

# Compiler Design

## Lecture 11: A Brief Tour of MIPS assembly

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# Overview

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# Assembly program template

## `.data`

Data segment: constant and variable definitions go here (including statically allocated arrays)

- format for declarations: `name: storage_type value`
- create storage for variable of specified type with given name and value
- `var1: .word 3 # one word of storage with initial value 3`
- `array1: .space 40 # 40 bytes of storage for array1`

## `.text`

Text segment: assembly instructions go here

# Components of an assembly program

Category	Example
Comment	<code># I am a comment</code>
Assembler directives	<code>.data, .ascii</code>
Operation mnemonic	<code>add, addi, lw, bne</code>
Register name	<code>\$zero, \$t3</code>
Address label (declaration)	<code>loop1:</code>
Address label (use)	<code>loop1</code>
Integer constant	<code>8, -4, 0xA9</code>
Character constant	<code>'h', '\t'</code>
String constant	<code>"Hello, world\n"</code>

# Hello world example

```
# Description: a simple hello world program
```

```
.data
```

```
hellostr: .asciiz "Hello, world\n"
```

```
.text
```

```
li $v0, 4           # setup print syscall  
la $a0, hellostr    # argument to print string  
syscall             # tell the OS to do the system call  
li $v0, 10          # setup exit syscall  
syscall             # tell the OS to perform the syscall
```

# Registers

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# Registers

- 32 general-purpose registers
- register preceded by \$ in assembly language
- two formats: name (\$zero) or number (\$0)
- holds 32 bits value (= 4 bytes = 1 word)
- stack grows from high memory to low memory



# Registers

Register number	Alternative name	Description
0	\$zero	the value 0
1	\$at	assembler temporary: reserved by the assembler
2-3	\$v0-\$v1	function return values
4-7	\$a0-\$a3	function arguments: first four parameters for function (no preserved across function call)
8-15	\$t0-\$t7	temporaries (not preserved across function calls)
16-23	\$s0-\$s7	saved temporaries (preserved across function calls)
24-25	\$t8-\$t9	temporaries: (not preserved across function calls)
26-27	\$k0-\$k1	reserved for use by the interrupt/trap handler
28	\$gp	global pointer : base of global data segment
29	\$sp	stack pointer : points to last location on stack
30	\$s8/\$fp	saved value / frame pointer (preserved across function call)
31	\$ra	return address

- Special Hi and Lo registers (not shown above) holds result of multiplication and division (see example later)

# Instructions

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# Instructions

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Arithmetic

# Arithmetic Instructions

- Most use three operands
- All operands are registers or immediate values (no memory access)
- All operands are 4 bytes (a word)

## Arithmetic & Move Instruction Examples

```
add    $t0,$t1,$t2
# $t0 = $t1 + $t2;
# add as signed (2's complement) integers

sub    $t2,$t3,$t4 # $t2 = $t3 - $t4
addi   $t2,$t3, 5  # $t2 = $t3 + 5;      "add immediate"
addu   $t1,$t6,$t7 # $t1 = $t6 + $t7;    add as unsigned integers
subu   $t1,$t6,$t7 # $t1 = $t6 + $t7;    subtract as unsigned integers
mul    $t1, $t2, $t3 # $t1 = $t2 * $t3; multiply

mult   $t3,$t4
# multiply 32-bit quantities in $t3 and $t4, and store 64-bit
# result in special registers Lo and Hi: (Hi,Lo) = $t3 * $t4

div    $t5,$t6
# Lo = $t5 / $t6 (integer quotient)
# Hi = $t5 mod $t6 (remainder)

mfhi   $t0
# move value from special register Hi to $t0: $t0 = Hi
mflo   $t1
# move value from special register Lo to $t1: $t1 = Lo

move   $t2,$t3    # $t2 = $t3
```

Instructions

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Memory

# Load / Store Instructions


- Memory access only allowed with explicit load and store instructions (load/store architecture)
- All other instructions use register operands
- Load
  - **lw** `register_destination, mem_source`  
copy a word (4 bytes) at source memory location to destination register
  - **lb** `register_destination, mem_source`  
copy a byte to low-order byte of destination register (sign extend higher-order bytes)
  - **li** `register_destination, value`  
load immediate value into destination register (pseudo-instruction)

- Store
  - **sw** **register\_source**, **mem\_destination**  
store a word (4 bytes) from source register to memory location
  - **sb** **register\_source**, **mem\_destination**  
store a byte (low-order) from source register to memory location

## Example

```
.data
var1: .word 23 # declare storage for var1; initial value is 23

.text
lw $t0, var1 # load content of mem location into register $t0: $t0 = 23
li $t1, 5    # $t1 = 5 ("load immediate")
sw $t1, var1 # store content of $t1 into mem: *var1 = 5
```

 **var1** represents a pointer to a word since it is an address.



# Indirect and Based Addressing

- load address:
  - `la $t0, var1`  
copy memory address of var1 into register \$t0
- indirect addressing:
  - `lw $t1, ($t0)`  
load word at memory address contained in \$t0 into \$t2
  - `sw $t2, ($t0)`  
store word in register \$t2 into memory at address contained in \$t0
- based/indexed addressing (useful for field access in struct):
  - `lw $t2, 4($t0)`  
load word at memory address ( $\$t0+4$ ) into register \$t2
  - `sw $t2, -12($t0)`  
store content of register \$t2 into memory at address ( $\$t0-12$ )

## Examples

`.data`

```
array1: .space 12 # declare 12 bytes of storage
```

`.text`

```
la $t0, array1 # load base address of array into $t0
```

```
li $t1, 5 # $t1 = 5 ("load immediate")
```

```
sw $t1, ($t0) # first array element set to 5
```

```
li $t1, 13 # $t1 = 13
```

```
sw $t1, 4($t0) # second array element set to 13
```

```
li $t1, -7 # $t1 = -7
```

```
sw $t1, 8($t0) # third array element set to -7
```

## Exercise

Write a MIPS assembly program corresponding to the following C code:

```
struct point_t {
    int x;
    int y;
};

struct point_t p;
int arr[12];

void foo() {
    p.x = 2;
    p.y = 4;
    arr[3] = 6;
}
```

# Instructions

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## Control Structures

# Control structures

- Branches:

```
b      target          # unconditional branch to target
beq   $t0,$t1,target # branch to target if $t0 = $t1
blt   $t0,$t1,target # branch to target if $t0 < $t1
ble   $t0,$t1,target # branch to target if $t0 <= $t1
bgt   $t0,$t1,target # branch to target if $t0 > $t1
bge   $t0,$t1,target # branch to target if $t0 >= $t1
bne   $t0,$t1,target # branch to target if $t0 <> $t1
```

## Example

```
addi $t0, $zero, 0 # t0 = 0
addi $t1, $zero, 10 # t1 = 10
```

loop:

```
addi $t0, $t0, 1 # t0 = t0+1
blt  $t0, $t1, loop # branch to loop if t0<t1 (t0<10)
```

# Control structures

- Jumps:

```
j          target  
# unconditional jump to program label target
```

```
jr        $t3  
# jump to address contained in $t3 ("jump register")
```

- Subroutine (function) call:

```
jal label # "jump and link"
```

- copy program counter (return address) to register \$ra (return address register)
- jump to program instruction at label

```
jr $ra # "jump register"
```

- jump to return address in \$ra (stored by jal instruction)

In case of nested function calls, the return address should be saved to the stack and restored accordingly.

# Instructions

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## System Calls

# System Calls (MIPS simulator)

System calls are used to interface with the operating systems. For instance input/output or dynamic memory allocation.

Using system calls:

1. load the service number in register `$v0`
2. load argument values in `$a0`, `$a1`, ...
3. issue the `syscall` instruction
4. retrieve return value if any

## Example: printing integer on the console

```
li $v0, 1
# service 1 is print integer

add $a0, $t0, $zero
# load desired value into argument register $a0

syscall
```



## System calls table

Service	\$v0	Arguments	Result
print integer	1	\$a0 = integer to print	
print string	4	\$a0 = address of null-terminated string to print	
print character	11	\$a0 = character to print	
read integer	5		\$v0 = integer read
read character	12		\$v0 = character read
allocate heap memory	9	\$a0 = number of bytes to allocate	\$v0 = address of allocated memory

Next lecture:

- Introduction to Code Generation