Symbol tables

Symbol tables are used to describe and analyse definitions and uses of identifiers. Grammars are too weak; the language:

\[ \{ waw | w \in \Sigma^* \} \]

is not context-free.

A symbol table is a map from identifiers to meanings:

<table>
<thead>
<tr>
<th>i</th>
<th>local</th>
<th>int</th>
</tr>
</thead>
<tbody>
<tr>
<td>done</td>
<td>local</td>
<td>boolean</td>
</tr>
<tr>
<td>insert</td>
<td>method</td>
<td>...</td>
</tr>
<tr>
<td>List</td>
<td>class</td>
<td>...</td>
</tr>
<tr>
<td>x</td>
<td>formal</td>
<td>list</td>
</tr>
<tr>
<td>:</td>
<td>:</td>
<td>:</td>
</tr>
</tbody>
</table>

We must construct a symbol table for every program point.

Using symbol tables to analyse JOOS:
- which classes are defined;
- what is the inheritance hierarchy;
- is the hierarchy well-formed;
- which fields are defined;
- which methods are defined;
- what are the signatures of methods;
- are identifiers defined twice;
- are identifiers defined when used; and
- are identifiers used properly?

Static, nested scope rules:

The standard of modern languages.
Old-style one-pass technology:

Still haunts some languages:

```c
void weedPROGRAM(PROGRAM *p);
void weedCLASSFILE(CLASSFILE *c);
void weedCLASS(CLASS *c);
```

Forward declarations enable recursion.

Use the most closely nested definition:

Identifiers at same level must be unique.

The symbol table behaves like a stack:

The symbol table can be implemented as a simple stack:

- `pushSymbol(SymbolTable *t, char *name, ...)`
- `popSymbol(SymbolTable *t)`
- `getSymbol(SymbolTable *t, char *name)`

But how do we detect multiple definitions of an identifier at the same level?

Use bookmarks and a cactus stack:

- `scopeSymbolTable(SymbolTable *t)`
- `putSymbol(SymbolTable *t, char *name, ...)`
- `unscopeSymbolTable(SymbolTable *t)`
- `getSymbol(SymbolTable *t, char *name)`

Still just linear search, though.
Implement symbol tables as a cactus stack of hash tables:

- each hash table contains the identifiers in a level;
- push a new hash table when a level is entered;
- each identifier is entered in the top hash table;
- it is an error if it is already there;
- a use of an identifier is looked up in the hash tables from top to bottom;
- it is an error if it is not found;
- pop a hash table when a level is left.

What is a good hash function on identifiers?
Use the initial letter:
- \texttt{codePROGRAM}, \texttt{codeMETHOD}, \texttt{codeEXP}, \ldots
Use the sum of the letters:
- doesn’t distinguish letter order
Use the shifted sum of the letters:

\[
\begin{align*}
\text{"j"} &= \text{106} = 0000000001101010 \\
\text{shift} &= 0000000011010100 \\
+ \text{"o"} &= \text{111} = 0000000001101111 \\
= &= 0000001011101101 \\
\text{shift} &= 0000010111101010 \\
+ \text{"o"} &= \text{111} = 0000000001101111 \\
= &= 0000011001011101 \\
= &= 1629
\end{align*}
\]

Hash tables for the JOOS source code:

\[
\begin{align*}
\text{hash} &= \star \text{str} \\
\text{while} (\star \text{str}) \text{ hash} &= \text{hash} + \star \text{str}++; \\
\text{while} (\star \text{str}) \text{ hash} &= (\text{hash} \ll 1) + \star \text{str}++;
\end{align*}
\]
How to handle mutual recursion:

A

...B...

B

...A...

A single traversal of the abstract syntax tree is not enough.

Make two traversals:
- collect definitions of identifiers; and
- analyse uses of identifiers.

For cases like recursive types, the definition is not completed before the second traversal.

Symbol information in JOOS:

$ cat tree.h

[...]

typedef enum{classSym,fieldSym,methodSym,
formalSym,localSym} SymbolKind;

typedef struct SYMBOL {
    char *name;
    SymbolKind kind;
    union {
        struct CLASS *classS;
        struct FIELD *fieldS;
        struct METHOD *methodS;
        struct FORMAL *formalS;
        struct LOCAL *localS;
    } val;
} SYMBOL;

[...]
Symbol tables are weaved together with abstract syntax trees:

```java
public class B extends A {
    protected A a;
    protected B b;
    public void m(A x, B y) {
        this.m(a, b);
    }
}
```

Complicated recursion in JOOS is resolved through multiple passes:

```bash
$ cat symbol.c
[...]
void symPROGRAM(PROGRAM *p)
{ classlib = initSymbolTable();
  symInterfacePROGRAM(p, classlib);
  symInterfaceTypesPROGRAM(p, classlib);
  symImplementationPROGRAM(p);
}
[...]
```

Each pass goes into further detail:

- **symInterfacePROGRAM**: define classes and their interfaces;
- **symInterfaceTypesPROGRAM**: build hierarchy and analyse interface types; and
- **symImplementationPROGRAM**: define locals and analyse method bodies.

### Defining a JOOS class:

```c
void symInterfaceCLASS(CLASS *c, SymbolTable *sym) {
    SYMBOL *s;
    if (defSymbol(sym, c->name)) {
        reportStrError("class name %s already defined",
                       c->name, c->lineno);
    } else {
        s = putSymbol(sym, c->name, classSym);
        s->val.classS = c;
        c->localsym = initSymbolTable();
        symInterfaceFIELD(c->fields, c->localsym);
        symInterfaceCONSTRUCTOR(c->constructors, c->localsym);
        symInterfaceMETHOD(c->methods, c->localsym);
    }
}
```

### Defining a JOOS method:

```c
void symInterfaceMETHOD(METHOD *m, SymbolTable *sym) {
    SYMBOL *s;
    if (m != NULL) {
        symInterfaceMETHOD(m->next, sym);
        if (defSymbol(sym, m->name)) {
            reportStrError("method name %s already defined",
                           m->name, m->lineno);
        } else {
            s = putSymbol(sym, m->name, methodSym);
            s->val.methodS = m;
        }
    }
}
```

and its signature:

```c
void symInterfaceTypesMETHOD(METHOD *m, SymbolTable *sym) {
    if (m != NULL) {
        symInterfaceTypesMETHOD(m->next, sym);
        symTYPE(m->returntype, sym);
        symInterfaceTypesFORMAL(m->formals, sym);
    }
}
```
Analysing a JOOS class implementation:

```c
void symImplementationCLASS(CLASS *c)
{
    SymbolTable *sym;
    sym = scopeSymbolTable(classlib);
    symImplementationFIELD(c->fields,sym);
    symImplementationCONSTRUCTOR(c->constructors,c,sym);
    symImplementationMETHOD(c->methods,c,sym);
}
```

Analysing a JOOS method implementation:

```c
void symImplementationMETHOD(METHOD *m, CLASS *this, SymbolTable *sym)
{
    SymbolTable *msym;
    if (m!=NULL) {
        symImplementationMETHOD(m->next,this,sym);
        msym = scopeSymbolTable(sym);
        symImplementationFORMAL(m->formals,msym);
        symImplementationSTATEMENT(m->statements,this,msym,
                                    m->modifier==staticMod);
    }
}
```

Analysing JOOS statements:

```c
void symImplementationSTATEMENT(STATEMENT *s, CLASS *this, SymbolTable *sym, int stat)
{
    SymbolTable *ssym;
    if (s!=NULL) {
        switch (s->kind) {
        ...
        case localK:
            symImplementationLOCAL(s->val.localS,sym);
            break;
        ...
        case blockK:
            ssym = scopeSymbolTable(sym);
            symImplementationSTATEMENT(s->val.blockS.body,
                                        this,ssym,stat);
            break;
        ...
        }
    }
}
```

Analysing JOOS local declarations:

```c
void symImplementationLOCAL(LOCAL *l, SymbolTable *sym)
{
    SYMBOL *s;
    if (l!=NULL) {
        symImplementationLOCAL(l->next,sym);
        symTYPE(l->type,sym);
        if (defSymbol(sym,l->name)) {
            reportStrError("local %s already declared",
                            l->name,l->lineno);
        } else {
            s = putSymbol(sym,l->name,localSym);
            s->val.localS = l;
        }
    }
}
```

Identifier lookup in the JOOS class hierarchy:

```c
SYMBOL *lookupHierarchy(char *name, CLASS *start)
{
    SYMBOL *s;
    if (start==NULL) return NULL;
    s = getSymbol(start->localsym,name);
    if (s!=NULL) return s;
    if (start->parent==NULL) return NULL;
    return lookupHierarchy(name,start->parent);
}
```

```c
CLASS *lookupHierarchyClass(char *name, CLASS *start)
{
    SYMBOL *s;
    if (start==NULL) return NULL;
    s = getSymbol(start->localsym,name);
    if (s!=NULL) return s;
    if (start->parent==NULL) return NULL;
    return lookupHierarchyClass(name,start->parent);
}
```

For which class do we return NULL on line 5 of each function?
Analysing expressions:

```c
void symImplementationEXP(EXP *e, CLASS *this,
    SymbolTable *sym, int stat)
{
    switch (e->kind) {
        case idK:
            e->val.idE.idsym = symVar(e->val.idE.name,sym,
                                      this,e->lineno,stat);
            break;
        case assignK:
            e->val.assignE.leftsym =
                symVar(e->val.assignE.left,sym,
                        this,e->lineno,stat);
            symImplementationEXP(e->val.assignE.right,
                                  this,sym,stat);
            break;
        [...]
    }
}
```

Analysing an identifier:

```c
SYMBOL *symVar(char *name, SymbolTable *sym,
                 CLASS *this, int lineno, int stat)
{
    SYMBOL *s;
    s = getSymbol(sym,name);
    if (s==NULL) {
        s = lookupHierarchy(name,this);
        if (s==NULL) {
            reportStrError("identifier %s not declared",
                            name,lineno);
        } else {
            if ((s->kind!=fieldSym) && (s->kind!=formalSym) &&
                (s->kind!=localSym))
                reportStrError("%s is not a variable as expected",
                               name,lineno);
            } else {
                if ((s->kind!=fieldSym) && (s->kind!=formalSym) &&
                    (s->kind!=localSym))
                    reportStrError("%s is not a variable as expected",
                                   name,lineno);
                } else {
                    if (s->kind==fieldSym && s->kind==formalSym &&
                        (s->kind==localSym))
                        reportStrError("%s is not a variable as expected",
                                       name,lineno);
                }
        } else {
            if ((s->kind!=fieldSym) && s->kind==formalSym &&
                (s->kind==localSym))
                reportStrError("illegal static reference to %s",
                               name,lineno);
            return s;
        }
    }
```

The testing strategy for the symbol tables involves an extension of the pretty printer.

A textual representation of the symbol table is printed once for every scope area.

- In Java, use `toString()`.

These tables are then compared to a corresponding manual construction for a sufficient collection of programs.

Furthermore, every error message should be provoked by some test program.