconditional jump to program label target

jr $t3
# jump to address contained in $t3 ("jump register")
```

· Subroutine (function) call:

```
jal label # "jump and link"
```

- copy program counter (return address) to register \$ra (return address register)
- jump to program instruction at label

```
jr $ra # "jump register"
```

• jump to return address in \$ra (stored by jal instruction)

In case of nested function calls, the return address should be saved to the stack and restored accordingly.

System Calls

System Calls (MIPS simulator)

System calls are used to interface with the operating systems. For instance input/output or dynamic memory allocation.

Using system calls:

- 1. load the service number in register \$v0
- 2. load argument values in \$a0, \$a1, ...
- 3. issue the syscall instruction
- 4. retrieve return value if any

Example: printing integer on the console

```
li $v0, 1
# service 1 is print integer
add $a0, $t0, $zero
# load desired value into argument register $a0
syscall
```

System calls table

Service	\$v0	Arguments	Result
print integer	1	\$a0 = integer to print	
print string	4	\$a0 = address of null- terminated string to print	
print character	11	\$a0 = character to print	
read integer	5		\$v0 = integer read
read character	12		\$v0 = character read
allocate heap	9	\$a0 = number of bytes to al-	\$v0 = address of
memory		locate	allocated memory

Next lecture:

· Introduction to Code Generation