

# Compiler Design

## Lecture 18: Instruction Selection via Peephole Matching

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Winter 2024

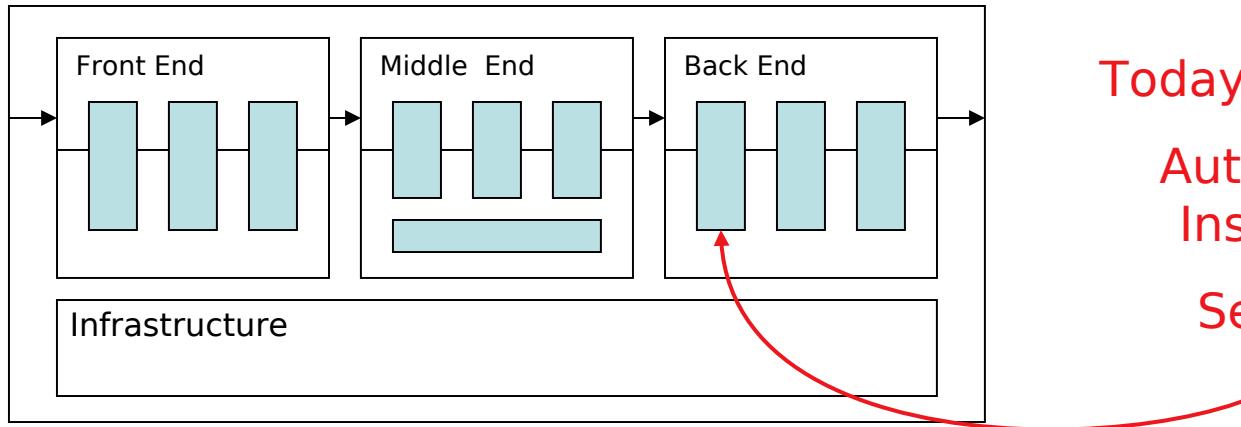
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# The Problem

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Writing a compiler is a lot of work

- Would like to reuse components whenever possible
- Would like to **automate** construction of components



Today's lecture:  
Automating  
Instruction  
Selection

- Front end construction is largely automated
- Middle is largely hand crafted
- (Parts of) back end can be automated

# Definitions

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## Instruction selection

- Process of mapping **IR** into assembly code
  - Assumes a fixed storage mapping & code shape
  - Combining operations, using address modes

## Instruction scheduling

- Process of reordering operations to hide latencies
  - Assumes a fixed program (*set of operations*)
  - Changes demand for registers

## Register allocation

- Process of deciding which values will reside in registers
  - Changes the storage mapping, may add false sharing
  - Concerns about placement of data & memory operations

# The Problem

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Modern computers have many ways to do anything

Consider register-to-register copy

- Obvious operation is  $\text{MOVE } r_j \leftarrow r_i$
- Many others exist

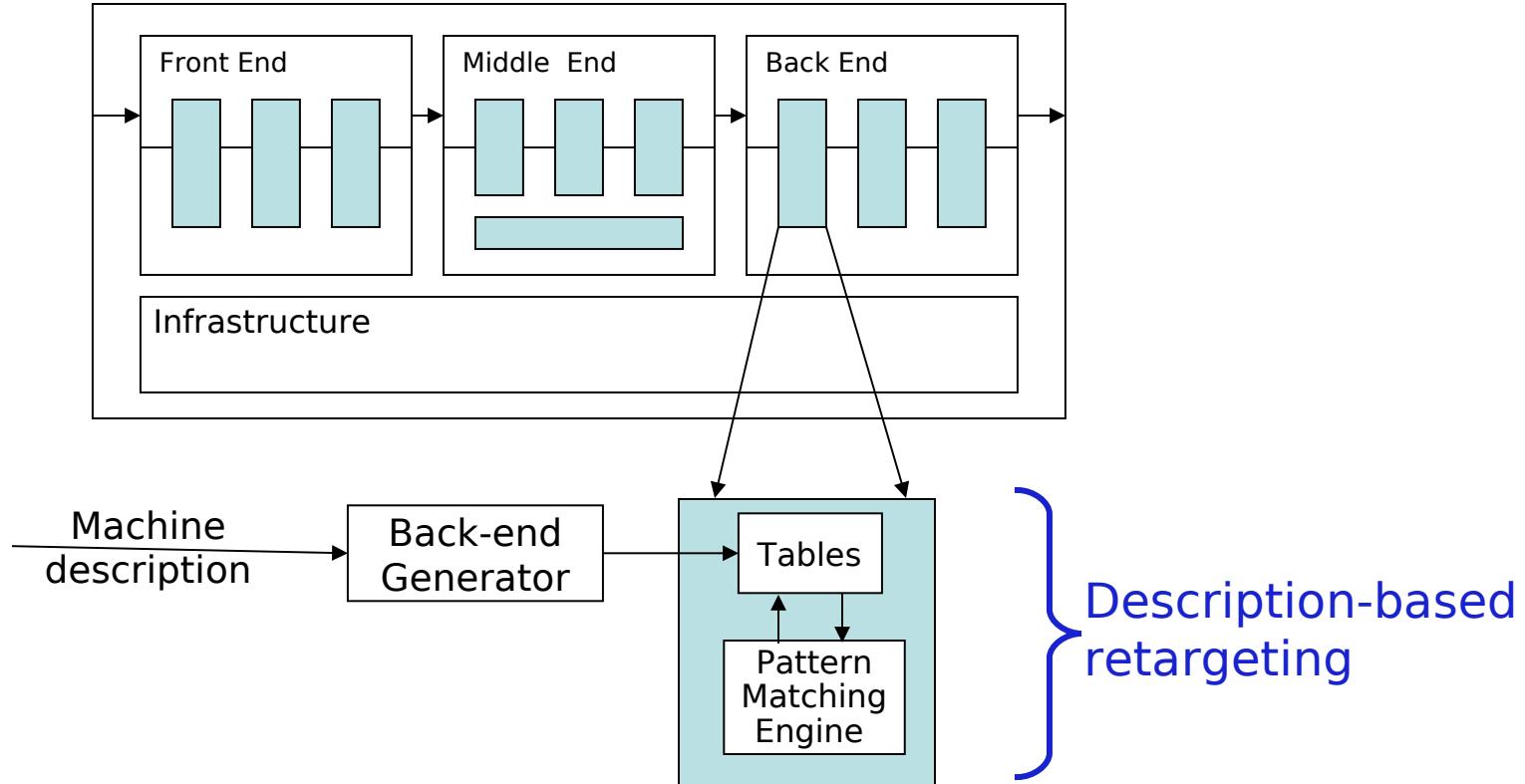
$\text{ADDI } r_j \leftarrow r_i, 0$	$\text{OR } r_j \leftarrow r_i, r_0$	$\text{SLL } r_j \leftarrow r_i, 0$
$\text{ADD } r_j \leftarrow r_i, r_0$	$\text{MULI } r_j \leftarrow r_i, 1$	$\text{RSHIFTI } r_j \leftarrow r_i, 0$
$\text{ORI } r_j \leftarrow r_i, 0$	$\text{XORI } r_j \leftarrow r_i, 0$	$\dots$ and others ...

- Human would ignore all of these
- Algorithm must look at all of them & find low-cost encoding
  - Take context into account *(busy functional unit?)*

And this is an overly-simplified example

# The Goal

Want to automate generation of instruction selectors



Machine description should also help with scheduling & allocation

# The Big Picture

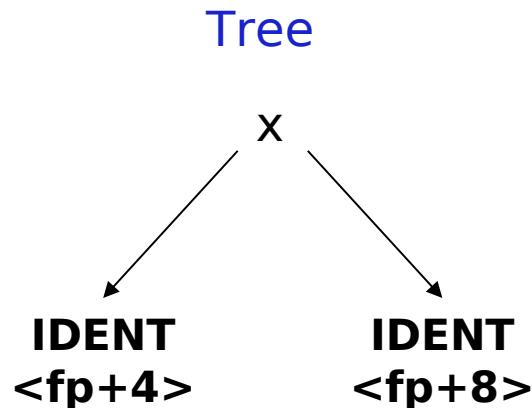
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Need pattern matching techniques

- Must produce good code *(some metric for good )*
- Must run quickly

Our treewalk (visitor) code generator ran quickly

How good was the code?



# The Big Picture

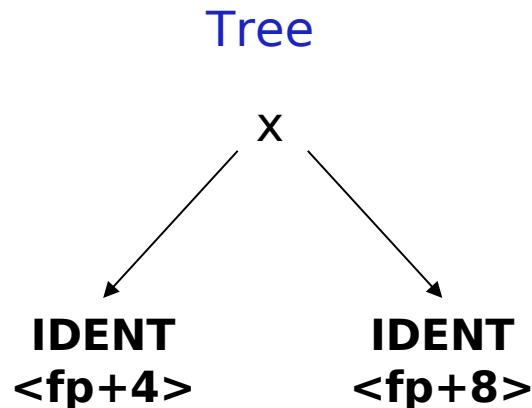
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Need pattern matching techniques

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How good was the code?



Treewalk Code

```
MOVE r1 ← 4  
ADD r2 ← r1 , $fp  
LW r3 ← 0(r2)  
MOVE r4 ← 8  
ADD r5 ← r4 , $fp  
LW r6 ← 0(r5)  
MUL r7 ← r3, r6
```

# The Big Picture

Need pattern matching techniques

- Must produce good code *(some metric for good )*
- Must run quickly

Our treewalk (visitor) code generator ran quickly

How good was the code?

Tree	Treewalk Code	Desired Code
<pre>graph TD; x[x] --&gt; IDENT1[IDENT&lt;fp+4&gt;]; x --&gt; IDENT2[IDENT&lt;fp+8&gt;]</pre>	<pre>MOVE r<sub>1</sub> ← 4 ADD r<sub>2</sub> ← r<sub>1</sub>, \$fp LW r<sub>3</sub> ← 0(r<sub>2</sub>) MOVE r<sub>4</sub> ← 8 ADD r<sub>5</sub> ← r<sub>4</sub>, \$fp LW r<sub>6</sub> ← 0(r<sub>5</sub>) MUL r<sub>7</sub> ← r<sub>3</sub>, r<sub>6</sub></pre>	<pre>LW r<sub>3</sub> ← 4(\$fp) LW r<sub>6</sub> ← 8(\$fp) MUL r<sub>7</sub> ← r<sub>3</sub>, r<sub>6</sub></pre>

# The Big Picture

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Okay wasn't too hard, could do it in the code generator

# The Big Picture

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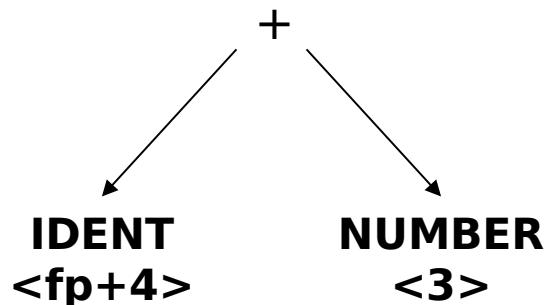
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How good was the code?

Tree



# The Big Picture

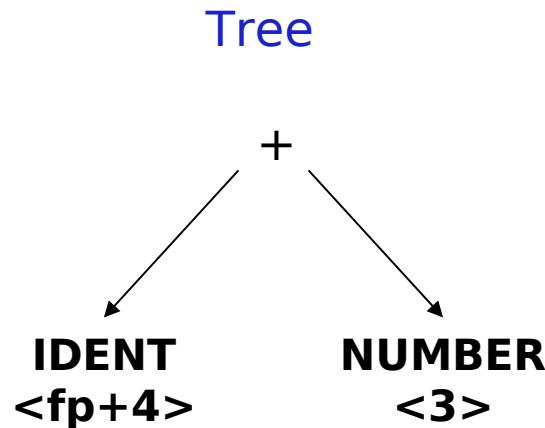
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Need pattern matching techniques

- Must produce good code *(some metric for good )*
- Must run quickly

Our treewalk (visitor) code generator ran quickly

How good was the code?



Treewalk Code

```
MOVE r1 ← 4  
ADD r2 ← r1 , $fp  
LW r3 ← 0(r2)  
MOVE r4 ← 3  
ADD r5 ← r3 , r4
```

# The Big Picture

---

Need pattern matching techniques

- Must produce good code *(some metric for good )*
- Must run quickly

Our treewalk (visitor) code generator ran quickly

How good was the code?

Tree	Treewalk Code	Desired Code
<pre>graph TD; "+" --- IDENT["IDENT&lt;fp+4&gt;"]; "+" --- NUMBER["NUMBER&lt;3&gt;"]</pre>	<pre>MOVE r<sub>1</sub> ← 4 ADD r<sub>2</sub> ← r<sub>1</sub> , \$fp LW r<sub>3</sub> ← 0(r<sub>2</sub>) MOVE r<sub>4</sub> ← 3 ADD r<sub>5</sub> ← r<sub>3</sub> , r<sub>4</sub></pre>	<pre>LW r<sub>3</sub> ← 4(\$fp) ADDI r<sub>5</sub> ← r<sub>3</sub> , 3</pre>

# The Big Picture

Need pattern matching techniques

- Must produce good code *(some metric for good )*
- Must run quickly

Our treewalk (visitor) code generator ran quickly

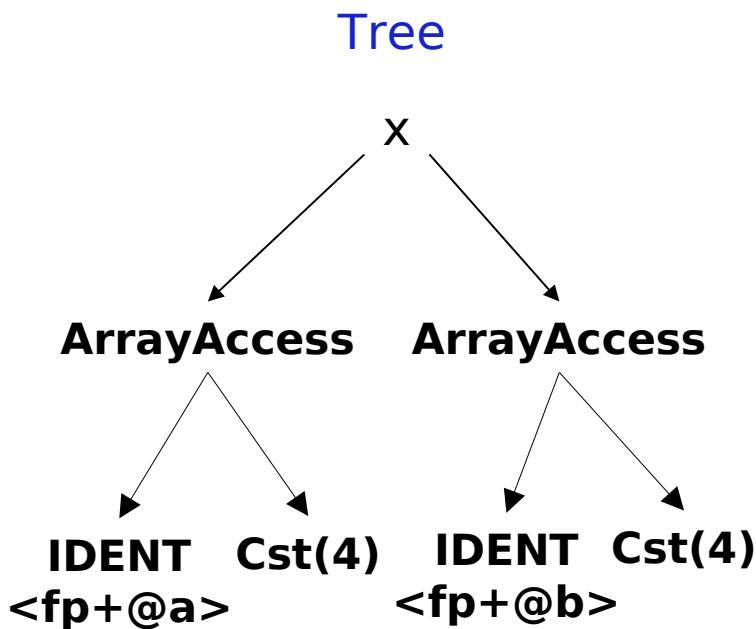
How good was the code?

Tree	Treewalk Code	Desired Code
<p><b>IDENT</b> <code>&lt;fp+4&gt;</code></p> <p><b>NUMBER</b> <code>&lt;3&gt;</code></p>	<pre>MOVE r<sub>1</sub> ← 4 ADD r<sub>2</sub> ← r<sub>1</sub> , \$fp LW r<sub>3</sub> ← 0(r<sub>2</sub>) MOVE r<sub>4</sub> ← 3 ADD r<sub>5</sub> ← r<sub>3</sub> , r<sub>4</sub></pre>	<pre>LW r<sub>3</sub> ← 4(\$fp) ADDI r<sub>5</sub> ← r<sub>3</sub> , 3</pre>

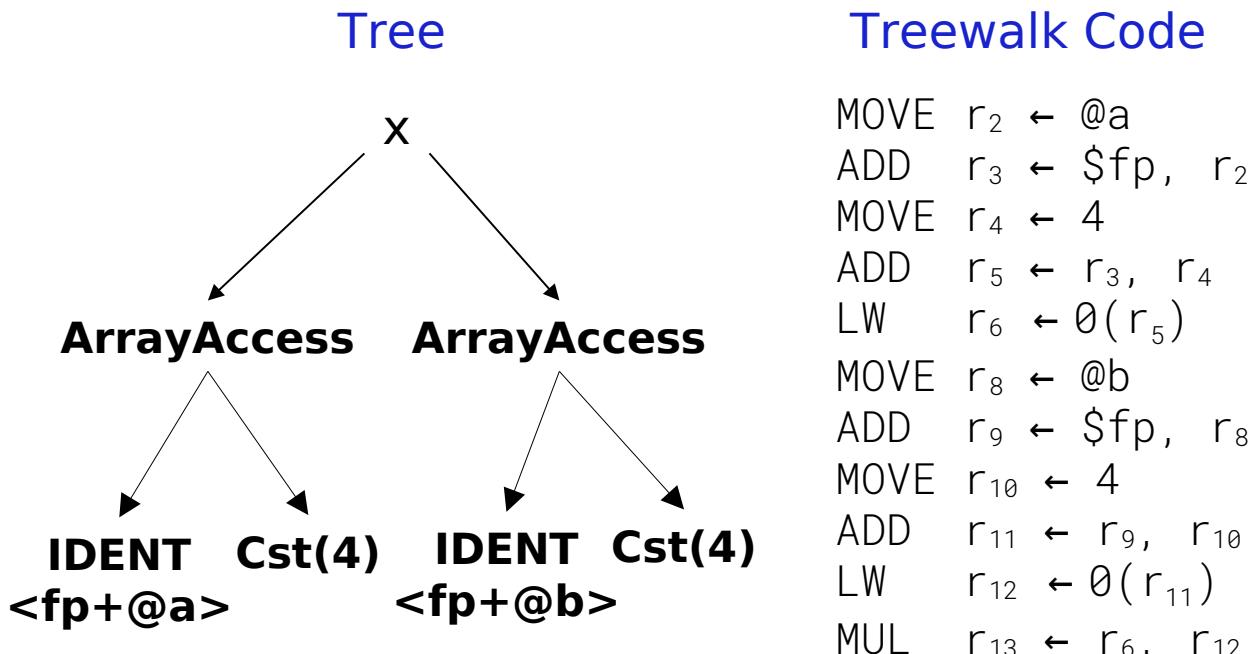
Must combine these.  
This is a **nonlocal** problem

# The Big Picture

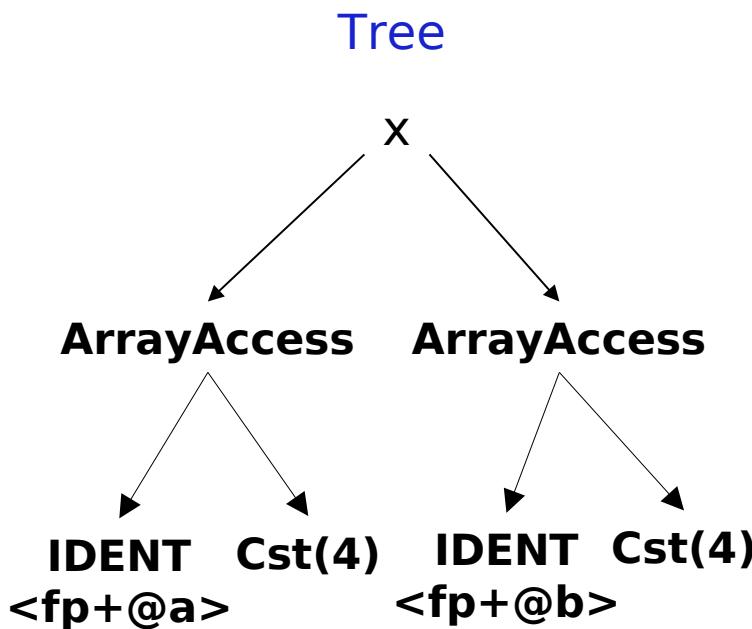
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# The Big Picture



# The Big Picture



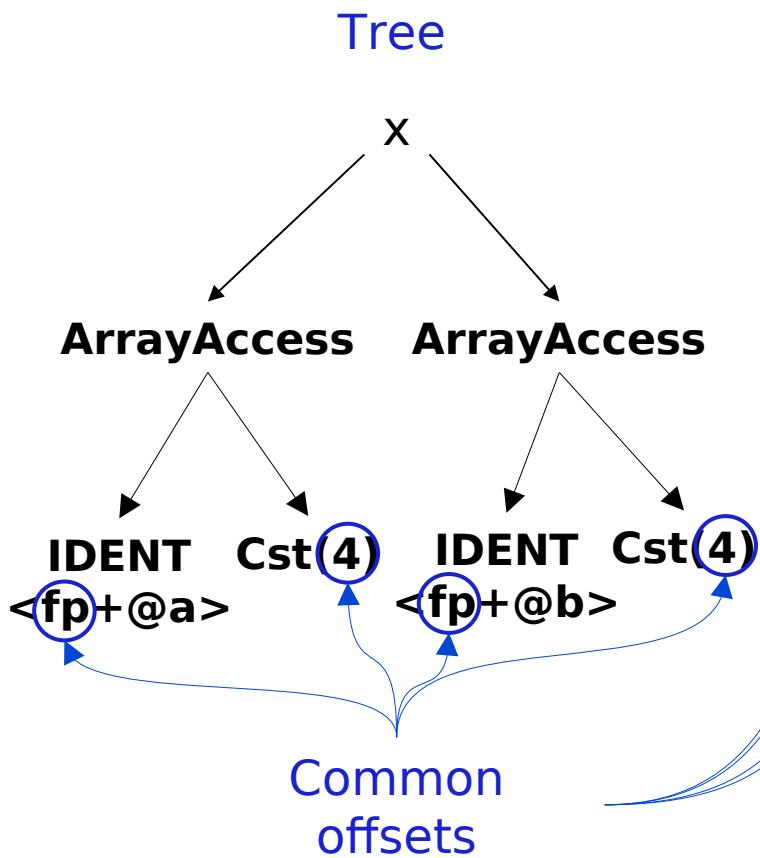
## Treewalk Code

```
MOVE r2 ← @a  
ADD r3 ← $fp, r2  
MOVE r4 ← 4  
ADD r5 ← r3, r4  
LW r6 ← 0(r5)  
MOVE r8 ← @b  
ADD r9 ← $fp, r8  
MOVE r10 ← 4  
ADD r11 ← r9, r10  
LW r12 ← 0(r11)  
MUL r13 ← r6, r12
```

## Desired Code

```
ADDI r3 ← $fp, 4  
LW r6 ← @a(r3)  
LW r12 ← @b(r3)  
MUL r13 ← r6, r12
```

# The Big Picture



## Treewalk Code

```

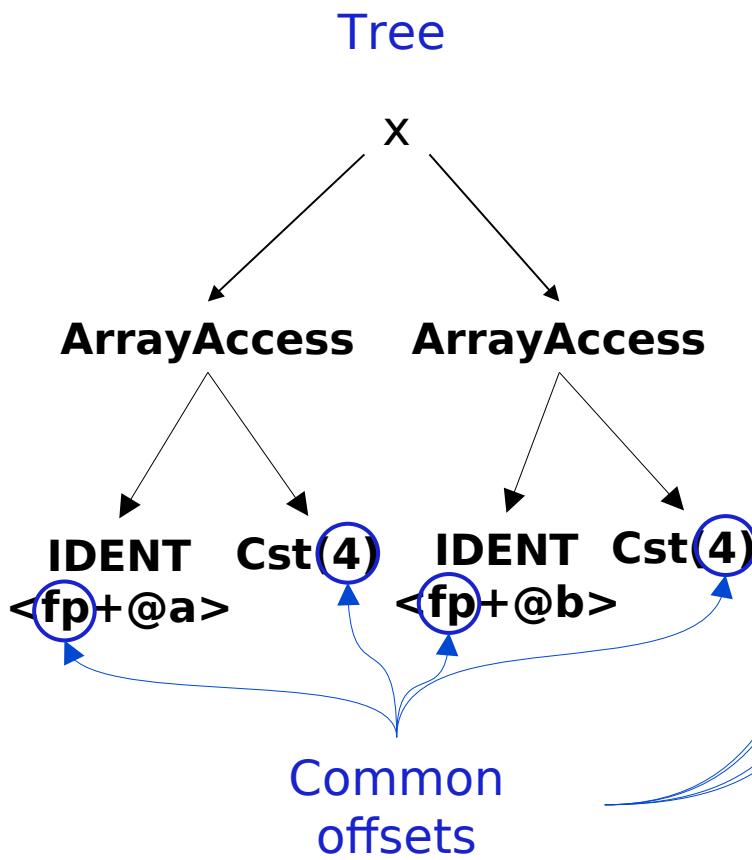
MOVE r2 ← @a
ADD r3 ← $fp, r2
MOVE r4 ← 4
ADD r5 ← r3, r4
LW r6 ← 0(r5)
MOVE r8 ← @b
ADD r9 ← $fp, r8
MOVE r10 ← 4
ADD r11 ← r9, r10
LW r12 ← 0(r11)
MUL r13 ← r6, r12
  
```

## Desired Code

```

ADDI r3 ← $fp, 4
LW r6 ← @a(r3)
LW r12 ← @b(r3)
MUL r13 ← r6, r12
  
```

# The Big Picture



## Treewalk Code

```

MOVE r2 ← @a
ADD r3 ← $fp, r2
MOVE r4 ← 4
ADD r5 ← r3, r4
LW r6 ← 0(r5)
MOVE r8 ← @b
ADD r9 ← $fp, r8
MOVE r10 ← 4
ADD r11 ← r9, r10
LW r12 ← 0(r11)
MUL r13 ← r6, r12
  
```

## Desired Code

```

ADDI r3 ← $fp, 4
LW r6 ← @a(r3)
LW r12 ← @b(r3)
MUL r13 ← r6, r12
  
```

$$($fp + @a) + 4 = ($fp + 4) + @a$$

$$($fp + @b) + 4 = ($fp + 4) + @b$$

Again, a nonlocal problem

# How do we perform this kind of matching ?

---

Tree-oriented IR suggests pattern matching on trees

- Tree-patterns as input, matcher as output
- Each pattern maps to a target-machine instruction sequence
- Use dynamic programming or bottom-up rewrite systems

Linear IR suggests using some sort of string matching

- Strings as input, matcher as output
- Each string maps to a target-machine instruction sequence
- Use text matching (Aho-Corasick) or peephole matching

In practice, both work well. Today we will look at matchers on Linear IR.

# Peephole Matching

---

- Basic idea:  
Compiler can discover local-*ish* improvements
  - Look at a small set of adjacent operations
  - Move a “peephole” over code & search for improvement



Marco Verch , CC BY 2.0 Source: <https://foto.wuestenigel.com/a-mans-hand-opens-the-peephole-on-the-door/>

# Peephole Matching

---

- Basic idea: Compiler can discover local improvements
  - Look at a small set of adjacent operations
  - Move a “peephole” over code & search for improvement
- Classic examples:
  - store followed by load

Original code

```
SW    r1 → 0(r2)
LW    r15 ← 0(r2)
```

Improved code

```
SW    r1 → 0(r2)
MOVE r15 ← r1
```

# Peephole Matching

---

- Basic idea: Compiler can discover local improvements
  - Look at a small set of adjacent operations
  - Move a “peephole” over code & search for improvement
- Classic examples:
  - store followed by load
  - simple algebraic identities

Original code

```
ADDI  r7 ← r2, 0  
MUL   r10 ← r4, r7
```

Improved code

```
MUL   r10 ← r4, r2
```

# Peephole Matching

---

- Basic idea: Compiler can discover local improvements
  - Look at a small set of adjacent operations
  - Move a “peephole” over code & search for improvement
- Classic examples:
  - store followed by load
  - simple algebraic identities
  - jump to a jump

Original code

Jump       $L_{10}$   
 $L_{10}$ : Jump       $L_{11}$

Improved code

$L_{10}$ : Jump       $L_{11}$

# Peephole Matching

## Implementing it

- Early systems used limited set of hand-coded patterns
- Window size ensured quick processing

## Modern peephole instruction selectors

- Break problem into three tasks



- Apply symbolic interpretation & simplification systematically

# Peephole Matching

## Expander

- Turns IR code into a low-level IR (LLIR) such as RTL\*
- Operation-by-operation, template-driven rewriting
- LLIR form includes all direct effects
- Significant, albeit constant, expansion of size



\*RTL = Register Transfer Language

# Peephole Matching

## Simplifier

- Looks at LLIR through window and rewrites it
- Uses forward substitution, algebraic simplification, local constant propagation, dead-effect elimination, ...
- Performs local optimization within window



- This is the heart of the peephole system
  - Benefit of peephole optimization shows up in this step

# Peephole Matching

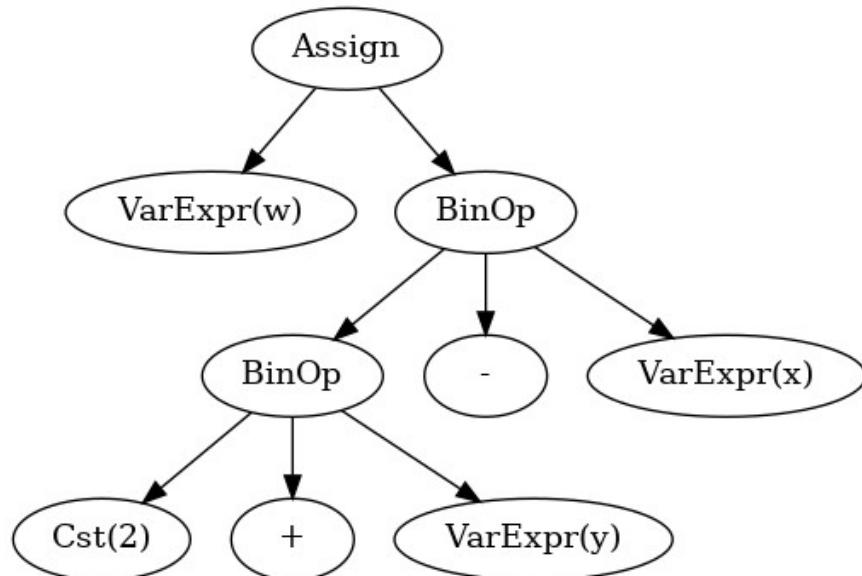
## Matcher

- Compares simplified LLIR against a library of patterns
- Picks low-cost pattern that captures effects
- Must preserve LLIR effects, may add new ones  
(e.g., *set condition code*)
- Generates the assembly code output



# Example: Expander

Original IR



Expand

LLIR Code

```
r10 ← 2  
r11 ← @y  
r12 ← fp + r11  
r13 ← MEM(r12)  
r14 ← r10 + r13  
r15 ← @x  
r16 ← fp + r15  
r17 ← MEM(r16)  
r18 ← r14 - r17  
r19 ← @w  
r20 ← fp + r19  
MEM(r20) ← r18
```

assumes x,y,w stack allocated

Each register is  
**single use**

@x, @y, @w = offsets from fp

## Example: Simplification rules

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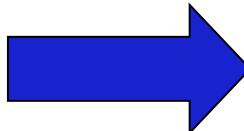
$$r_i \leftarrow @x$$
$$r_j \leftarrow r_k + r_i$$
$$r_j \leftarrow r_k + @x$$
$$r_i \leftarrow r_j + @x$$
$$r_k \leftarrow \text{MEM}(r_i)$$
$$r_k \leftarrow \text{MEM}(r_j + @x)$$
$$r_i \leftarrow \text{cst}$$
$$\text{insn}_x$$
$$r_k \leftarrow r_i + r_j$$
$$\text{insn}_x$$
$$r_k \leftarrow \text{cst} + r_j$$

# Example: Simplifier

## LLIR Code

```
r10 ← 2  
r11 ← @y  
r12 ← fp + r11  
r13 ← MEM(r12)  
r14 ← r10 + r13  
r15 ← @x  
r16 ← fp + r15  
r17 ← MEM(r16)  
r18 ← r14 - r17  
r19 ← @w  
r20 ← fp + r19  
MEM(r20) ← r18
```

Simplify



LLIR Code

```
r13 ← MEM(fp + @y)  
r14 ← 2 + r13  
r17 ← MEM(fp + @x)  
r18 ← r14 - r17  
MEM(fp + @w) ← r18
```

# Steps of the Simplifier *(3-operation window)*

---

## LLIR Code

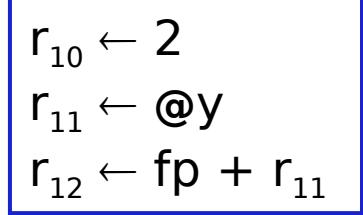
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r10 ← 2
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r19 ← @w
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```

# Steps of the Simplifier *(3-operation window)*

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## LLIR Code

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```



# Steps of the Simplifier

(3-operation window)

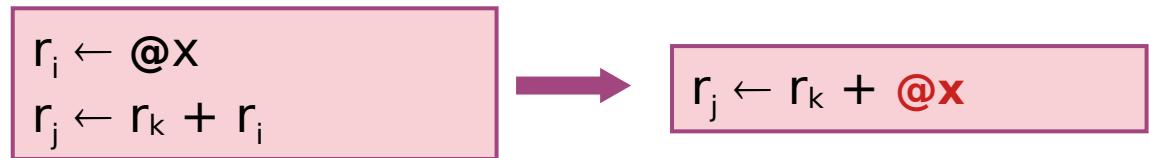
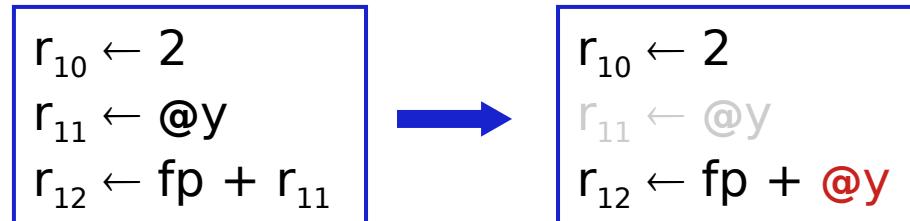
---

## LLIR Code

```

 $r_{10} \leftarrow 2$ 
 $r_{11} \leftarrow @y$ 
 $r_{12} \leftarrow fp + r_{11}$ 
 $r_{13} \leftarrow \mathbf{MEM}(r_{12})$ 
 $r_{14} \leftarrow r_{10} + r_{13}$ 
 $r_{15} \leftarrow @x$ 
 $r_{16} \leftarrow fp + r_{15}$ 
 $r_{17} \leftarrow \mathbf{MEM}(r_{16})$ 
 $r_{18} \leftarrow r_{14} - r_{17}$ 
 $r_{19} \leftarrow @w$ 
 $r_{20} \leftarrow fp + r_{19}$ 
 $\mathbf{MEM}(r_{20}) \leftarrow r_{18}$ 

```



$r_{10} \leftarrow 2$

$r_{11} \leftarrow @y$

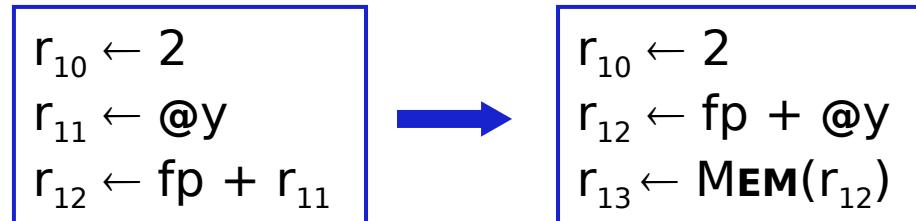
$r_{12} \leftarrow fp + @y$

$r_j \leftarrow r_k + @x$

# Steps of the Simplifier *(3-operation window)*

## LLIR Code

```
r10 ← 2  
r11 ← @y  
r12 ← fp + r11  
r13 ← MEM(r12)  
  
r14 ← r10 + r13  
r15 ← @x  
r16 ← fp + r15  
r17 ← MEM(r16)  
  
r18 ← r14 - r17  
r19 ← @w  
r20 ← fp + r19  
MEM(r20) ← r18
```



# Steps of the Simplifier

*(3-operation window)*

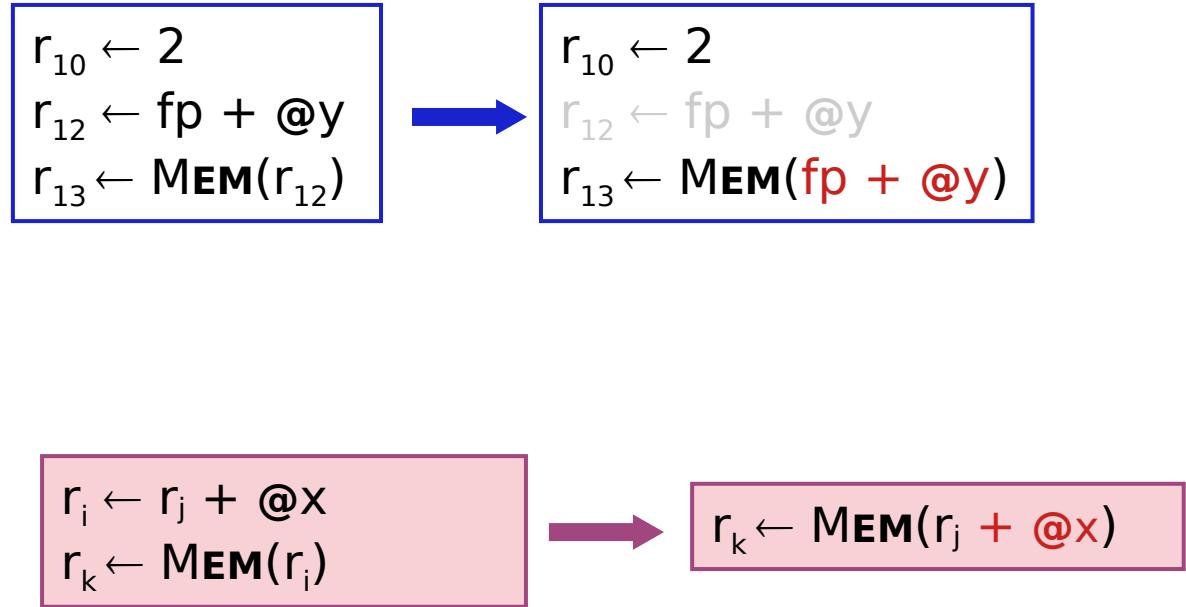
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## LLIR Code

```

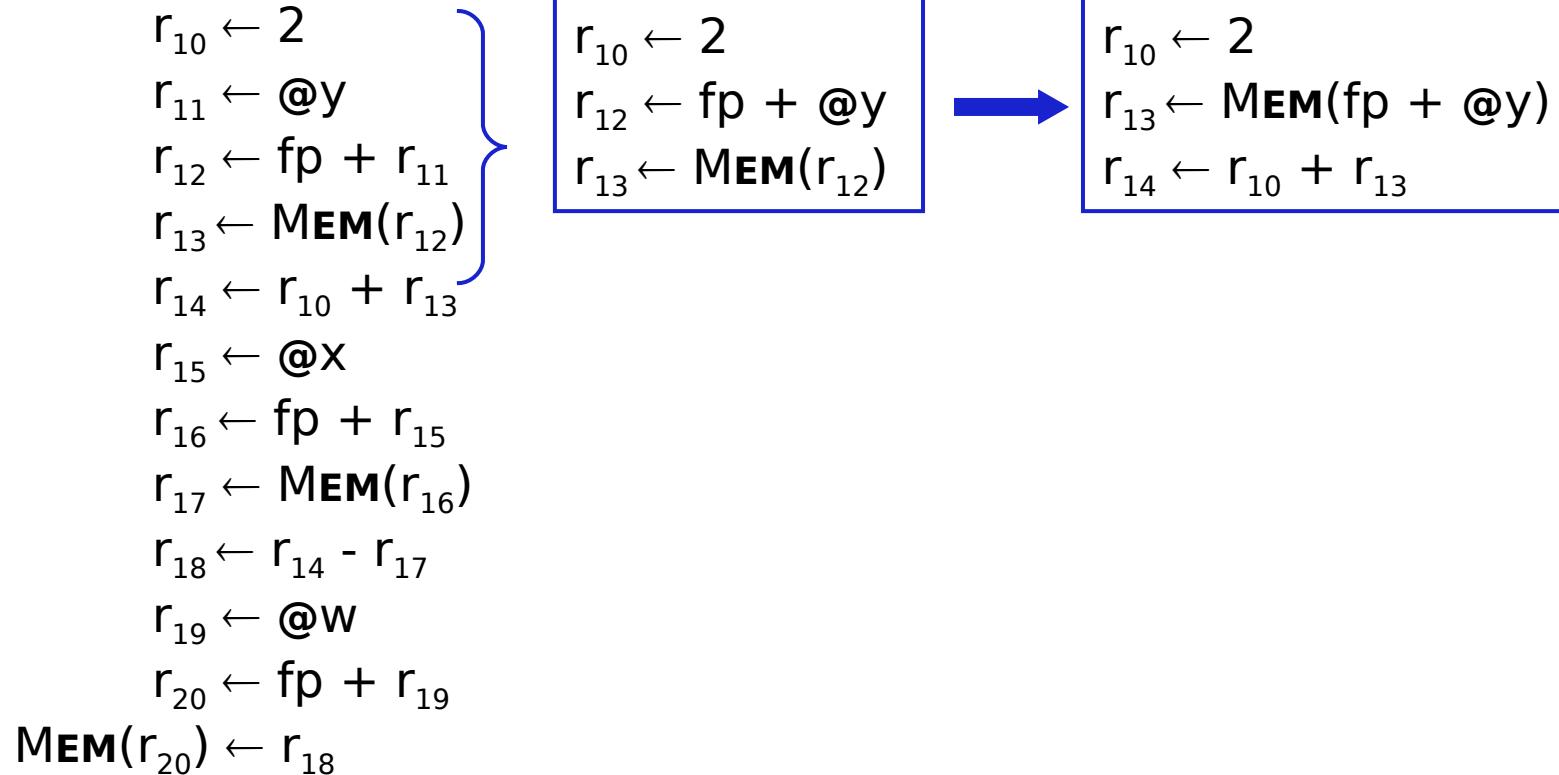
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 $r_{11} \leftarrow @y$ 
 $r_{12} \leftarrow fp + r_{11}$ 
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 $r_{19} \leftarrow @w$ 
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 $\mathbf{MEM}(r_{20}) \leftarrow r_{18}$ 

```



# Steps of the Simplifier *(3-operation window)*

## LLIR Code

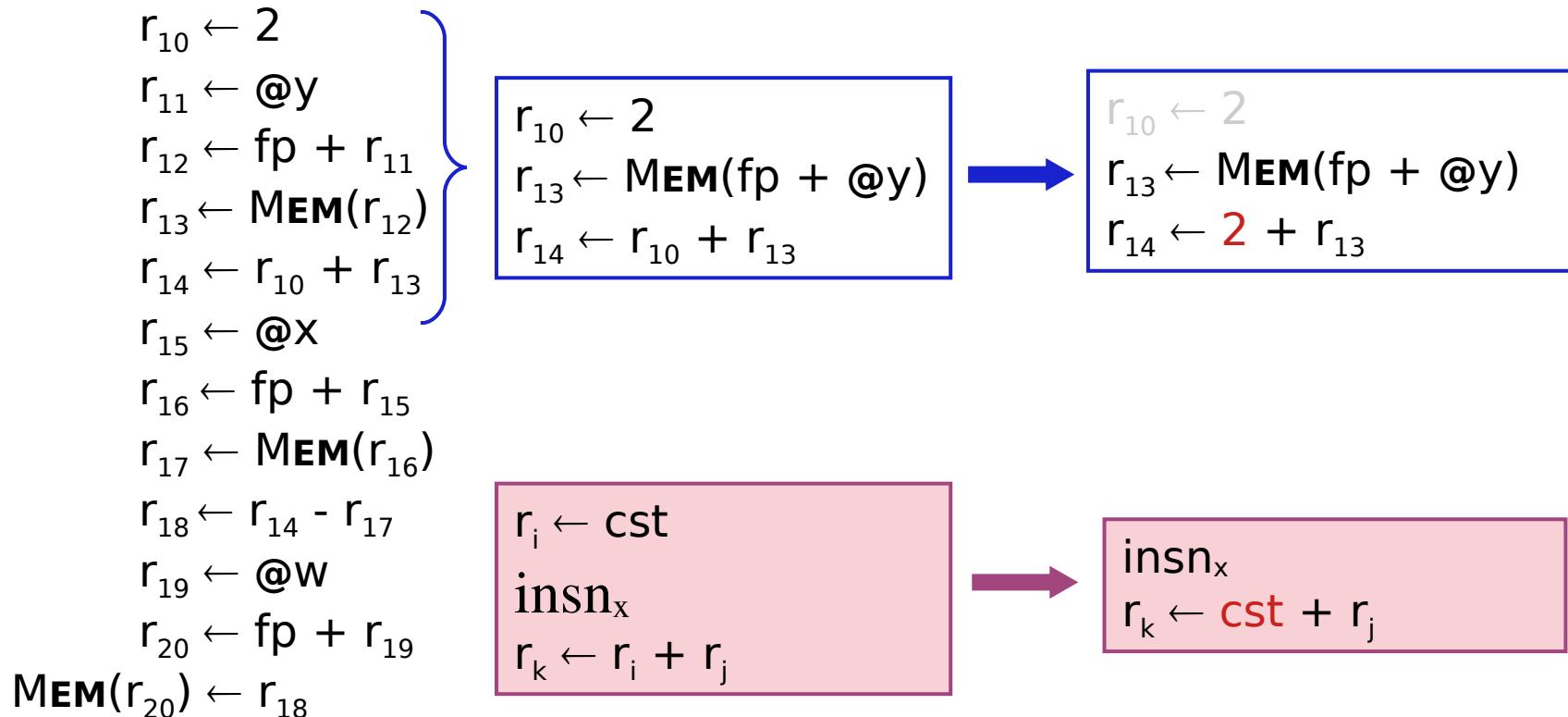


# Steps of the Simplifier

*(3-operation window)*

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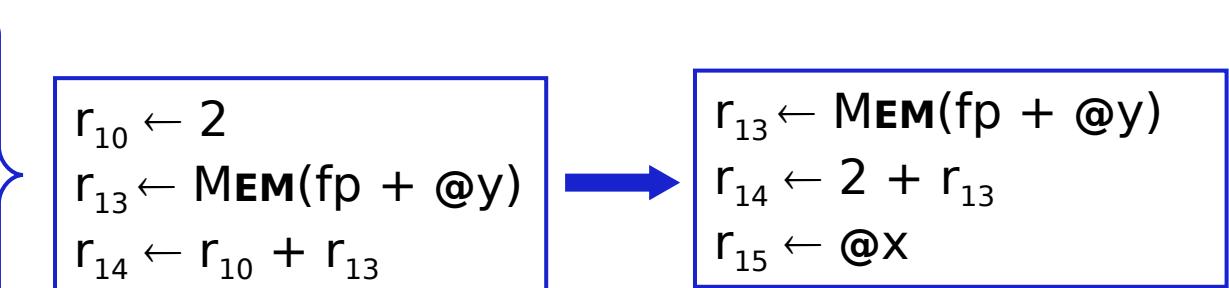
## LLIR Code



# Steps of the Simplifier *(3-operation window)*

## LLIR Code

$r_{10} \leftarrow 2$   
 $r_{11} \leftarrow @y$   
 $r_{12} \leftarrow fp + r_{11}$   
 $r_{13} \leftarrow \mathbf{MEM}(r_{12})$   
 $r_{14} \leftarrow r_{10} + r_{13}$   
 $r_{15} \leftarrow @x$   
 $r_{16} \leftarrow fp + r_{15}$   
 $r_{17} \leftarrow \mathbf{MEM}(r_{16})$   
 $r_{18} \leftarrow r_{14} - r_{17}$   
 $r_{19} \leftarrow @w$   
 $r_{20} \leftarrow fp + r_{19}$   
 $\mathbf{MEM}(r_{20}) \leftarrow r_{18}$



The diagram illustrates the simplification process. It starts with a brace grouping the first five assignments. This is followed by a blue box containing the next three assignments:  $r_{10} \leftarrow 2$ ,  $r_{13} \leftarrow \mathbf{MEM}(fp + @y)$ , and  $r_{14} \leftarrow r_{10} + r_{13}$ . A large blue arrow points from this box to another blue box containing the simplified form:  $r_{13} \leftarrow \mathbf{MEM}(fp + @y)$ ,  $r_{14} \leftarrow 2 + r_{13}$ , and  $r_{15} \leftarrow @x$ .

# Steps of the Simplifier

*(3-operation window)*

LLIR Code

```

 $r_{10} \leftarrow 2$ 
 $r_{11} \leftarrow @y$ 
 $r_{12} \leftarrow fp + r_{11}$ 
 $r_{13} \leftarrow \mathbf{MEM}(r_{12})$ 
 $r_{14} \leftarrow r_{10} + r_{13}$ 
 $r_{15} \leftarrow @x$ 
 $r_{16} \leftarrow fp + r_{15}$ 
 $r_{17} \leftarrow \mathbf{MEM}(r_{16})$ 
 $r_{18} \leftarrow r_{14} - r_{17}$ 
 $r_{19} \leftarrow @w$ 
 $r_{20} \leftarrow fp + r_{19}$ 
 $\mathbf{MEM}(r_{20}) \leftarrow r_{18}$ 

```

**1<sup>st</sup> op rolling out of window**

$r_{13} \leftarrow \mathbf{MEM}(fp + @y)$

$r_{13} \leftarrow \mathbf{MEM}(fp + @y)$   
 $r_{14} \leftarrow 2 + r_{13}$   
 $r_{15} \leftarrow @x$

$r_{14} \leftarrow 2 + r_{13}$   
 $r_{15} \leftarrow @x$   
 $r_{16} \leftarrow fp + r_{15}$

# Steps of the Simplifier

*(3-operation window)*

---

## LLIR Code

```

 $r_{10} \leftarrow 2$ 
 $r_{11} \leftarrow @y$ 
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 $r_{13} \leftarrow \mathbf{MEM}(r_{12})$ 
 $r_{14} \leftarrow r_{10} + r_{13}$ 
 $r_{15} \leftarrow @x$ 
 $r_{16} \leftarrow fp + r_{15}$ 
 $r_{17} \leftarrow \mathbf{MEM}(r_{16})$ 
 $r_{18} \leftarrow r_{14} - r_{17}$ 
 $r_{19} \leftarrow @w$ 
 $r_{20} \leftarrow fp + r_{19}$ 
 $\mathbf{MEM}(r_{20}) \leftarrow r_{18}$ 

```

$r_{14} \leftarrow 2 + r_{13}$   
 $r_{15} \leftarrow @x$   
 $r_{16} \leftarrow fp + r_{15}$



$r_{14} \leftarrow 2 + r_{13}$   
 $r_{15} \leftarrow @x$   
 $r_{16} \leftarrow fp + @x$

$r_i \leftarrow @x$   
 $r_j \leftarrow r_k + r_i$



$r_j \leftarrow r_k + @x$

# Steps of the Simplifier *(3-operation window)*

## LLIR Code

```
r10 ← 2  
r11 ← @y  
r12 ← fp + r11  
r13 ← MEM(r12)  
r14 ← r10 + r13  
r15 ← @x  
r16 ← fp + r15  
r17 ← MEM(r16)  
r18 ← r14 - r17  
r19 ← @w  
r20 ← fp + r19  
MEM(r20) ← r18
```

$r_{14} \leftarrow 2 + r_{13}$   
 $r_{15} \leftarrow @x$   
 $r_{16} \leftarrow fp + r_{15}$



$r_{14} \leftarrow 2 + r_{13}$   
 $r_{16} \leftarrow fp + @x$   
 $r_{17} \leftarrow **MEM**(r<sub>16</sub>)$

# Steps of the Simplifier (3-operation window)

## LLIR Code

```

 $r_{10} \leftarrow 2$ 
 $r_{11} \leftarrow @y$ 
 $r_{12} \leftarrow fp + r_{11}$ 
 $r_{13} \leftarrow \mathbf{MEM}(r_{12})$ 
 $r_{14} \leftarrow r_{10} + r_{13}$ 
 $r_{15} \leftarrow @x$ 
 $r_{16} \leftarrow fp + r_{15}$ 
 $r_{17} \leftarrow \mathbf{MEM}(r_{16})$ 
 $r_{18} \leftarrow r_{14} - r_{17}$ 
 $r_{19} \leftarrow @w$ 
 $r_{20} \leftarrow fp + r_{19}$ 
 $\mathbf{MEM}(r_{20}) \leftarrow r_{18}$ 

```

$r_{14} \leftarrow 2 + r_{13}$   
 $r_{16} \leftarrow fp + @x$   
 $r_{17} \leftarrow \mathbf{MEM}(r_{16})$

$r_{14} \leftarrow 2 + r_{13}$   
 $r_{17} \leftarrow \mathbf{MEM}(fp + @x)$   
 $r_{18} \leftarrow r_{14} - r_{17}$

$r_i \leftarrow r_j + @x$   
 $r_k \leftarrow \mathbf{MEM}(r_i)$

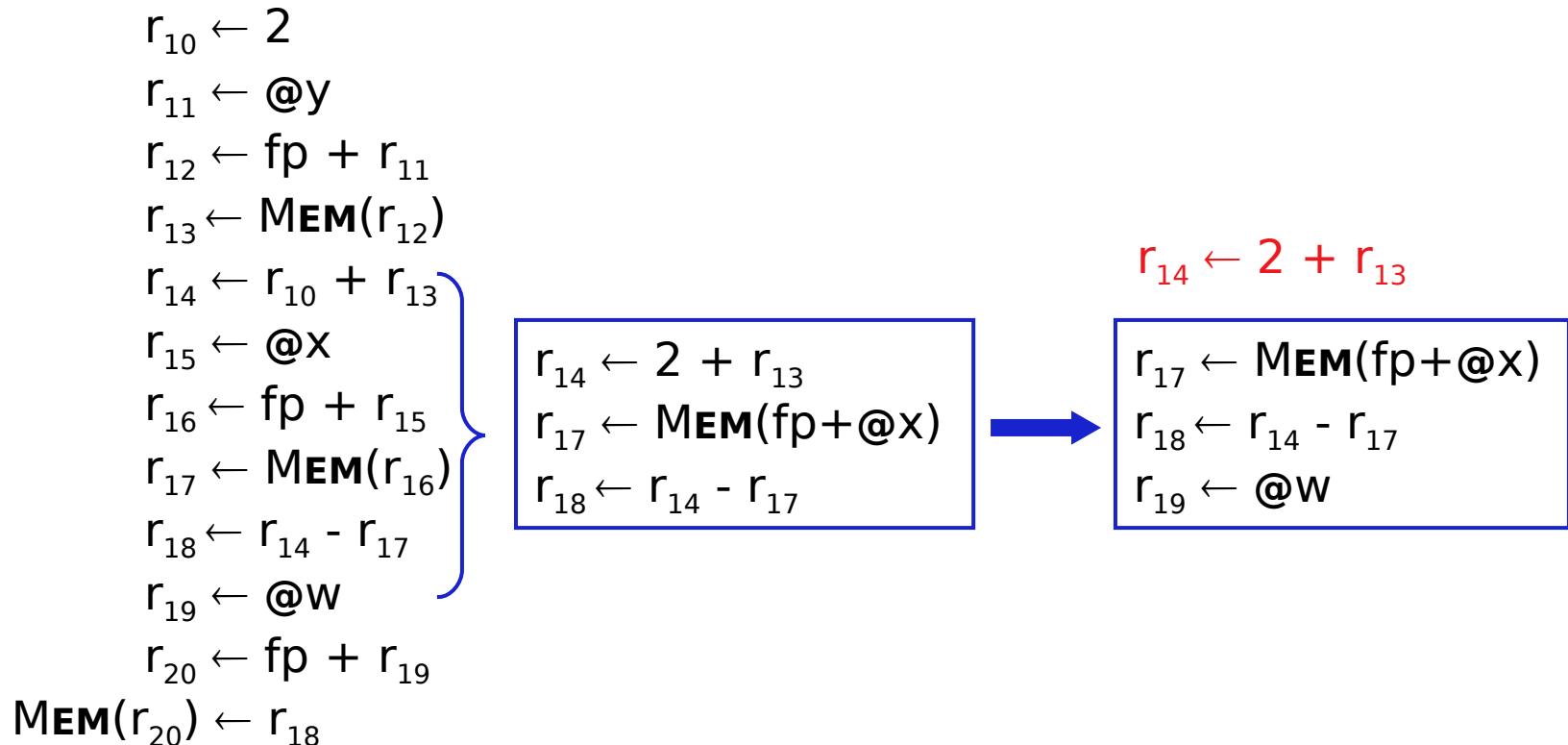
$r_k \leftarrow \mathbf{MEM}(r_j + @x)$

# Steps of the Simplifier

*(3-operation window)*

---

## LLIR Code



# Steps of the Simplifier *(3-operation window)*

## LLIR Code

$r_{10} \leftarrow 2$   
 $r_{11} \leftarrow @y$   
 $r_{12} \leftarrow fp + r_{11}$   
 $r_{13} \leftarrow \mathbf{MEM}(r_{12})$   
 $r_{14} \leftarrow r_{10} + r_{13}$   
 $r_{15} \leftarrow @x$   
 $r_{16} \leftarrow fp + r_{15}$   
 $r_{17} \leftarrow \mathbf{MEM}(r_{16})$   
 $r_{18} \leftarrow r_{14} - r_{17}$   
 $r_{19} \leftarrow @w$   
 $r_{20} \leftarrow fp + r_{19}$   
 $\mathbf{MEM}(r_{20}) \leftarrow r_{18}$

$r_{17} \leftarrow \mathbf{MEM}(fp+@x)$

$r_{17} \leftarrow \mathbf{MEM}(fp+@x)$   
 $r_{18} \leftarrow r_{14} - r_{17}$   
 $r_{19} \leftarrow @w$

$r_{18} \leftarrow r_{14} - r_{17}$   
 $r_{19} \leftarrow @w$   
 $r_{20} \leftarrow fp + r_{19}$   
 $r_{17} \leftarrow \mathbf{MEM}(fp+@x)$

# Steps of the Simplifier

(3-operation window)

## LLIR Code

```
r10 ← 2  
r11 ← @y  
r12 ← fp + r11  
r13 ← MEM(r12)  
r14 ← r10 + r13  
r15 ← @x  
r16 ← fp + r15  
r17 ← MEM(r16)  
r18 ← r14 - r17  
r19 ← @w  
r20 ← fp + r19  
MEM(r20) ← r18
```

r<sub>i</sub> ← @x  
r<sub>j</sub> ← r<sub>k</sub> + r<sub>i</sub>



r<sub>j</sub> ← r<sub>k</sub> + @x

r<sub>18</sub> ← r<sub>14</sub> - r<sub>17</sub>  
r<sub>19</sub> ← @w  
r<sub>20</sub> ← fp + r<sub>19</sub>



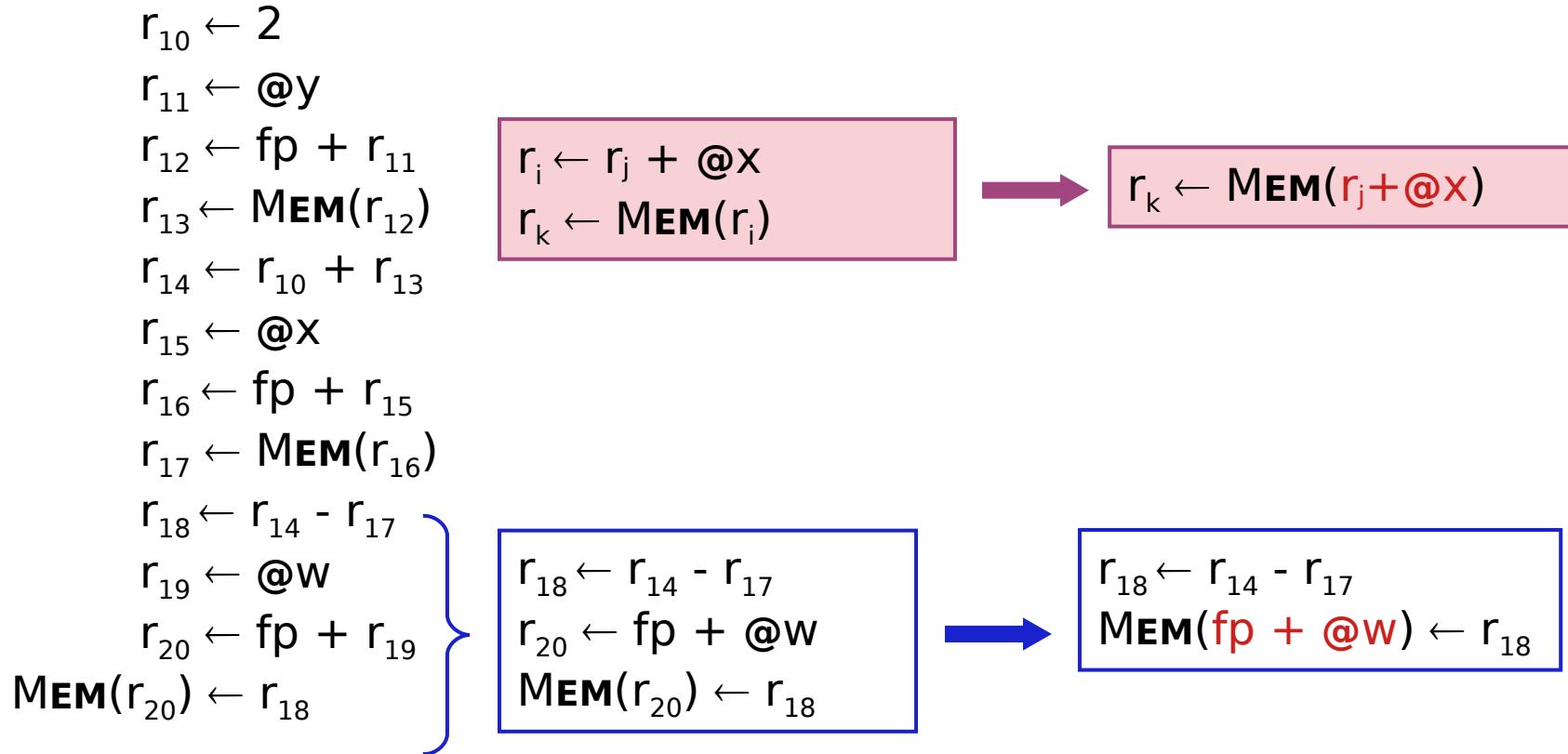
r<sub>18</sub> ← r<sub>14</sub> - r<sub>17</sub>  
r<sub>20</sub> ← fp + @w  
**MEM**(r<sub>20</sub>) ← r<sub>18</sub>

# Steps of the Simplifier

*(3-operation window)*

---

## LLIR Code

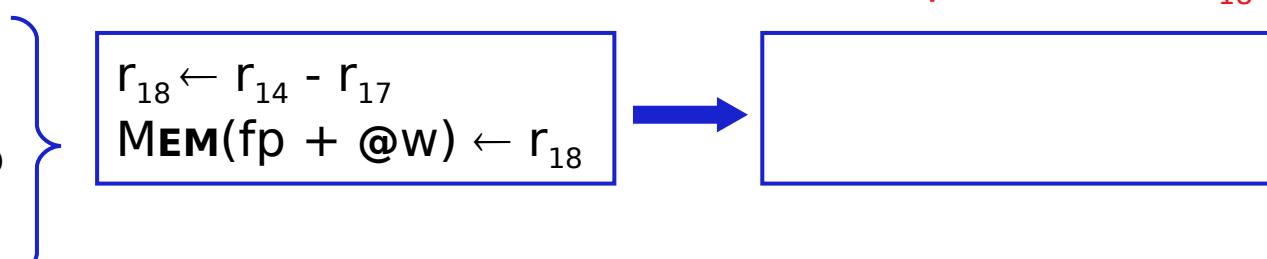


# Steps of the Simplifier *(3-operation window)*

## LLIR Code

```
r10 ← 2
r11 ← @y
r12 ← fp + r11
r13 ← MEM(r12)
r14 ← r10 + r13
r15 ← @x
r16 ← fp + r15
r17 ← MEM(r16)
r18 ← r14 - r17
r19 ← @w
r20 ← fp + r19
MEM(r20) ← r18
```

$r_{18} \leftarrow r_{14} - r_{17}$   
 $\mathbf{MEM}(\text{fp} + @w) \leftarrow r_{18}$



# Example

---

## LLIR Code

```
r10 ← 2  
r11 ← @y  
r12 ← fp + r11  
r13 ← MEM(r12)  
r14 ← r10 + r13  
r15 ← @x  
r16 ← fp + r15  
r17 ← MEM(r16)  
r18 ← r14 - r17  
r19 ← @w  
r20 ← fp + r19  
MEM(r20) ← r18
```

Simplify



LLIR Code

```
r13 ← MEM(fp + @y)  
r14 ← 2 + r13  
r17 ← MEM(fp + @x)  
r18 ← r14 - r17  
MEM(fp + @w) ← r18
```

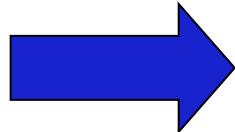
# Example: Matcher

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## LLIR Code

```
r13 ← MEM(fp + @y)  
r14 ← 2 + r13  
r17 ← MEM(fp + @x)  
r18 ← r14 - r17  
MEM(fp + @w)← r18
```

Match



## MIPS Code

LW	r <sub>13</sub> , @y(\$fp)
ADDI	r <sub>14</sub> , r <sub>13</sub> , 2
LW	r <sub>17</sub> , @x(\$fp)
SUB	r <sub>18</sub> , r <sub>14</sub> , r <sub>17</sub>
SW	r <sub>18</sub> , @w(\$fp)

- Produces pretty good code

# Making It All Work

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## Details

- LLIR is largely machine independent (RTL)
- Target machine described as LLIR → ASM pattern
- Actual pattern matching
  - Use a hand-coded pattern matcher (GCC)
  - Turn patterns into grammar & use LR parser (VPO)
- Several important compilers use this technology
- It seems to produce good portable instruction selectors

Key strength appears to be late low-level optimization

# Next lecture

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## Instruction selection

- Tree-based pattern matching (LLVM)