

Assignment 1 – COMP 523

Language-based security

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Exercise 1(20pts): In the lecture and in Pierce’s book, we define the following operational semantics for a small language of terms. For convenience, we repeat the evaluation rules here.

$$\begin{array}{c}
 \frac{}{\text{pred } (\text{succ } nv) \rightarrow nv} \text{ E-PRED-SUCC} \quad \frac{}{\text{pred } z \rightarrow z} \text{ E-PRED-ZERO} \\
 \\
 \frac{t \rightarrow t'}{\text{succ } t \rightarrow \text{succ } t'} \text{ E-SUCC} \quad \frac{t \rightarrow t'}{\text{pred } t \rightarrow \text{pred } t'} \text{ E-PRED} \\
 \\
 \frac{}{\text{if true then } t_1 \text{ else } t_2 \rightarrow t_1} \text{ E-IF-TRUE} \\
 \\
 \frac{}{\text{if false then } t_1 \text{ else } t_2 \rightarrow t_2} \text{ E-IF-FALSE} \\
 \\
 \frac{t \rightarrow t'}{\text{if } t \text{ then } t_1 \text{ else } t_2 \rightarrow \text{if } t' \text{ then } t_1 \text{ else } t_2} \text{ E-IF} \\
 \\
 \frac{}{\text{iszero } z \rightarrow \text{true}} \text{ E-ISZERO-ZERO} \\
 \\
 \frac{}{\text{iszero } (\text{succ } nv) \rightarrow \text{false}} \text{ E-ISZERO-SUCC} \\
 \\
 \frac{t \rightarrow t'}{\text{iszero } t \rightarrow \text{iszero } t'} \text{ E-ISZERO}
 \end{array}$$

A friend of yours suggests to replace the evaluation rule E-PRED-SUCC with the rule

$$\frac{}{\text{succ } (\text{pred } (nv)) \rightarrow nv}$$

Is this a good idea? What would you say to her or him? Which basic property discussed in Ch 3 breaks down? – If you think the above rule is good, verify that all the theorems in Ch 3 still hold. If you think the rule is bad, then give a counterexample and explain which theorem does not hold.

Exercise 2(30pts): Show that for the small-step semantics, we have that all values evaluate to themselves.

If v is a value and $v \rightarrow^* v'$ then $v = v'$.

Exercise 3(50pts) : An alternative style to the small-step semantics seen in class is the *big-step* semantics. The judgment $e \Downarrow v$ describes the complete evaluation of the expression e to some final value v . We concentrate here on the fragment for natural numbers. The rules for big-step evaluation for the small fragment consisting of **z**, **succ** e , **pred** e , and **iszero** e are given below.

$$\begin{array}{c}
\frac{}{\mathbf{z} \Downarrow \mathbf{z}} \text{ B-Z} \qquad \frac{e \Downarrow v}{\mathbf{succ} \ e \Downarrow \mathbf{succ} \ v} \text{ B-SUCC} \\
\\
\frac{e \Downarrow \mathbf{z}}{\mathbf{pred} \ e \Downarrow \mathbf{z}} \text{ B-PRED-ZERO} \qquad \frac{e \Downarrow \mathbf{succ} \ v}{\mathbf{pred} \ e \Downarrow v} \text{ B-PRED-SUCC} \\
\\
\frac{e \Downarrow \mathbf{z}}{\mathbf{iszero} \ e \Downarrow \mathbf{true}} \text{ B-ISZERO} \qquad \frac{e \Downarrow \mathbf{succ} \ v}{\mathbf{iszero} \ e \Downarrow \mathbf{false}} \text{ B-ISSUCC}
\end{array}$$

Prove that the small-step and big-step semantics for this language coincide, i.e. $e \Downarrow v$ iff $e \rightarrow^* v$. State and prove all necessary lemmas.

Hint: Read Exercise 3.5.17 in TAPL page 42 and the corresponding solution page 498.