

# COMP 364: Computer Tools for Life Sciences

Algorithm design: Selection and Insertion Sort

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# Key course information

## Quiz #4 is available today!

- ▶ available on MyCourses (multiple choice questions)
- ▶ Quiz #4 closes at 11:59:59 pm on Monday, October 16th
- ▶ questions cover topics from the last two weeks

## Midterm review sessions

- ▶ 1:00-2:00 pm on Thursday, October 19th
  - ▶ hosted by Roman (TR 3104)
- ▶ 5:00-7:00 pm on Friday, October 20th
  - ▶ hosted by Pouriya (TR 3104)
- ▶ TAs will not prepare material, be sure to arrive with questions

# Sorting algorithms

A **sorting algorithm** is an algorithm that takes

- ▶ a list/array as input
- ▶ performs specified operations on the list/array
- ▶ outputs a sorted list/array

For example:

- ▶  $[a, c, d, b]$  could be sorted alphabetically to  $[a, b, c, d]$
- ▶  $[1, 3, 2, 0]$  could be sorted:
  - ▶ increasing order:  $[0, 1, 2, 3]$
  - ▶ or decreasing order:  $[3, 2, 1, 0]$

# Why is it useful to sort data?

Sorted data searching can be optimized to a very high level

- ▶ also used to represent data in more readable formats

## Contacts

- ▶ your mobile phone stores the telephone numbers of contacts by names
- ▶ names can easily be searched to find a desired number

## Dictionary

- ▶ dictionaries store words in alphabetical order to allow for easy searching of any word

Remember **binary search**?

# Adding more algorithms to your toolbox

In the last lecture, we covered searching algorithms, specifically:

- ▶ linear search
- ▶ binary search

Today, we will cover the following sorting algorithms:

- ▶ selection sort
- ▶ insertion sort

Images are taken from the following online tutorial: [https://www.tutorialspoint.com/data\\_structures\\_algorithms/](https://www.tutorialspoint.com/data_structures_algorithms/)

# Selection sort

Conceptually the most simple of all the sorting algorithms

Start by selecting the smallest (or largest) item in a list

- ▶ then place this item at the start of the list
- ▶ repeat for the remaining items in the list
  - ▶ move next smallest/largest item to the second position
  - ▶ then the next
  - ▶ and so on and so on...
  - ▶ until the list is sorted

Let's consider the following unsorted list:



## Selection sort #2

For the first position in the resulting sorted list

- ▶ the whole list is scanned sequentially
- ▶ the first position is where 14 is currently stored

We search the whole list

- ▶ to find that 10 is the lowest value in the list



## Selection sort #3

We then replace 14 with 10



After one iteration

- ▶ 10, which happens to be the minimum value in the list
- ▶ appears in the first position of the sorted list

For the second position

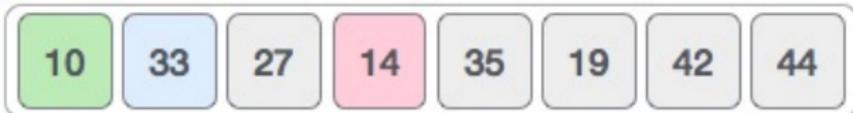
- ▶ where 33 is residing
- ▶ we start scanning the rest of the list in a linear manner



## Selection sort #4

14 is found to be the second lowest value in the list

- ▶ and should appear at the second place
- ▶ we swap these values.



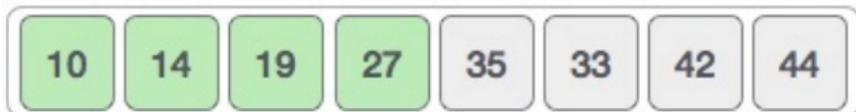
After two iterations

- ▶ the two items with the least values
- ▶ are positioned at the beginning in a sorted manner



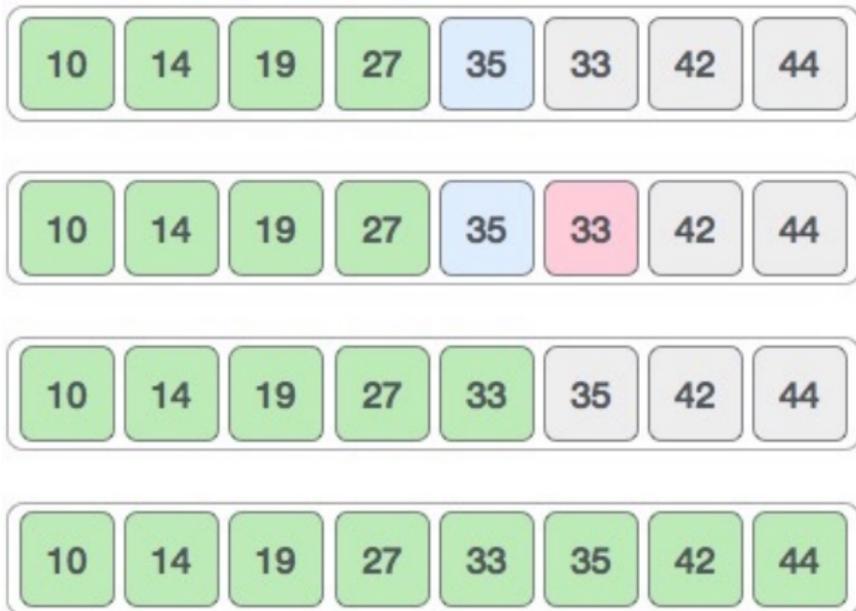
## Selection sort #5

The same process is applied to the rest of the items in the list



## Selection sort #6

Until the list is sorted



# Selection sort algorithm

## Selection sort (*sequence*)

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- Step 1 - find the item with the smallest value in *sequence*
- Step 2 - swap it with the first item in *sequence*
- Step 3 - find the item with the second smallest value in *sequence*
- Step 4 - swap it with the second item in *sequence*
- Step 5 - find the item with the third smallest value in *sequence*
- Step 6 - swap it with the third item in *sequence*
- Step 7 - repeat finding the item with the next smallest value
- Step 8 - then swap it with the correct item until *sequence* is sorted

# Selection sort: pseudocode

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**Algorithm 1** Selection sort

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```
1: procedure SELECTION_SORT(sequence)
2:    $N \leftarrow$  length of sequence
3:   for  $i \leftarrow 0$  to  $N - 1$  do
4:      $min \leftarrow i$ 
5:     for  $j \leftarrow i + 1$  to  $N - 1$  do
6:       if  $sequence[j] \leq sequence[min]$  then
7:          $min \leftarrow j$ 
8:       end if
9:     end for
10:    SWAP( $sequence[i], sequence[min]$ )
11:  end for
12: end procedure
```

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## Selection sort: Python implementation

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```
1 import random
2
3 sequence = list(range(0,10))
4 random.shuffle(sequence) # shuffles items
5 N = len(sequence)
6 for i in range(0,N,1): # why not N-1?
7     min = i
8     for j in range(i+1,N,1):
9         if sequence[j] <= sequence[min]:
10            min = j
11     sequence[i],sequence[min] = sequence[min],sequence[i]
12 print(sequence) # prints ???
```

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# Insertion sort

Insertion sort does what you might expect

- ▶ inserts each item of the list into its proper position
- ▶ resulting in progressively larger sequences of a sorted list

Start with a sorted list of 1 element on the left and  $N-1$  unsorted items on the right

- ▶ take the first unsorted item
- ▶ insert it into the sorted list, moving elements as necessary
- ▶ now have a sorted list of size 2, and  $N - 2$  unsorted elements
- ▶ repeat for all items

## Insertion sort #2

Let's reuse our unsorted list from before and sort it in ascending order:



Start by comparing the first two items



## Insertion sort #3

We find that both 14 and 33 are already in ascending order.

- ▶ for now, 14 is in sorted sub-list



Insertion sort then moves ahead and compares 33 with 27



## Insertion sort #4

33 is not in the correct position



Swap 33 with 27

- ▶ also check that all the elements of sorted sub-list
- ▶ we see that the sorted sub-list has only one element 14
- ▶ 27 is greater than 14
- ▶ therefore, the sorted sub-list remains sorted after swapping



## Insertion sort #5

Now that we have 14 and 27 in the sorted sub-list

- ▶ compare 33 with 10



Values are not in a sorted order



So we swap them



## Insertion sort #6

However, swapping makes 27 and 10 unsorted



We swap them too



We find 14 and 10 in an unsorted order



## Insertion sort #7

We swap them again

- ▶ by the end of third iteration, we have a sorted sub-list of 4 items



This process goes on until all the unsorted values are covered in a sorted sub-list

# Insertion sort algorithm

## Insertion sort (*sequence*)

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Step 1 - If it is the first element, item is already sorted

Step 2 - select next item

Step 3 - compare against all other items in the sorted sub-list

Step 4 - shift all the elements in the sorted sub-list that are greater than the value to be sorted

Step 5 - Insert the value in the sorted sub-list

Step 6 - repeat until list is sorted

## Insertion sort: pseudocode

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### Algorithm 2 Insertion sort

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```
1: procedure INSERTION_SORT(sequence)
2:   for  $i \leftarrow 1$  to  $N - 1$  do
3:      $key \leftarrow sequence[i]$ 
4:     // inset key into the sorted sub-list
5:      $j \leftarrow i - 1$ 
6:     while  $j \geq 0$  and  $sequence[j] > key$  do
7:        $sequence[j + 1] \leftarrow sequence[j]$ 
8:        $j \leftarrow j - 1$ 
9:     end while
10:     $sequence[j + 1] \leftarrow key$ 
11:  end for
12: end procedure
```

---

# Insertion sort: Python implementation

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```
1 import random
2
3 sequence = list(range(0,10))
4 random.shuffle(sequence)
5 N = len(sequence)
6 for i in range(1,N,1):
7     j = i-1
8     key = sequence[i]
9     while(j >= 0 and sequence[j] > key):
10         sequence[j+1] = sequence[j]
11         j -= 1
12     sequence[j+1] = key
13 print(sequence)           # prints ???
```

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# Summary

Why learn both selection and insertion sort?

- ▶ insertion sort is expected to be faster
- ▶ selection sort makes more comparisons than movements
  - ▶ insertion sort is the opposite
- ▶ if less movement is needed
  - ▶ e.g., list is almost sorted
  - ▶ then selection sort is the better choice

**Question:** based on the algorithms you have already learned, how could you further improve insertion sort?